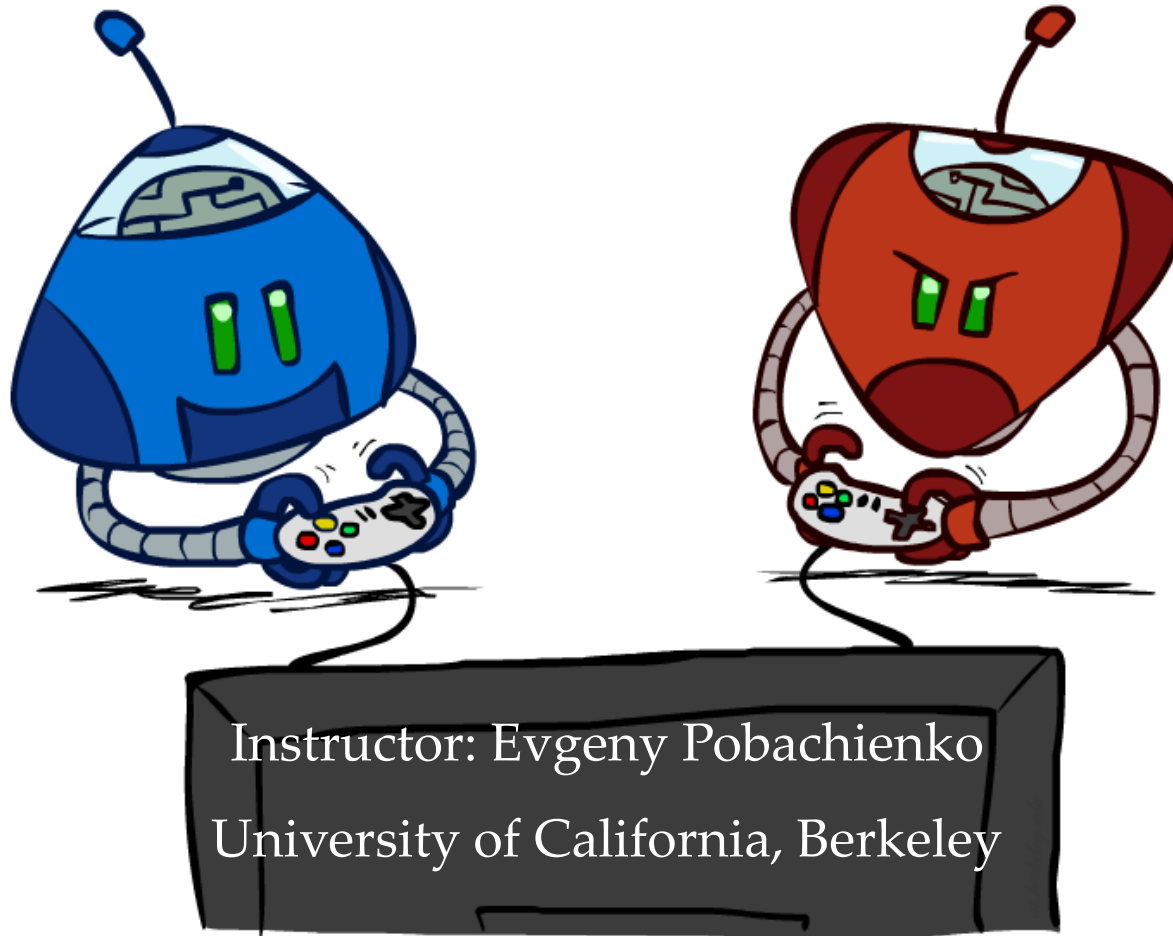


CS 188: Artificial Intelligence

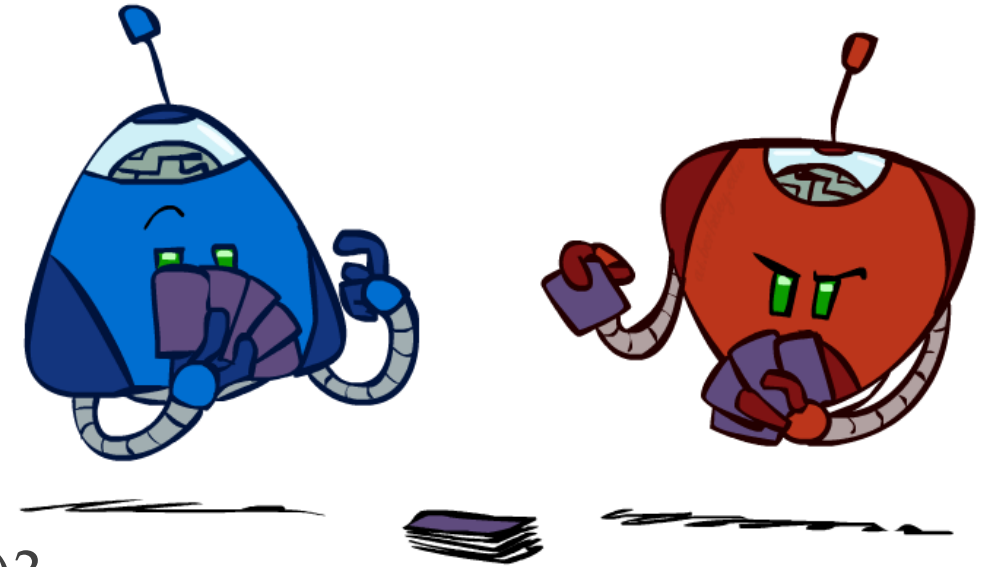
Search with Other Agents



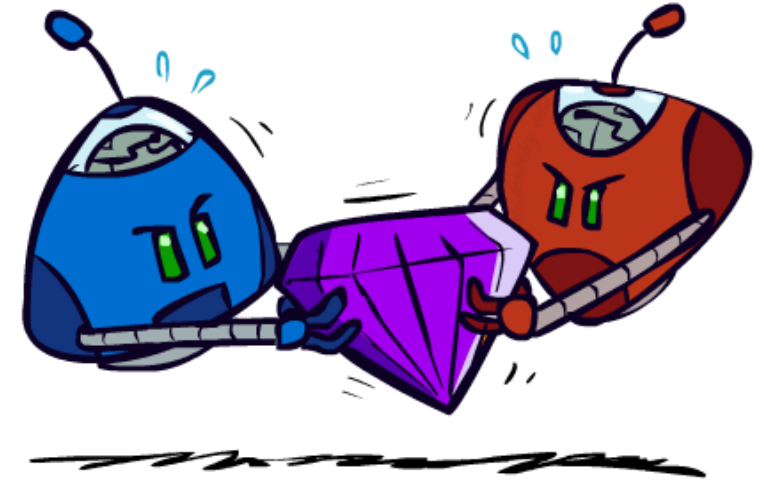
Instructor: Evgeny Pobachienko
University of California, Berkeley

Types of Games

- Many different kinds of games!
- Axes:
 - Deterministic or stochastic?
 - One, two, or more players?
 - Zero sum?
 - Perfect information (can you see the state)?



Types of Games



○ General Games

- Agents have independent utilities (values on outcomes)
- Cooperation, indifference, competition, and more are all possible
 - We don't make AI to act in isolation, it should a) work around people and b) help people
 - That means that every AI agent needs to solve a game

○ Zero-Sum Games

- Agents have opposite utilities (values on outcomes)
- Lets us think of a single value that one maximizes and the other minimizes
- Adversarial, pure competition

Zero-Sum Games 😊

- **Checkers**

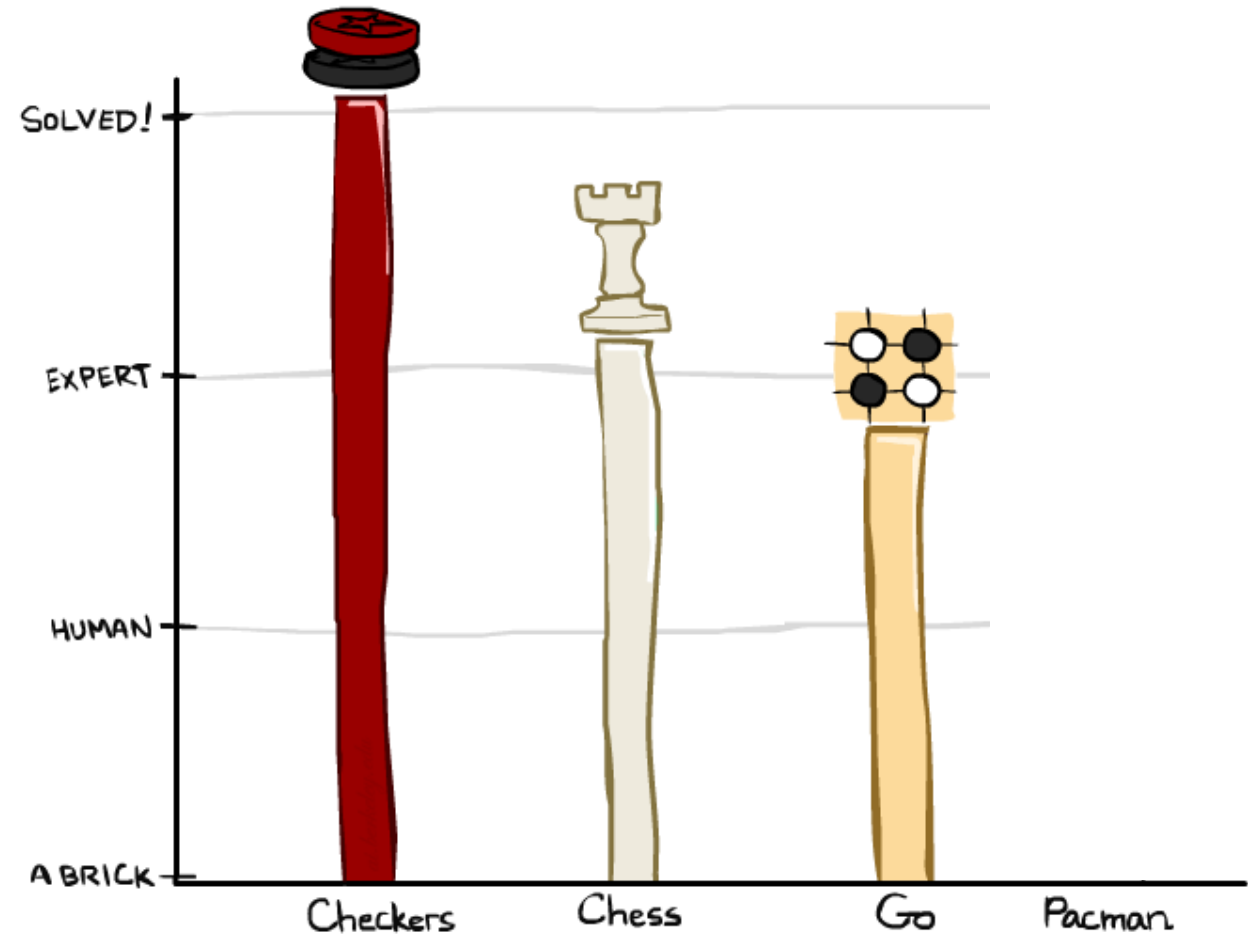
- (1950): First computer player.
- (1994): First computer champion: Chinook ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame.
- (2007): Checkers solved!

- **Chess**

- (1997): Deep Blue defeats human champion Gary Kasparov in a six-game match. Current programs are even better, if less historic.

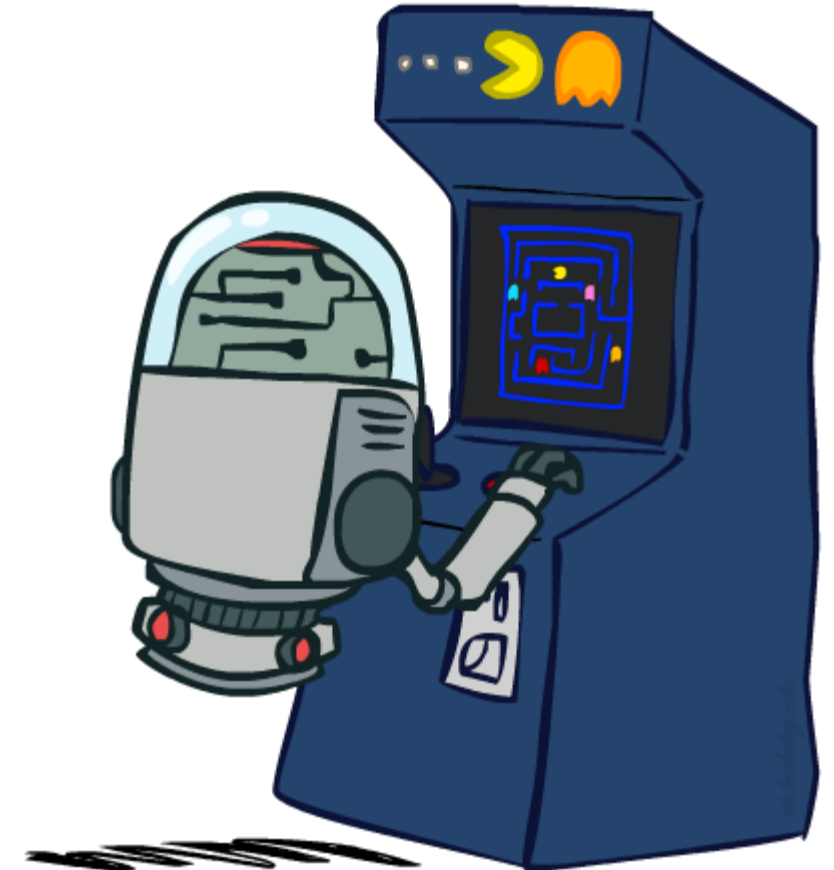
- **Go**

- (2016): AlphaGo defeats human champion Lee Sedol. Uses Monte Carlo Tree Search, learned evaluation function.

