

Q1. They See Me Rolling (Search Problem)

Pacman buys a car to start Rolling in the Pac-City! But driving a car requires a new set of controls because he can now travel faster than 1 grid per turn (second). Instead of solely moving [North, South, East, West, Stop], Pacman's car has two distinct integers to control: throttle, and steering.

Throttle: $t_i \in \{1, 0, -1\}$, corresponding to {Gas, Coast, Brake}. This controls the **speed** of the car by determining its acceleration. The integer chosen here will be added to his velocity for the next state. For example, if Pacman is currently driving at 5 grid/s and chooses Gas he will be traveling at 6 grid/s in the next turn.

Steering: $s_i \in \{1, 0, -1\}$, corresponding to {Turn Left, Neutral, Turn Right}. This controls the **direction** of the car. For example, if he is facing North and chooses Turn Left he will be facing West in the next turn.

(a) Suppose Pac-city has dimension m by n , but only $k < mn$ squares are legal roads. The speed limit of Pac-city is 3 grid/s. For this sub-part only, suppose Pacman is a law-abiding citizen, so $0 \leq v \leq 3$ at all time, and he only drives on legal roads.

(i) Without any additional information, what is the **tightest upper bound** on the size of state space, if he wants to search a route (not necessarily the shortest) from his current location to anywhere in the city. Please note that your state space representation must be able to represent **all** states in the search space.

- $4mn$
 $4k$
 $12mn$
 $12k$
 $16mn$
 $16k$
 $48mn$
 $48k$

(ii) What is the maximum branching factor? The answer should be an integer.

(iii) Which algorithm(s) is/are guaranteed to return a path between two points, if one exists?

- Depth First Tree Search Breadth First Tree Search
 Depth First Graph Search Breadth First Graph Search

(iv) Is Breadth First Graph Search guaranteed to return the path with the shortest grid distance?

- Yes No

(b) Now let's remove the constraint that Pacman follows the speed limit. Now Pacman's speed is limited by the mechanical constraints of the car, which is 6 grid/s, double the speed limit.

Pacman is now able to drive twice as fast on the route to his destination. How do the following properties of the search problem change as a result of being able to drive twice as fast?

(i) Size of State Space:

- Increases Stays the same Decreases

(ii) Maximum Branching Factor:

- Increases Stays the same Decreases

For the following part, **mark all choices that could happen on any graph**

(iii) The number of nodes expanded with **Depth** First Graph Search:

- Increases Stays the same Decreases

- (c) Now we need to consider that there are $p > 0$ police cars waiting at $p > 0$ distinct locations trying to catch Pacman riding dirty!! All police cars are stationary, but once Pacman takes an action which lands him in the same grid as one police car, Pacman will be arrested and the game ends.

Pacman wants to find a route to his destination, without being arrested. How do the following properties of the search problem change as a result of avoiding the police?

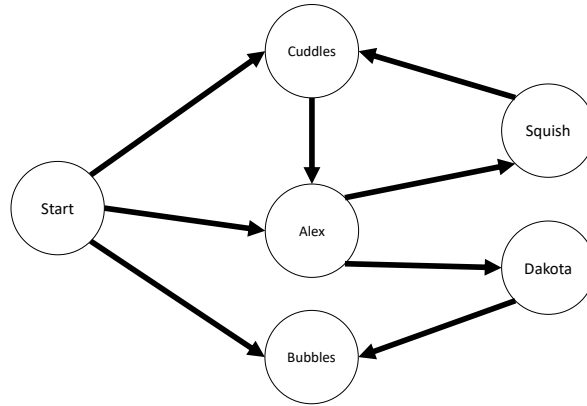
- (i) Size of State Space:
 Increases Stays the same Decreases
- (ii) Maximum Branching Factor:
 Increases Stays the same Decreases

For the following part, **mark all choices that could happen on any graph**

- (iii) Number of nodes expanded with **Breadth** First Graph Search:
 Increases Stays the same Decreases

Q2. Search: Snail search for love

Scorpborg the snail is looking for a mate. It can visit different potential mates based on a trail of ooze to nearby snails, and then test them for chemistry, as represented in the below graph, where each node represents a snail. In all cases, nodes with equal priority should be visited in alphabetical order.