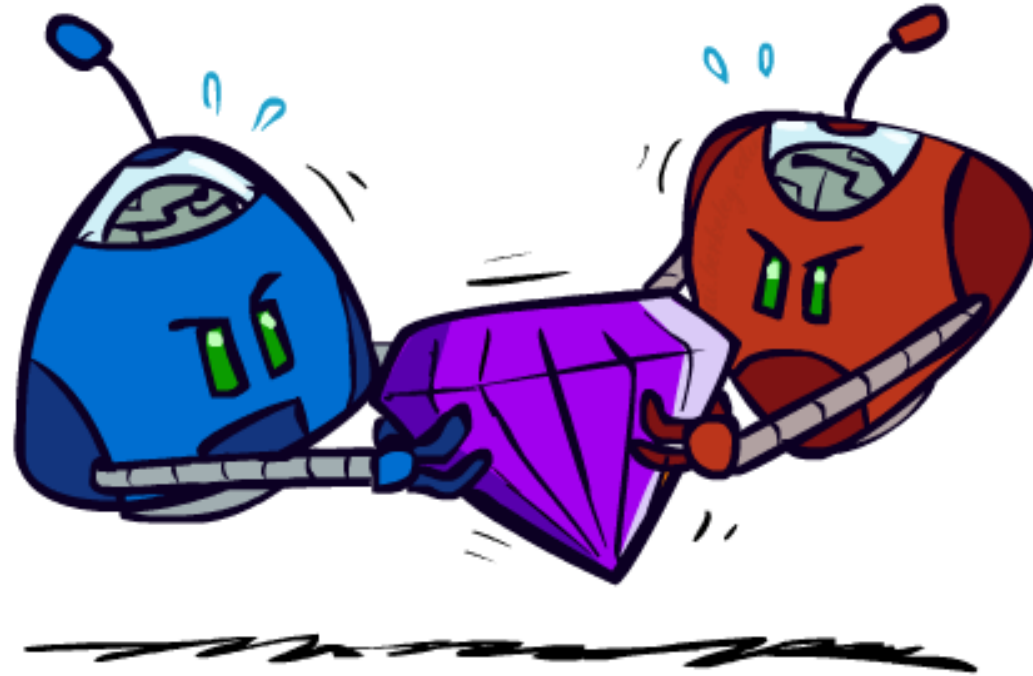


Deterministic Games with Terminal Utilities

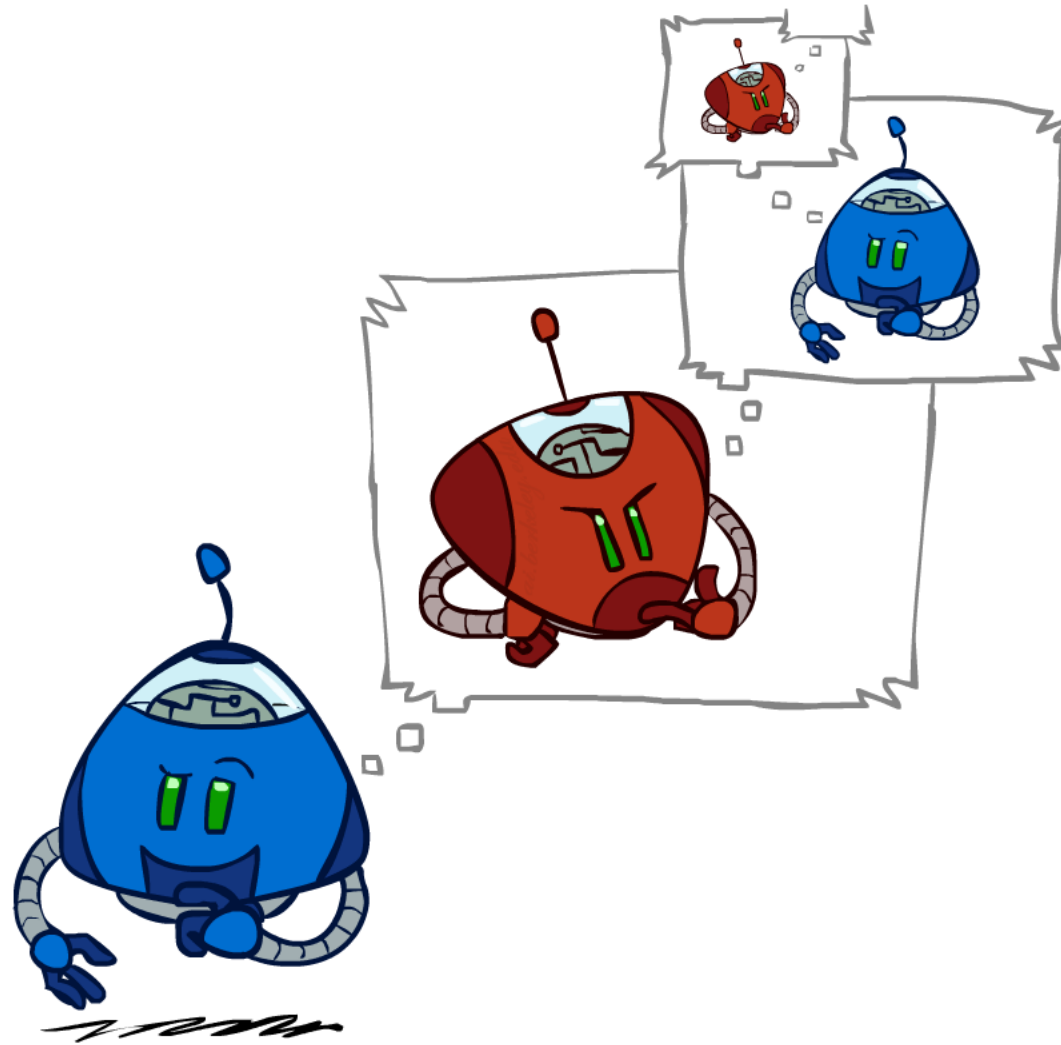
- Many possible formalizations, one is:
 - States: S (start at s_0)
 - Players: $P = \{1 \dots N\}$ (usually take turns)
 - Actions: A (may depend on player / state)
 - Transition Function: $S \times A \rightarrow S$
 - Terminal Test: $S \rightarrow \{t, f\}$
 - Terminal Utilities: $S \times P \rightarrow \mathbb{R}$
- Solution for a player is a **policy**: $S \rightarrow A$



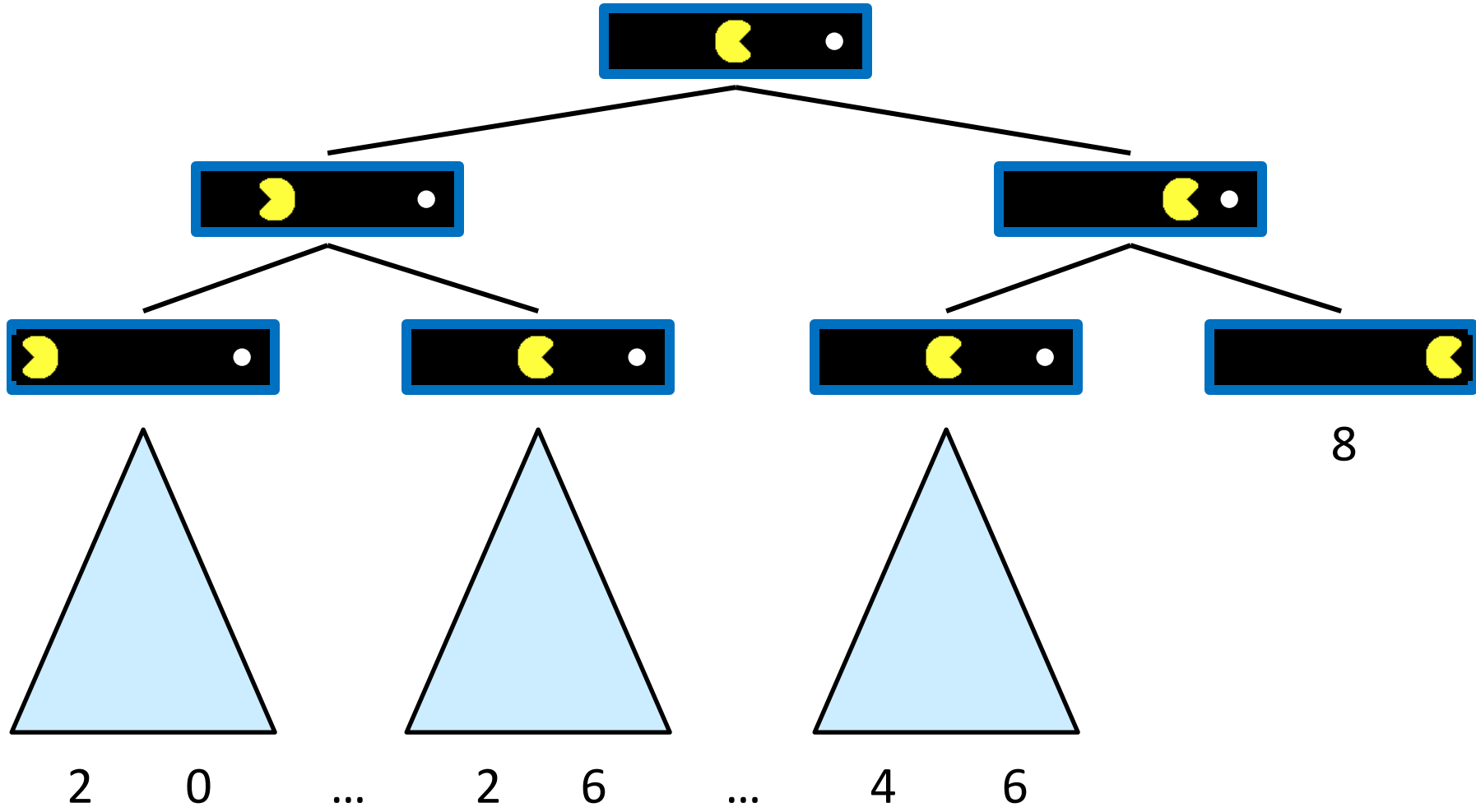
Adversarial Games



Adversarial Search

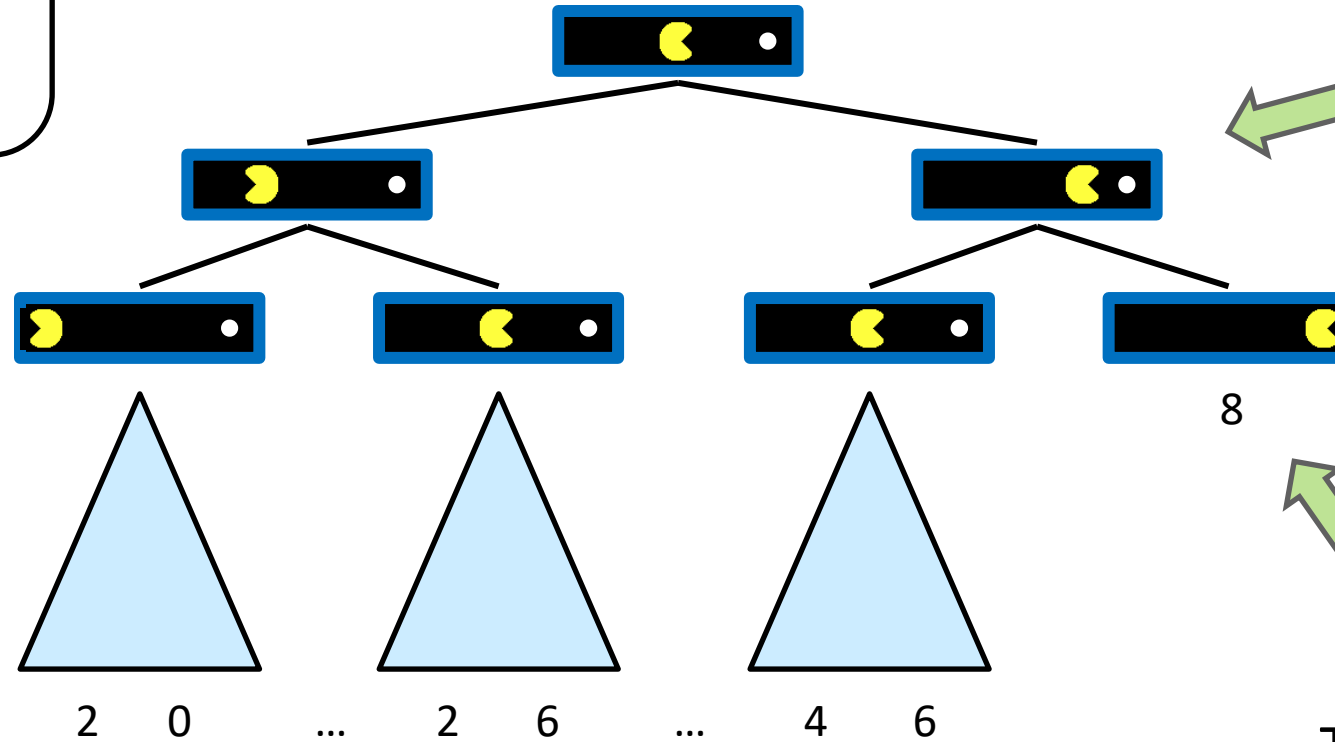


Single-Agent Trees



Value of a State

Value of a state:
The best achievable
outcome (utility)
from that state



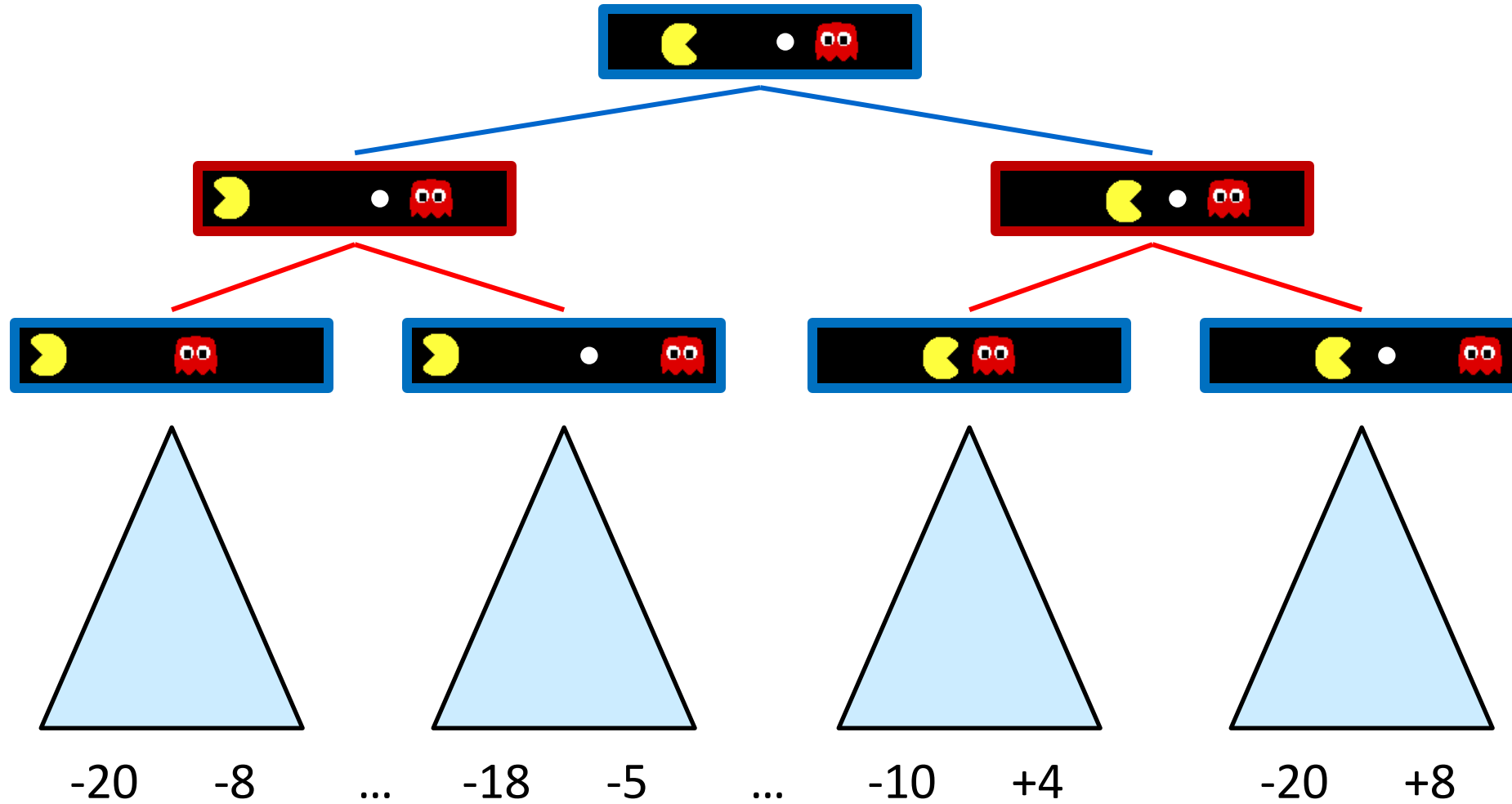
Non-Terminal States:

$$V(s) = \max_{s' \in \text{children}(s)} V(s')$$

Terminal States:

$$V(s) = \text{known}$$

Adversarial Game Trees



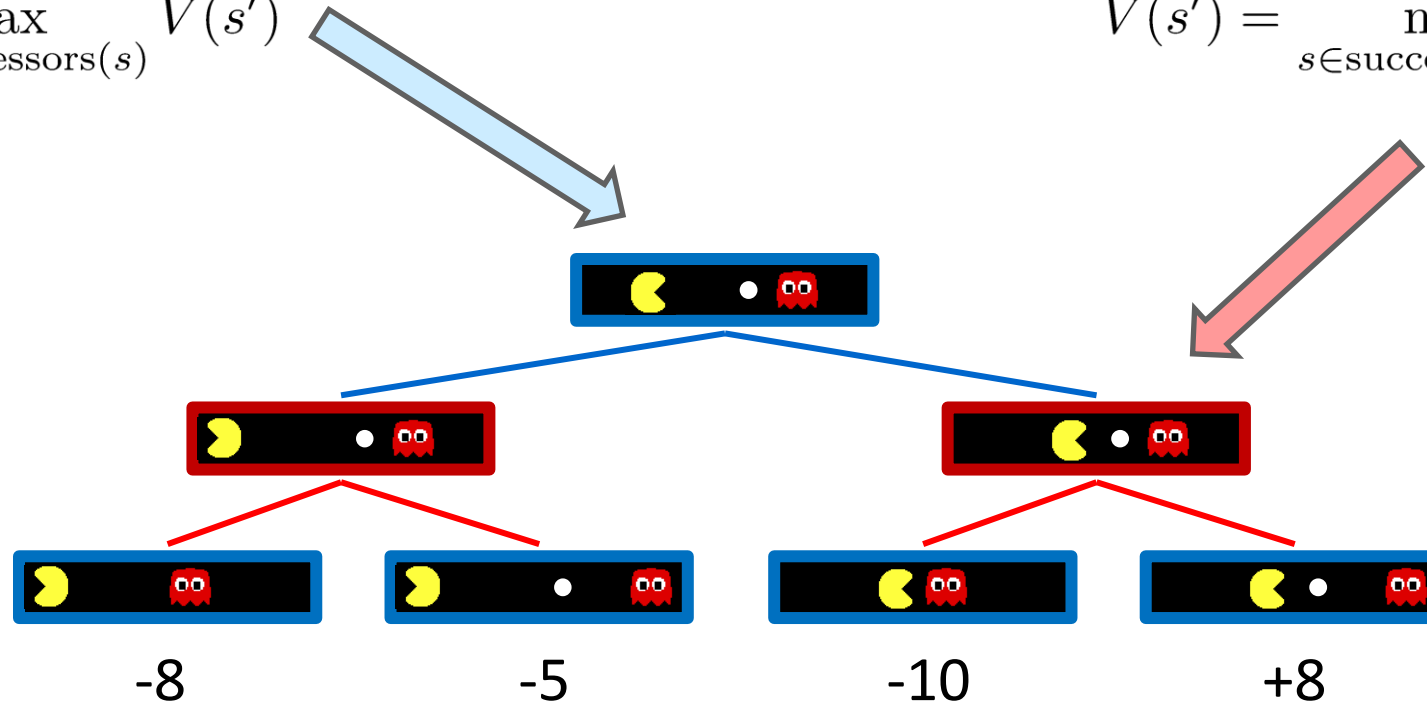
Minimax Values

States Under Agent's Control:

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$

States Under Opponent's Control:

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$



Terminal States:

$$V(s) = \text{known}$$

Tic-Tac-Toe Game Tree



MAX (X)



MIN (O)



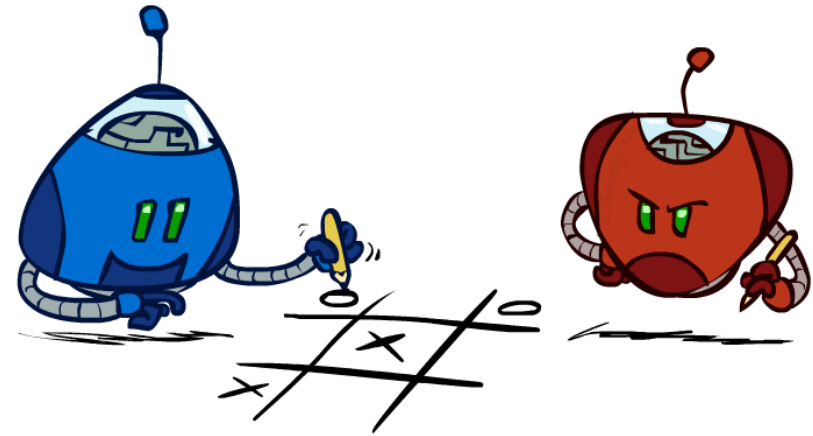
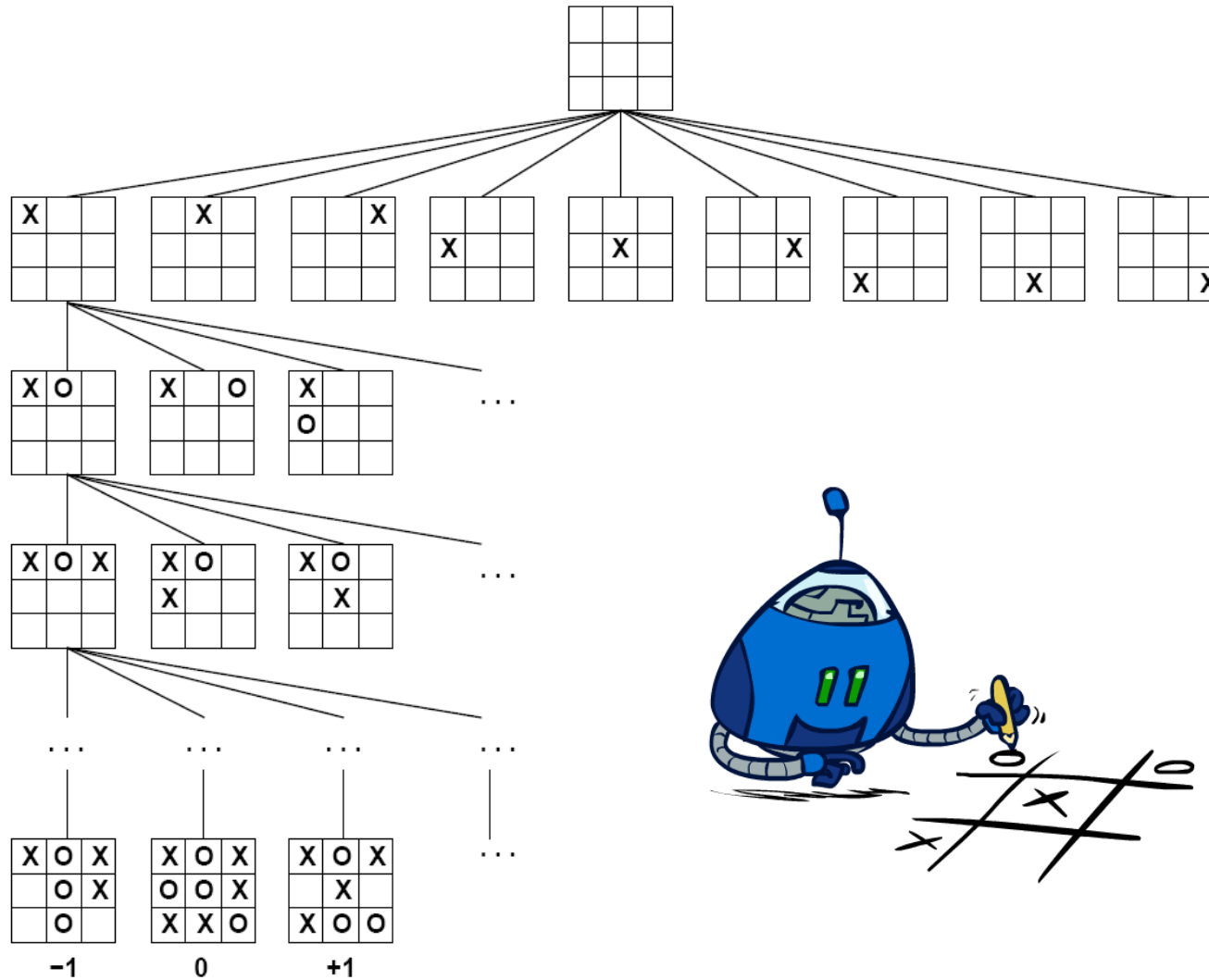
MAX (X)



MIN (O)

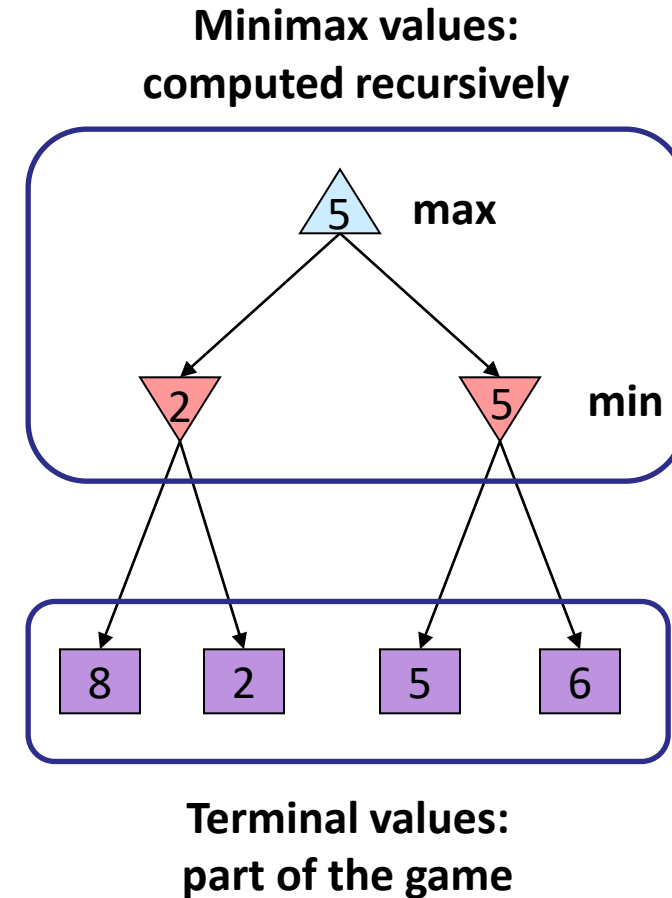
TERMINAL

Utility



Adversarial Search (Minimax)

- Deterministic, zero-sum games:
 - Tic-tac-toe, chess, checkers
 - One player maximizes result
 - The other minimizes result
- Minimax search:
 - A state-space search tree
 - Players alternate turns
 - Compute each node's **minimax value**: the best achievable utility against a rational (optimal) adversary



Minimax Implementation (Dispatch)

```
def value(state):
```

```
    if the state is terminal: return the state's utility
```

```
    if the next agent is MAX: return max-value(state)
```

```
    if the next agent is MIN: return min-value(state)
```

```
def max-value(state):
```

```
    initialize v =  $-\infty$ 
```

```
    for each successor of state:
```

```
        v = max(v, value(successor))
```

```
    return v
```

```
def min-value(state):
```

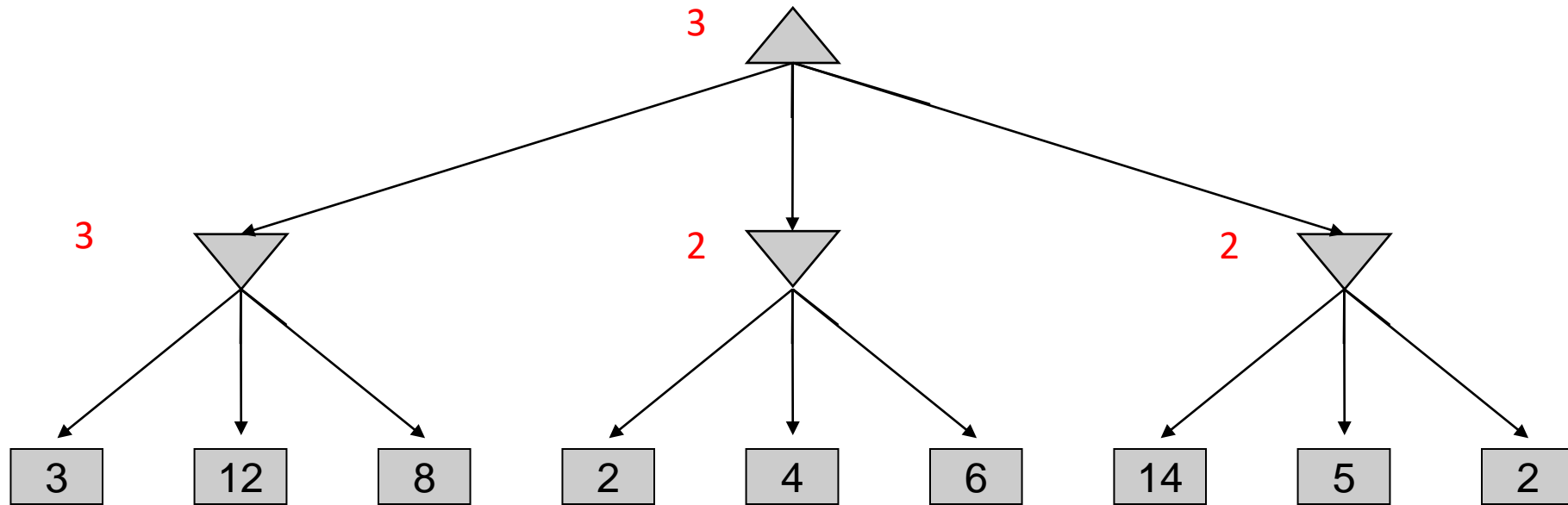
```
    initialize v =  $+\infty$ 
```

```
    for each successor of state:
```

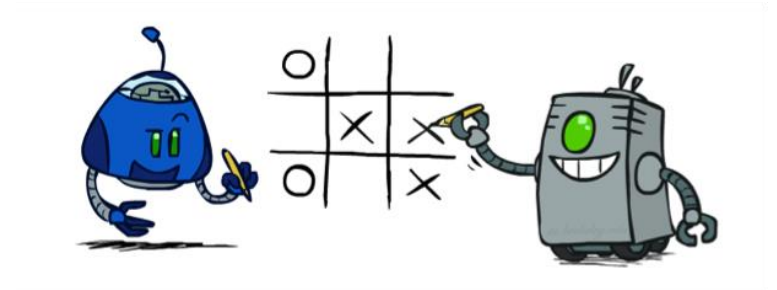
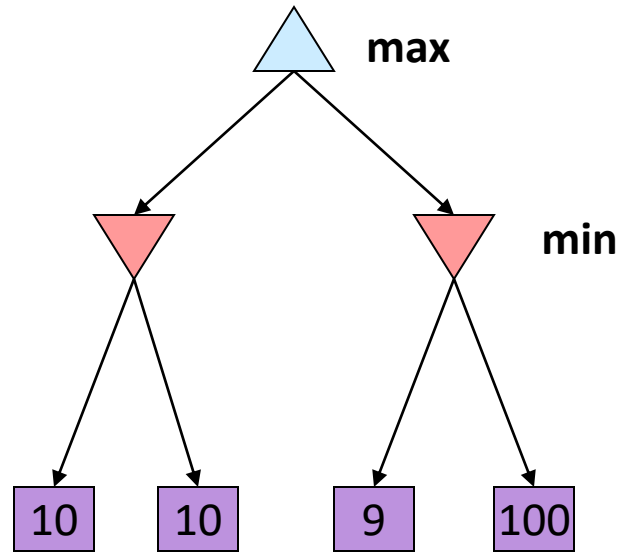
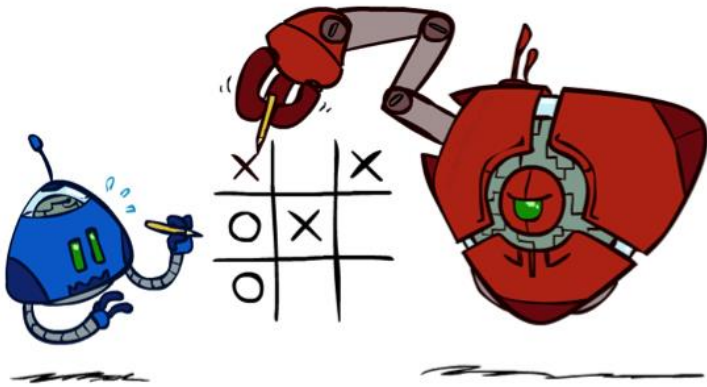
```
        v = min(v, value(successor))
```

```
    return v
```

Minimax Example



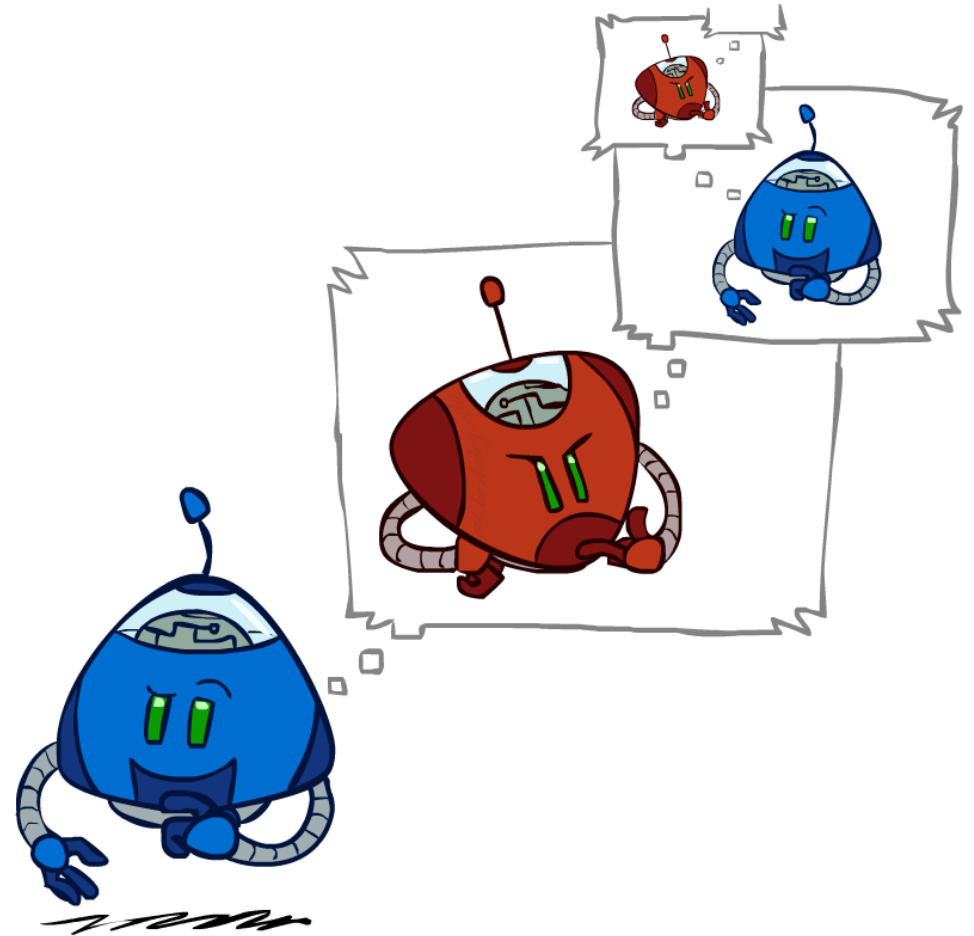
Minimax Properties



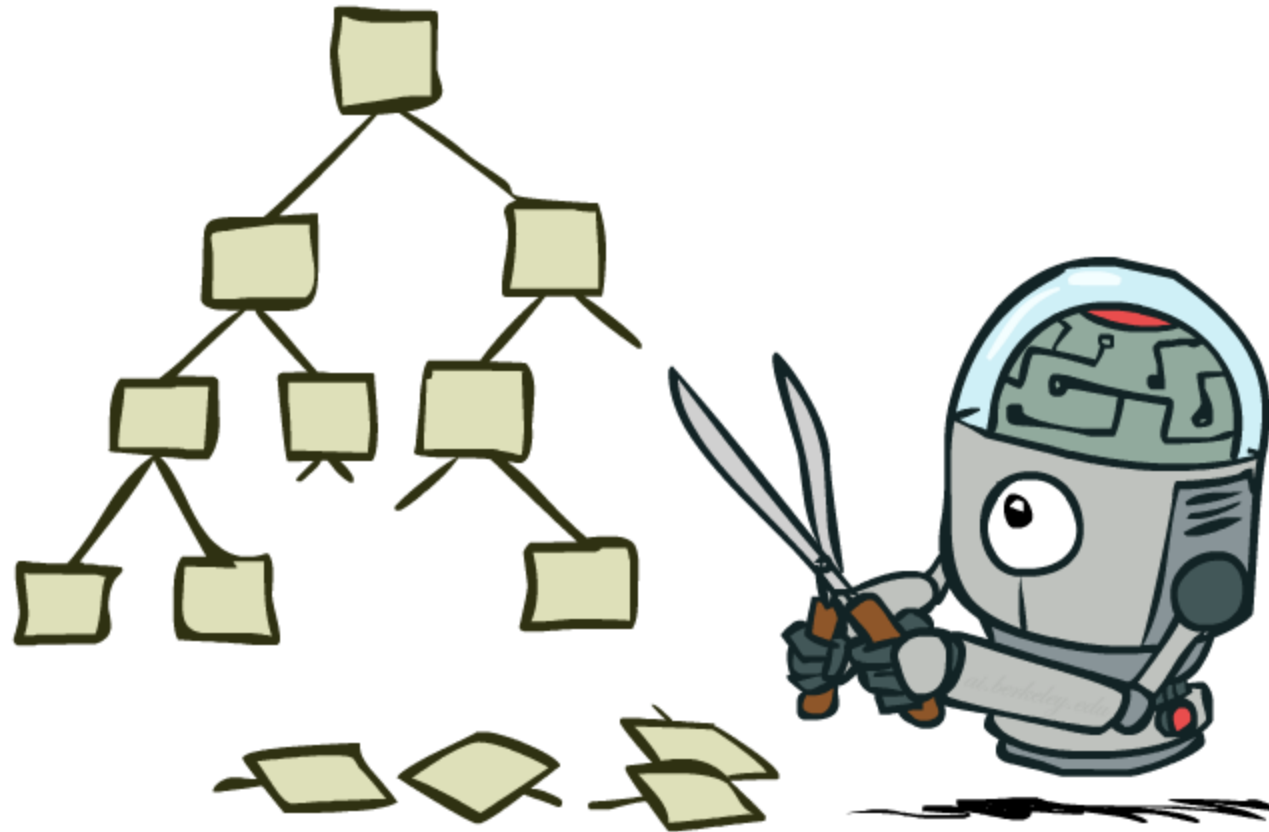
Optimal against a perfect player. Otherwise?

Minimax Efficiency

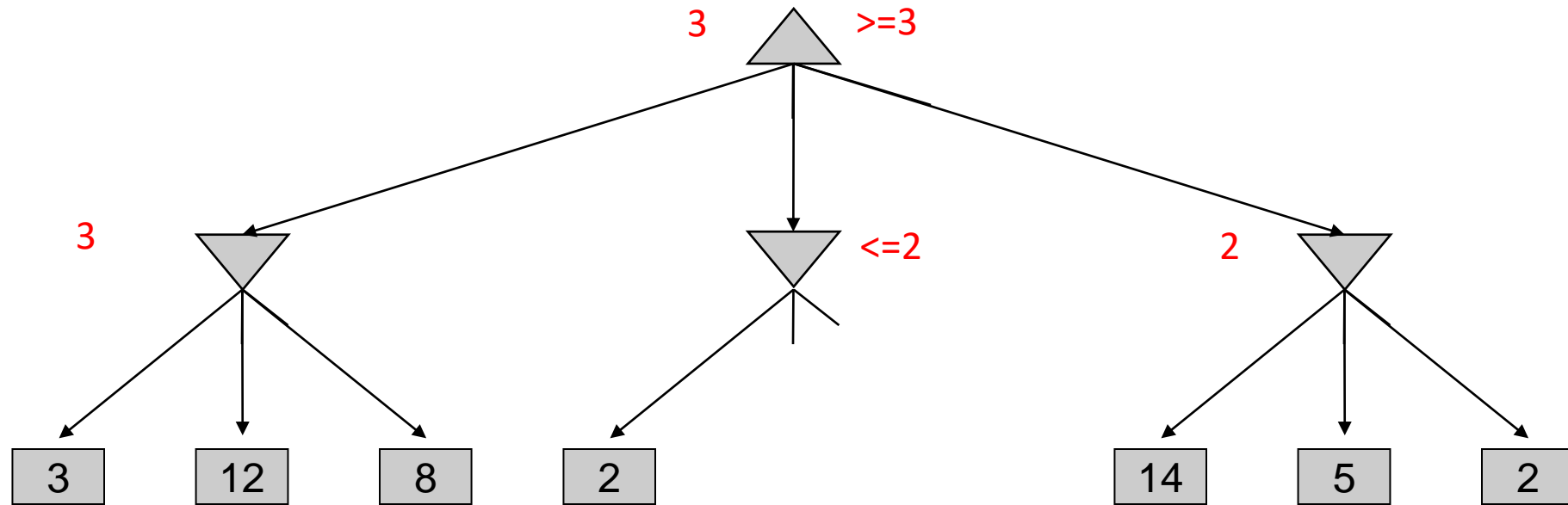
- How efficient is minimax?
 - Just like (exhaustive) DFS
 - Time: $O(b^m)$
 - Space: $O(bm)$
- Example: For chess, $b \approx 35$, $m \approx 100$
 - Exact solution is completely infeasible
 - But, do we need to explore the whole tree?



Game Tree Pruning



Minimax Example: Metareasoning



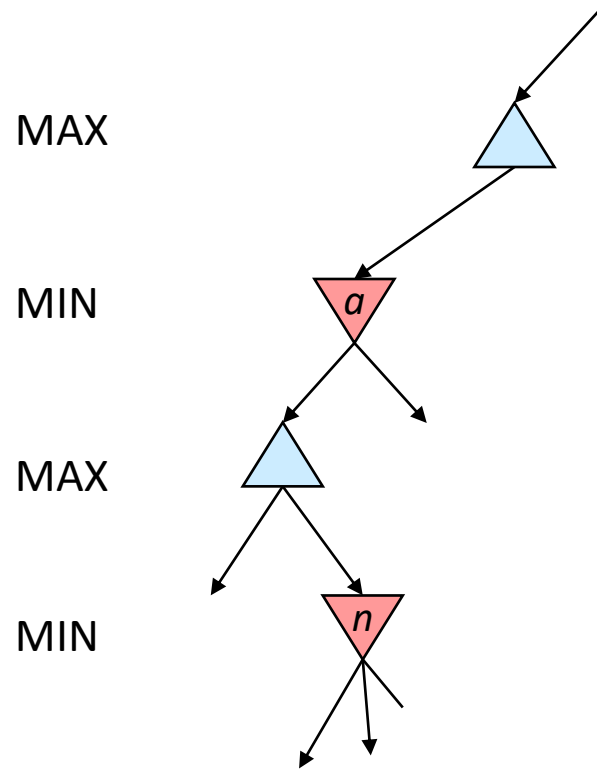
Alpha-Beta Implementation

α : MAX's best option on path to root
 β : MIN's best option on path to root

```
def max-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = -\infty$   
    for each successor of state:  
         $v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \geq \beta$  return  $v$   
         $\alpha = \max(\alpha, v)$   
    return  $v$ 
```

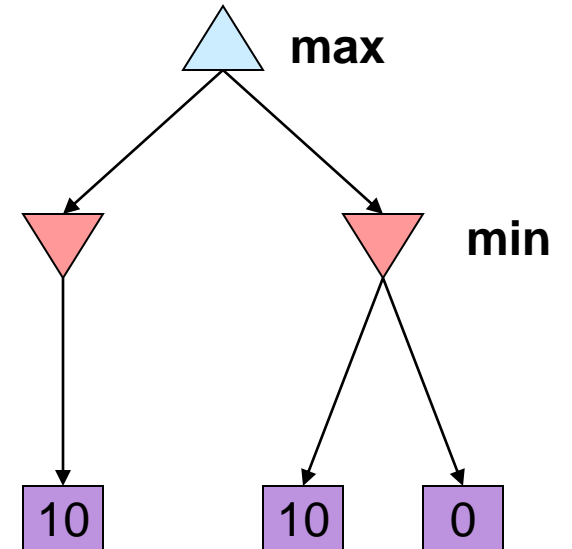
```
def min-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = +\infty$   
    for each successor of state:  
         $v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \leq \alpha$  return  $v$   
         $\beta = \min(\beta, v)$   
    return  $v$ 
```

Why on Path?