(a) Simple search

In this part, assume that the only match for Scorpborg is Squish (i.e. Squish is the goal state). Which of the following are true **when searching the above graph**?

- (i) BFS Tree Search expands more nodes than DFS Tree Search True False
- (ii) DFS Tree Search finds a path to the goal for this graph True False
- (iii) DFS Graph Search finds the shortest path to the goal for this graph True False
- (iv) If we remove the connection from Cuddles \rightarrow Alex, can DFS Graph Search find a path to the goal for the altered graph? Yes No

(b) Third Time's A Charm

Now we assume that Scorpborg's mate preferences have changed. The new criteria she is looking for in a mate is that she has **visited the mate twice before** (i.e. when she visits any state for the third time, she has found a path to the goal).

- (i) What should the most simple yet sufficient new state space representation include?
 - The current location of Scorpborg
 - The total number of edges travelled so far
 - An array of booleans indicating whether each snail has been visited so far
 - An array of numbers indicating how many times each snail has been visited so far
 - The number of distinct snails visited so far
- (ii) DFS Tree Search finds a path to the goal for this graph True False
- (iii) BFS Graph Search finds a path to the goal for this graph True False
- (iv) If we remove the connection from Cuddles \rightarrow Alex, can DFS Graph Search find a path to the goal for the altered graph? Yes No

We continue as in part (b) where the goal is still to find a mate who is visited for the third time.

(c) Costs for visiting snails

Assume we are using Uniform cost search and we can now add costs to the actions in the graph.

- (i) Can one assign (non-negative) costs to the actions in the graph such that the goal state returned by UCS (Tree-search) changes? Yes No
- (ii) Can one assign (potentially negative) costs to the actions in the graph such that UCS (Tree-search) will never find a goal state? Yes No

Q3. Power Pellets

Consider a Pacman game where Pacman can eat 3 types of pellets:

- Normal pellets (n-pellets), which are worth one point.
- Decaying pellets (d-pellets), which are worth $\max(0, 5 - t)$ points, where t is time.
- Growing pellets (g-pellets), which are worth t points, where t is time.

The location and type of each pellet is fixed. The pellet's point value stops changing once eaten. For example, if Pacman eats one g-pellet at $t = 1$ and one d-pellet at $t = 2$, Pacman will have won $1 + 3 = 4$ points.

Pacman needs to find a path to win at least 10 points but he wants to minimize distance travelled. The cost between states is equal to distance travelled.

(a) Which of the following must be including for a minimum, sufficient state space?

- Pacman's location
- Location and type of each pellet
- How far Pacman has travelled
- Current time
- How many pellets Pacman has eaten and the point value of each eaten pellet
- Total points Pacman has won
- Which pellets Pacman has eaten

(b) Which of the following are admissible heuristics? Let x be the number of points won so far.

- Distance to closest pellet, except if in the goal state, in which case the heuristic value is 0.
- Distance needed to win $10 - x$ points, determining the value of all pellets as if they were n-pellets.
- Distance needed to win $10 - x$ points, determining the value of all pellets as if they were g-pellets (i.e. all pellet values will be t .)
- Distance needed to win $10 - x$ points, determining the value of all pellets as if they were d-pellets (i.e. all pellet values will be $\max(0, 5 - t)$.)
- Distance needed to win $10 - x$ points assuming all pellets maintain current point value (g-pellets stop increasing in value and d-pellets stop decreasing in value)
- None of the above

(c) Instead of finding a path which minimizes distance, Pacman would like to find a path which minimizes the following:

$$C_{new} = a * t + b * d$$

where t is the amount of time elapsed, d is the distance travelled, and a and b are non-negative constants such that $a + b = 1$. Pacman knows an admissible heuristic when he is trying to minimize time (i.e. when $a = 1, b = 0$), h_t , and when he is trying to minimize distance, h_d (i.e. when $a = 0, b = 1$).

Which of the following heuristics is guaranteed to be admissible when minimizing C_{new} ?

- $mean(h_t, h_d)$
- $min(h_t, h_d)$
- $max(h_t, h_d)$
- $a * h_t + b * h_d$
- None of the above