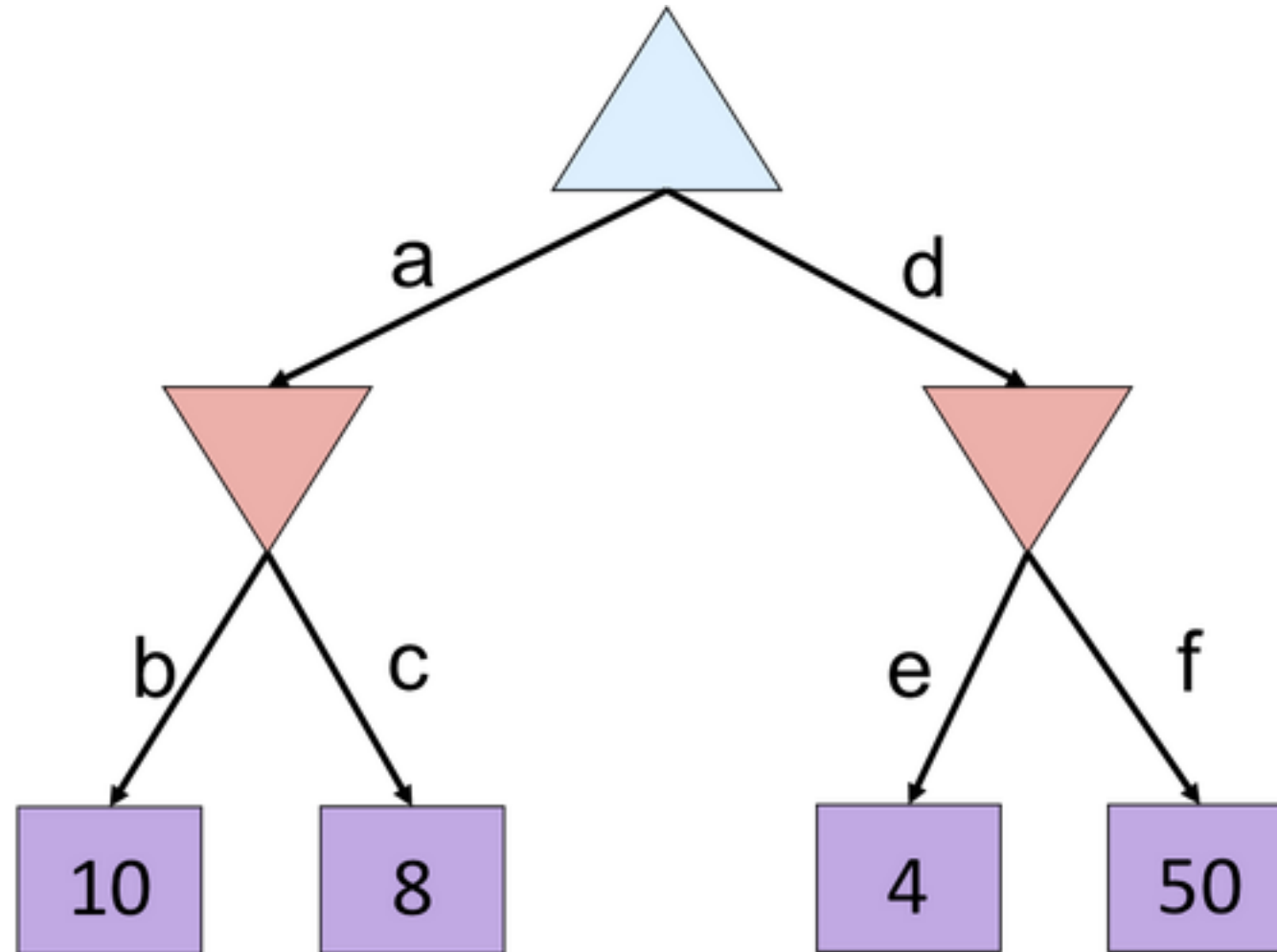


Alpha-Beta Pruning Properties

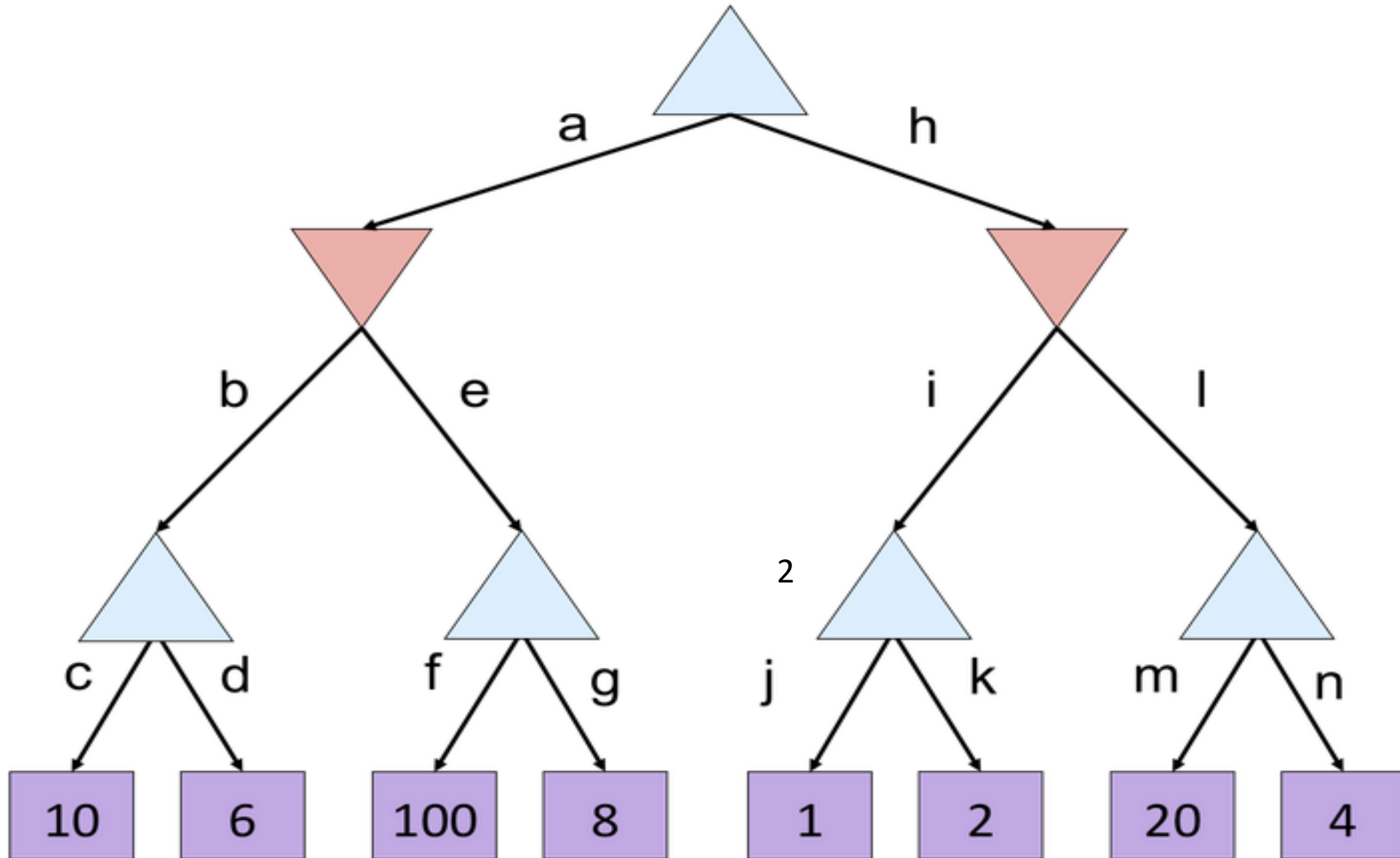
- This pruning has **no effect** on minimax value computed for the root!
- Values of intermediate nodes might be wrong
 - Important: children of the root may have the wrong value
 - So the most naïve version won't let you do action selection
- Good child ordering improves effectiveness of pruning
- With “perfect ordering”:
 - Time complexity drops to $O(b^{m/2})$
 - Doubles solvable depth!
 - Full search of, e.g. chess, is still hopeless...
- This is a simple example of **metareasoning** (computing about what to compute)



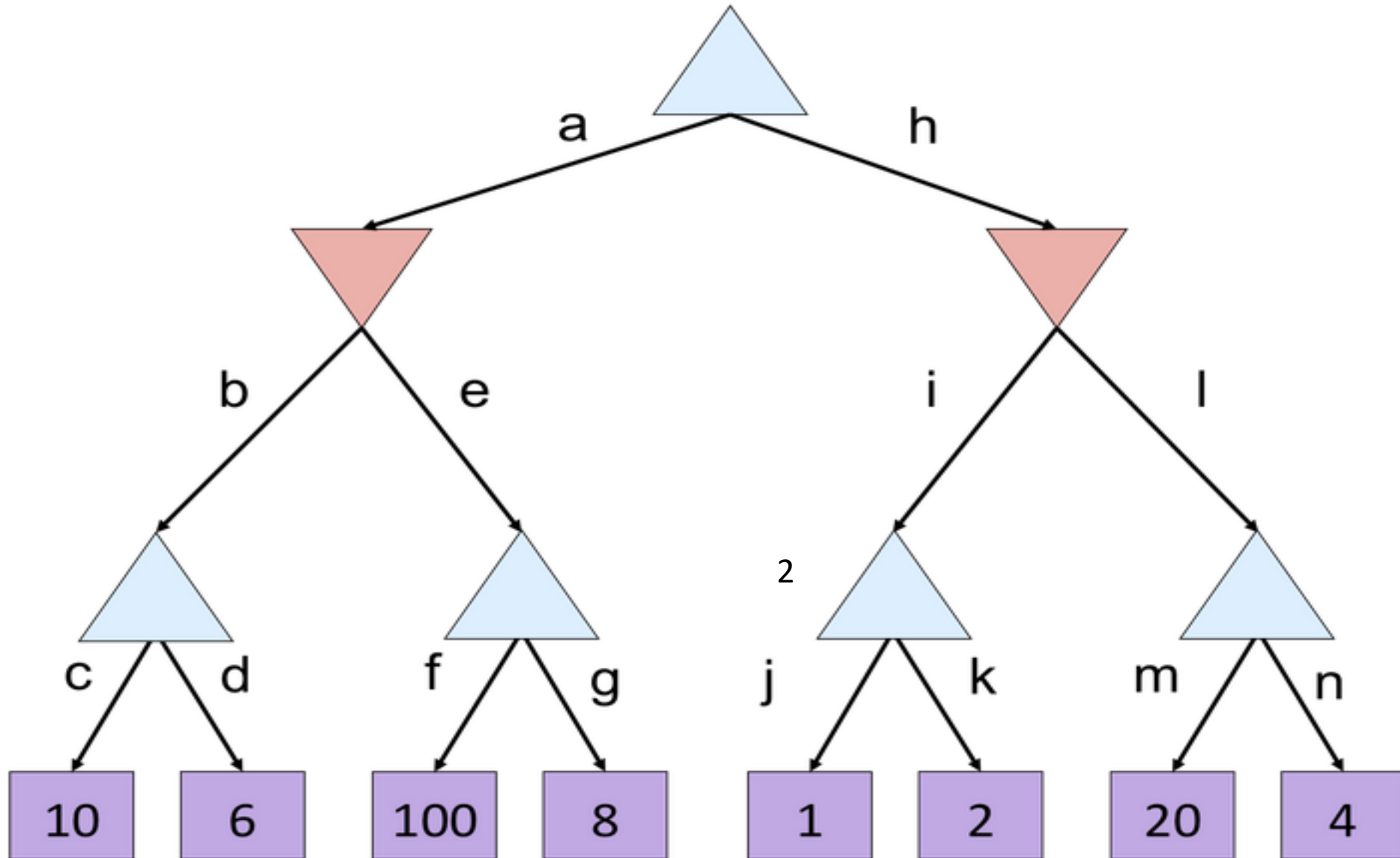
Alpha-Beta Quiz



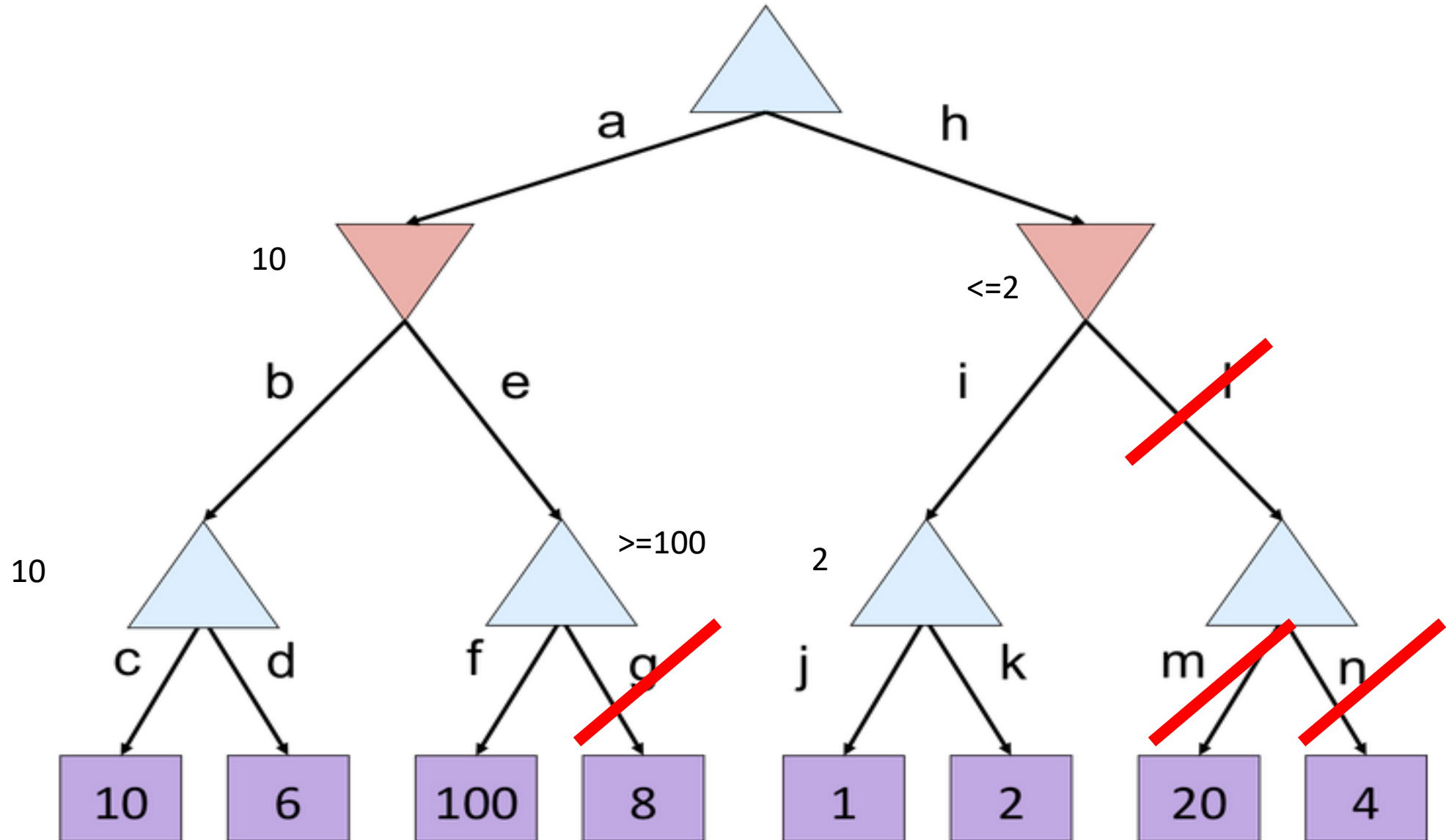
Alpha-Beta Quiz 2



Alpha-Beta Quiz 2



Alpha-Beta Quiz 2

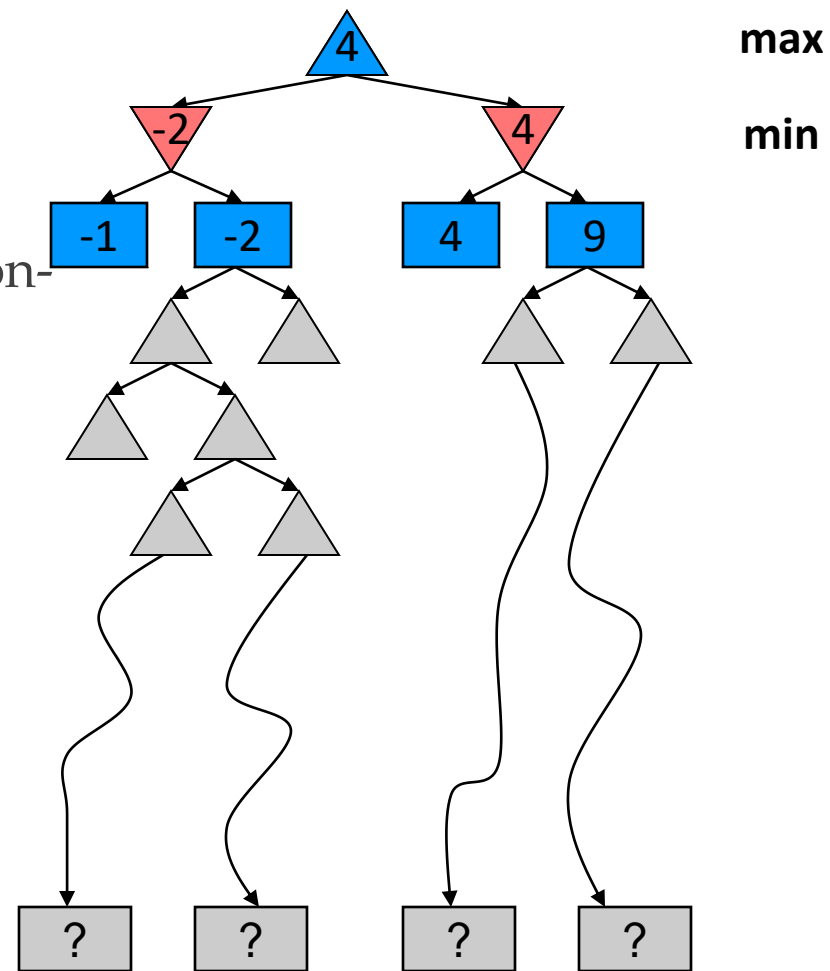


Resource Limits



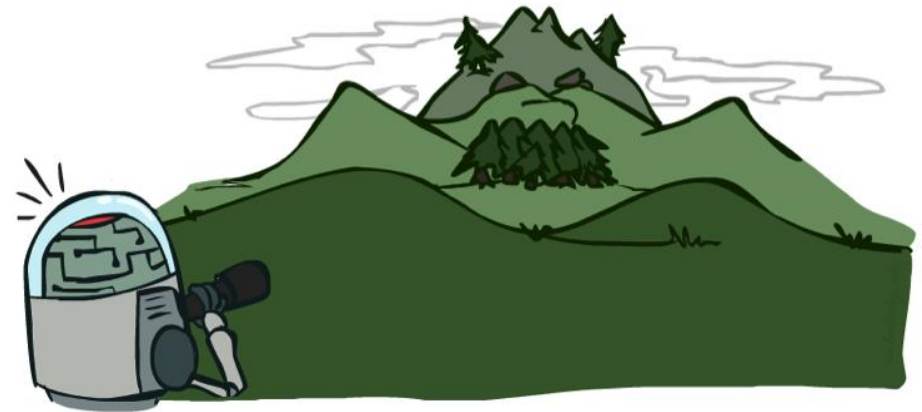
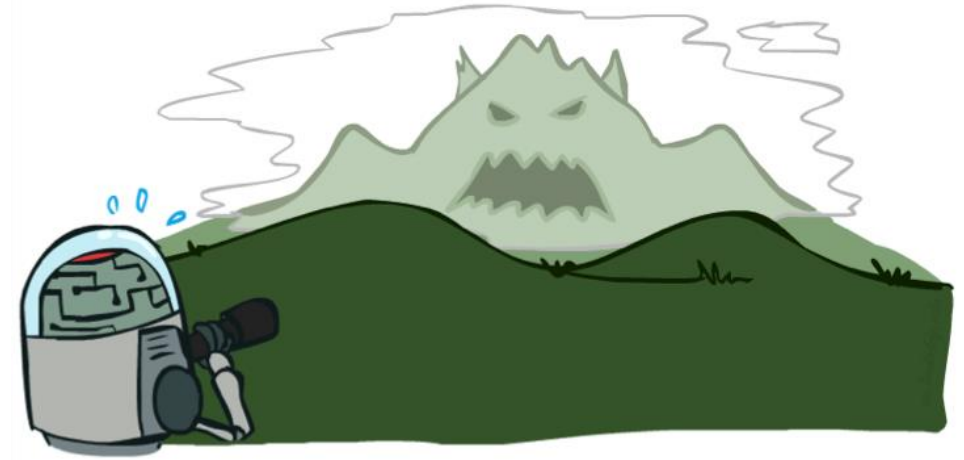
Resource Limits

- Problem: In realistic games, cannot search to leaves!
- Solution: Depth-limited search
 - Instead, search only to a limited depth in the tree
 - Replace terminal utilities with **an evaluation function** for non-terminal positions
- Example:
 - Suppose we have 100 seconds, can explore 10K nodes / sec
 - So can check 1M nodes per move
 - α - β reaches about depth 8 – decent chess program
- Guarantee of optimal play is gone
- More plies makes a BIG difference
- Use iterative deepening for an anytime algorithm

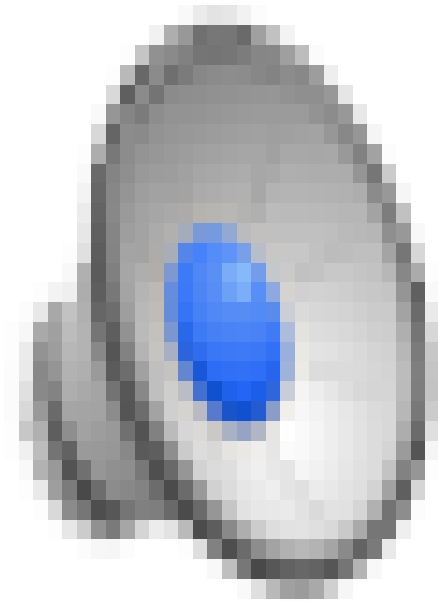


Depth Matters

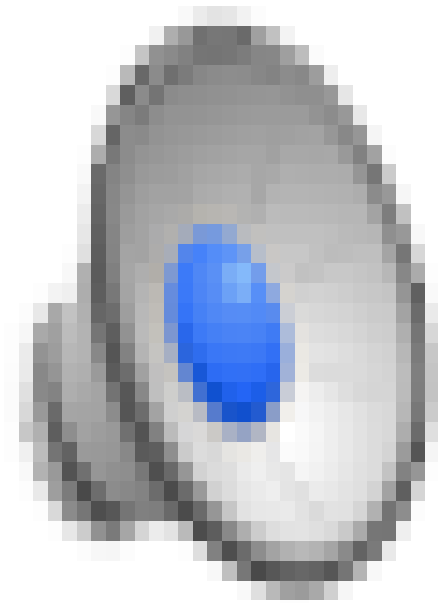
- Evaluation functions are always imperfect
- The deeper in the tree the evaluation function is buried, the less the quality of the evaluation function matters
- An important example of the tradeoff between complexity of features and complexity of computation



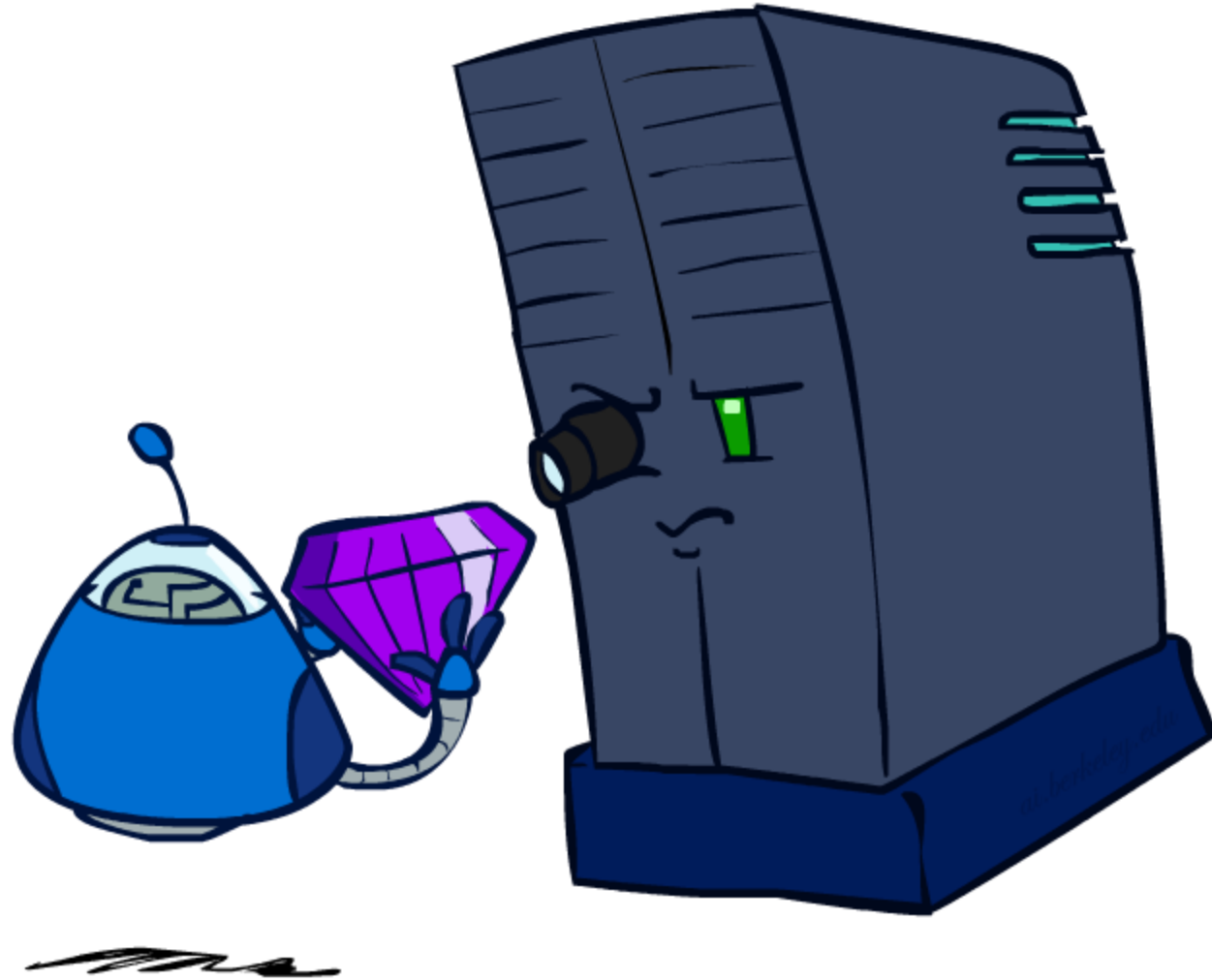
Video of Demo Limited Depth (2)



Video of Demo Limited Depth (10)

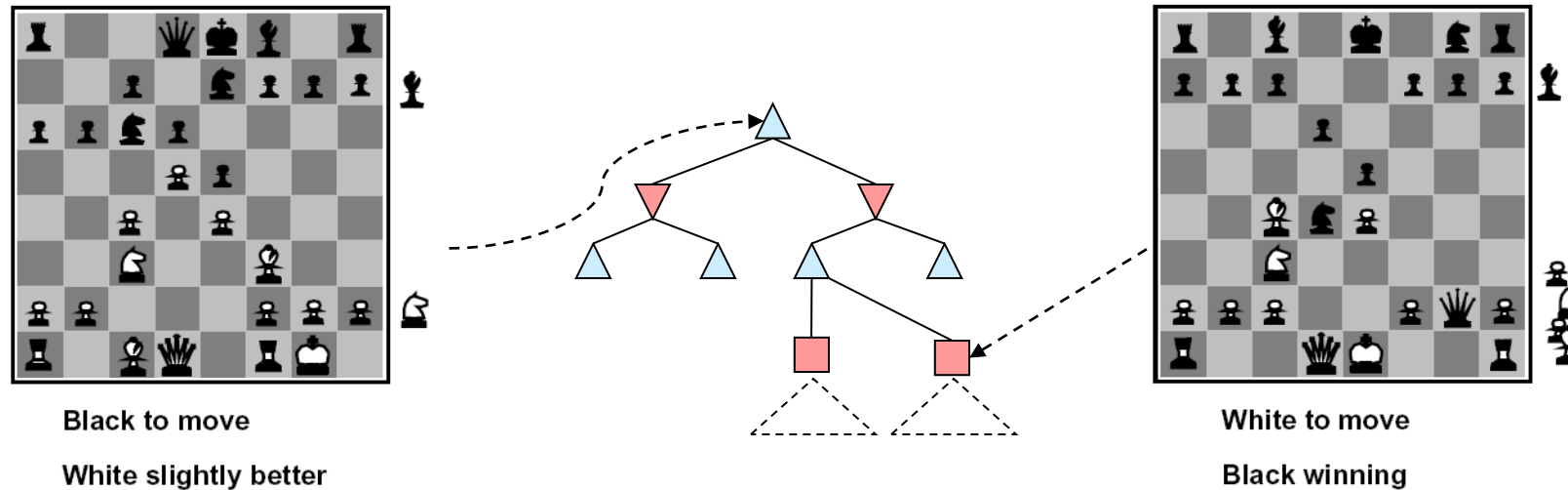


Evaluation Functions



Evaluation Functions

- Evaluation functions score non-terminals in depth-limited search

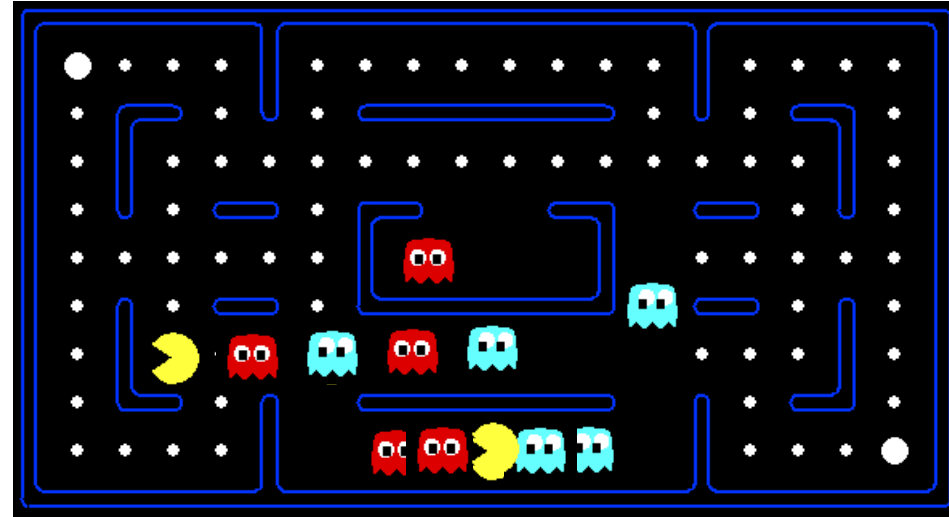


- Ideal function: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

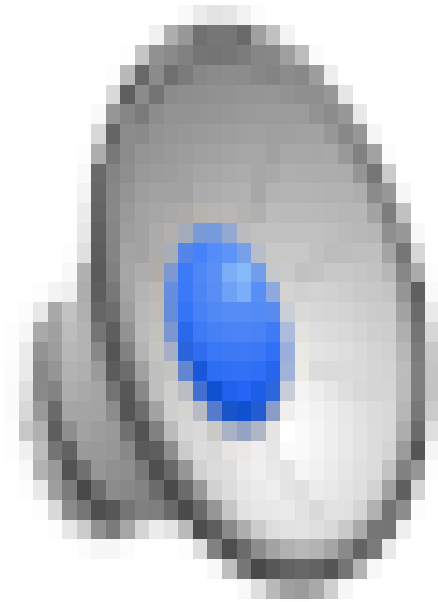
- e.g. $f_1(s) = (\text{num white queens} - \text{num black queens})$, etc.

Evaluation for Pacman

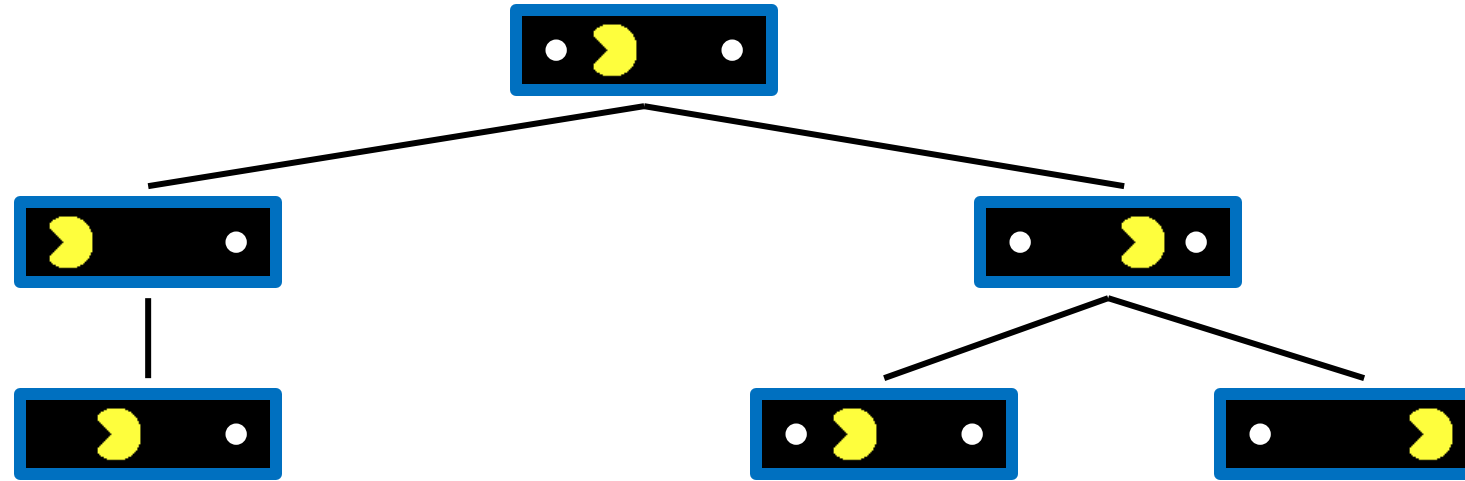


[Demo: thrashing $d=2$, thrashing $d=2$ (fixed evaluation function), smart ghosts coordinate (L6D6,7,8,10)]

Video of Demo Thrashing (d=2)

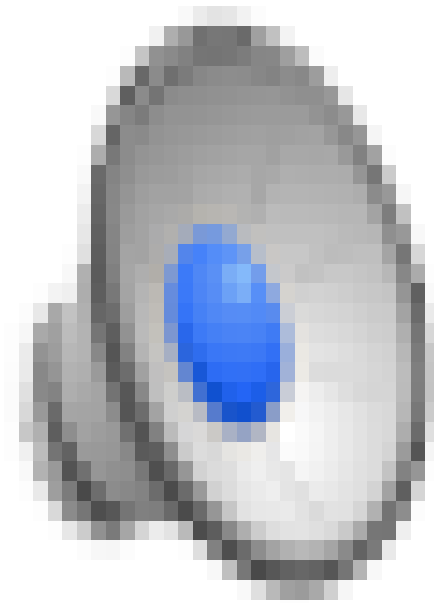


Why Pacman Starves

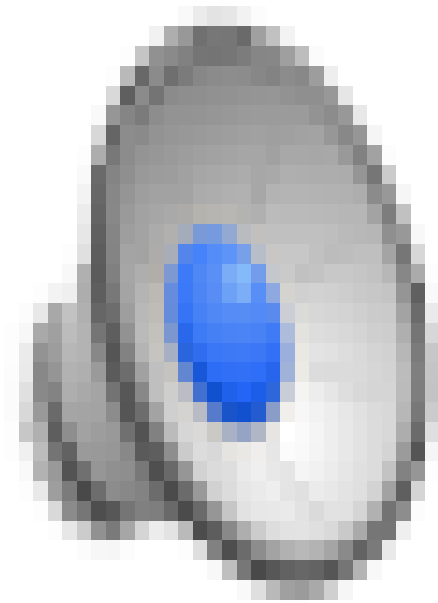


- A danger of replanning agents!
 - He knows his score will go up by eating the dot now (west, east)
 - He knows his score will go up just as much by eating the dot later (east, west)
 - There are no point-scoring opportunities after eating the dot (within the horizon, two here)
 - Therefore, waiting seems just as good as eating: he may go east, then back west in the next round of replanning!

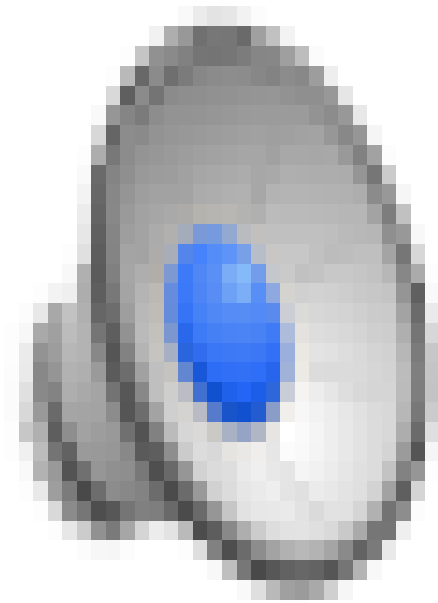
Video of Demo Thrashing -- Fixed (d=2)



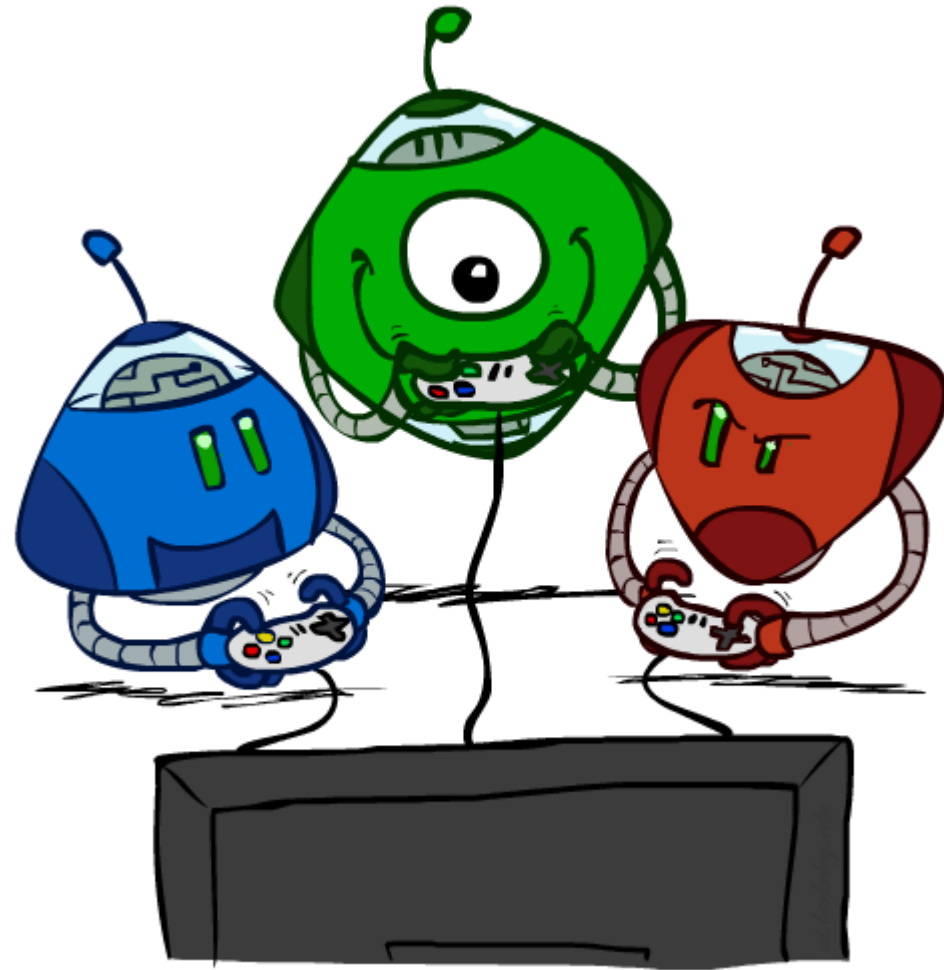
Video of Demo Smart Ghosts (Coordination)



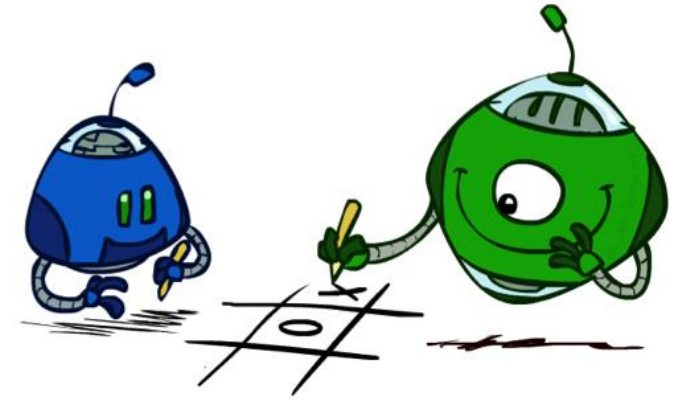
Video of Demo Smart Ghosts (Coordination) – Zoomed In



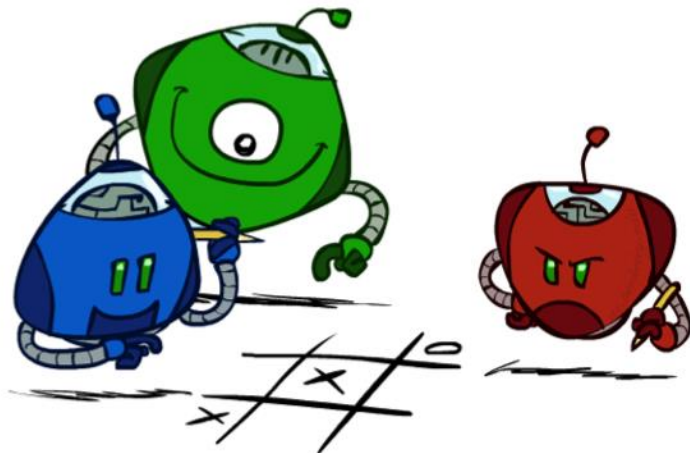
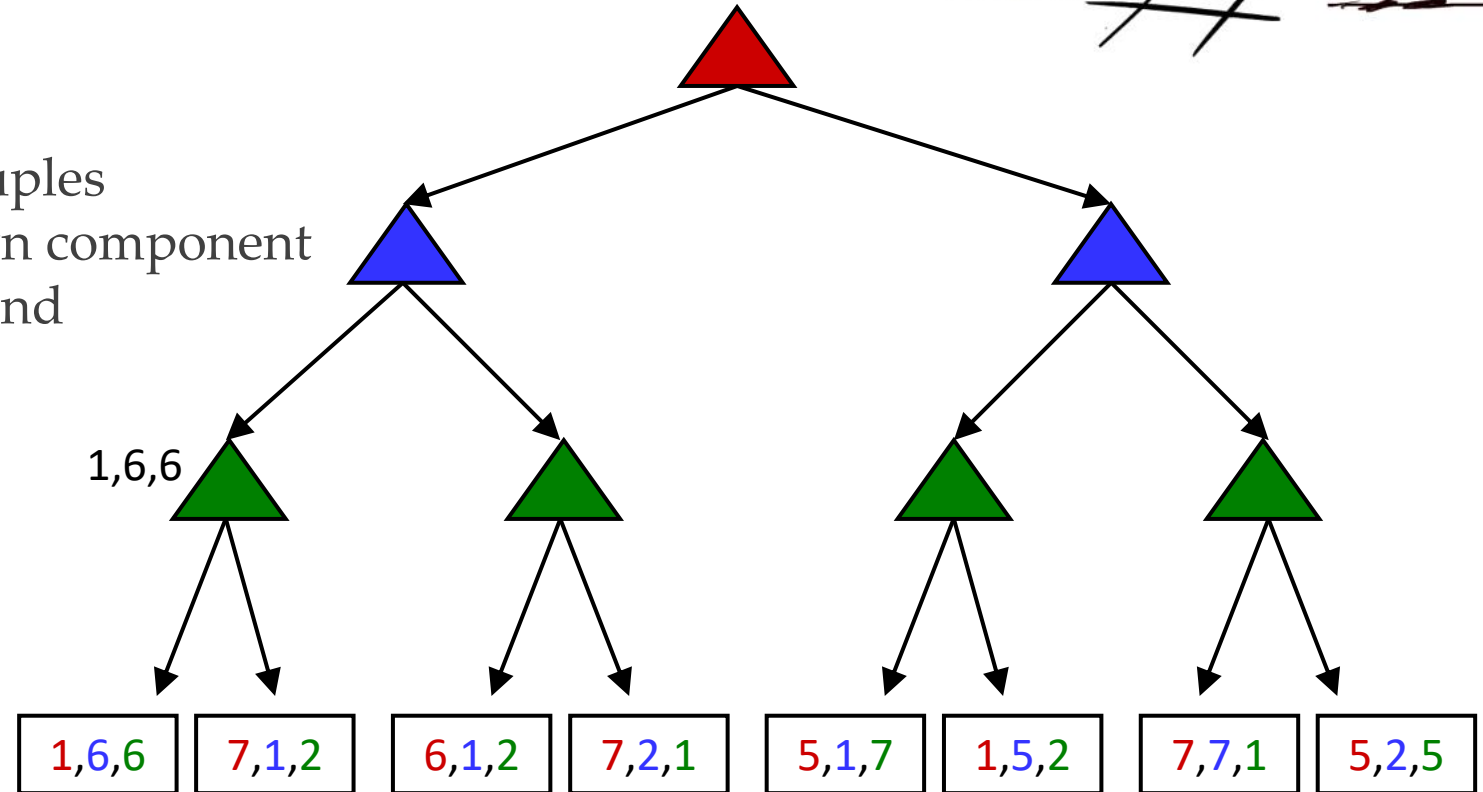
Other Game Types



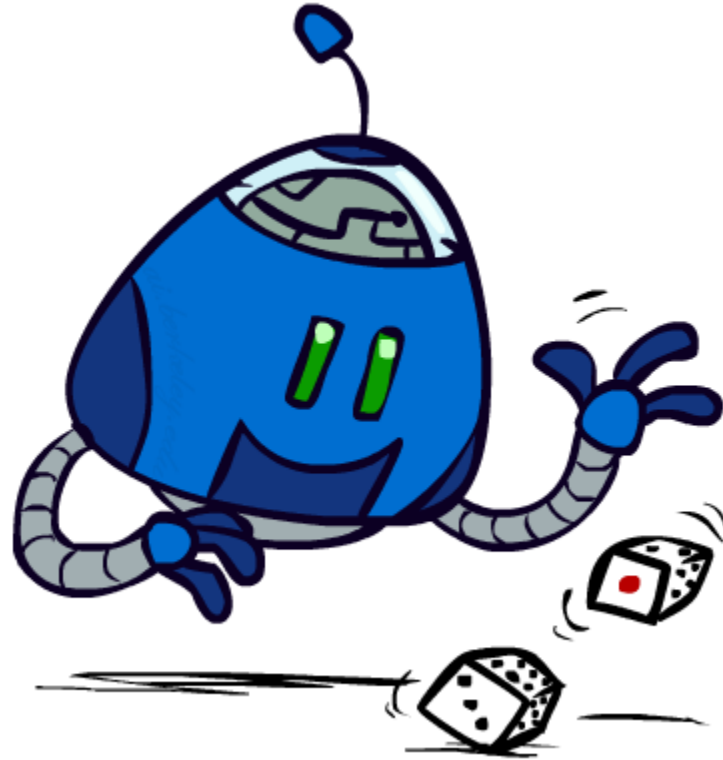
Multi-Agent Utilities



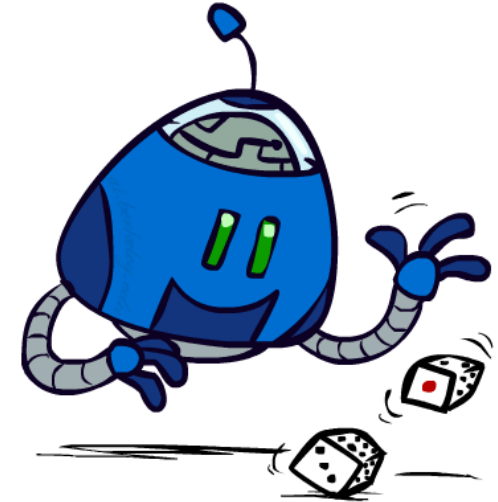
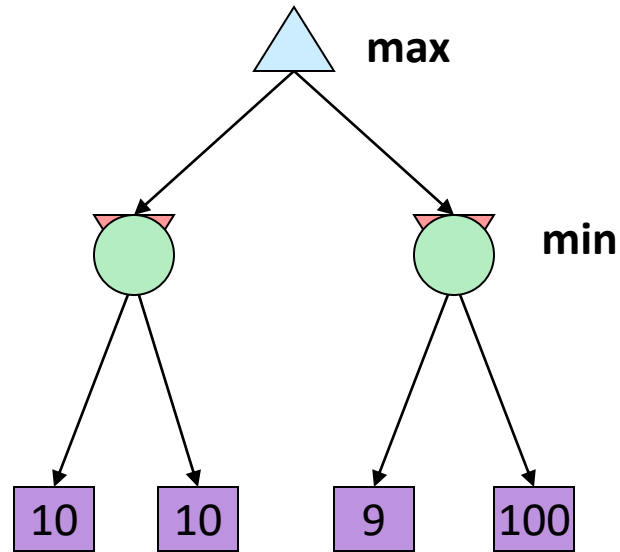
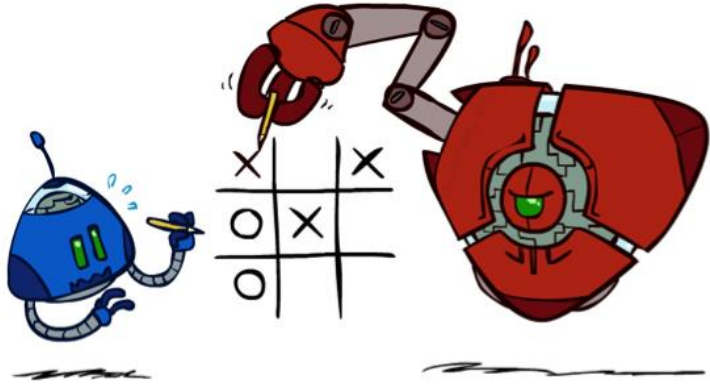
- What if the game is not zero-sum, or has multiple players?
- Generalization of minimax:
 - Terminals have utility tuples
 - Node values are also utility tuples
 - Each player maximizes its own component
 - Can give rise to cooperation and competition dynamically...



Uncertain Outcomes



Worst-Case vs. Average Case



Idea: Uncertain outcomes controlled by chance, not an adversary!

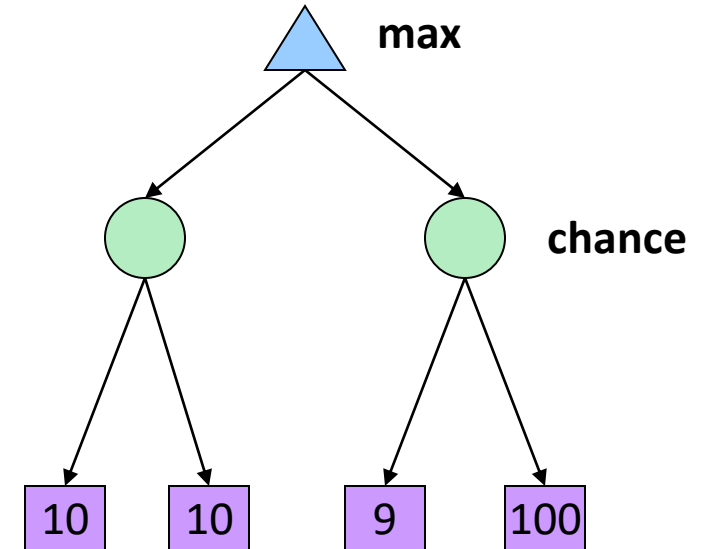
Why not minimax?

- Worst case reasoning is too conservative
- Need average case reasoning

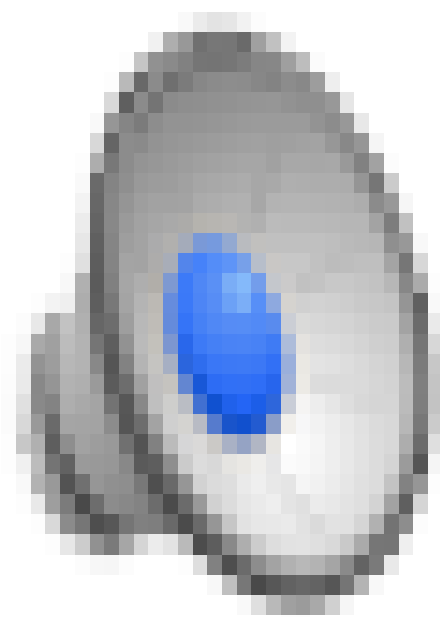


Expectimax Search

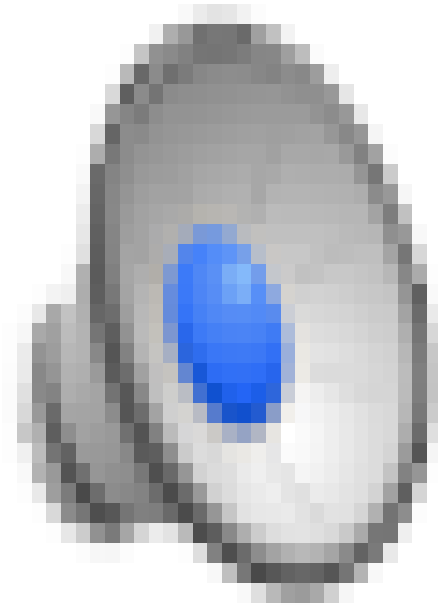
- Why wouldn't we know what the result of an action will be?
 - Explicit randomness: rolling dice
 - Unpredictable opponents: the ghosts respond randomly
 - Unpredictable humans: humans are not perfect
 - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (expectimax) outcomes, not worst-case (minimax) outcomes
- **Expectimax search**: compute the average score under optimal play
 - Max nodes as in minimax search
 - Chance nodes are like min nodes but the outcome is uncertain
 - Calculate their **expected utilities**
 - I.e. take weighted average (expectation) of children
- Later, we'll learn how to formalize the underlying uncertain-result problems as **Markov Decision Processes**



Video of Demo Minimax vs Expectimax (Min)



Video of Demo Minimax vs Expectimax (Exp)



Expectimax Pseudocode

```
def value(state):
```

```
    if the state is a terminal state: return the state's utility
```

```
    if the next agent is MAX: return max-value(state)
```

```
    if the next agent is EXP: return exp-value(state)
```

```
def max-value(state):
```

```
    initialize v =  $-\infty$ 
```

```
    for each successor of state:
```

```
        v = max(v, value(successor))
```

```
    return v
```

```
def exp-value(state):
```

```
    initialize v = 0
```

```
    for each successor of state:
```

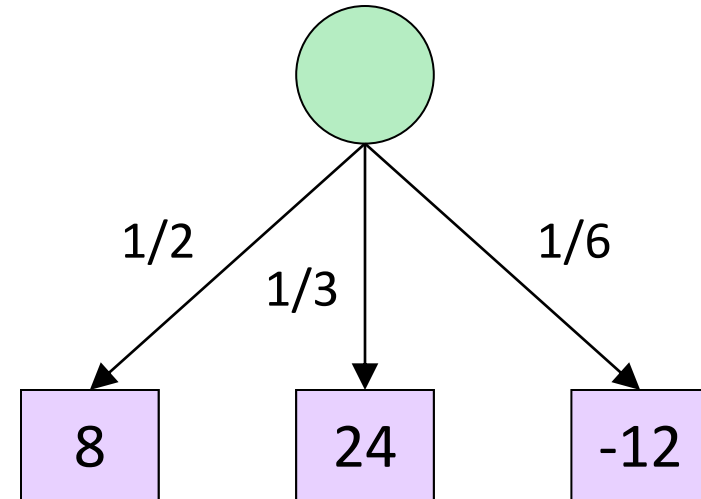
```
        p = probability(successor)
```

```
        v += p * value(successor)
```

```
    return v
```

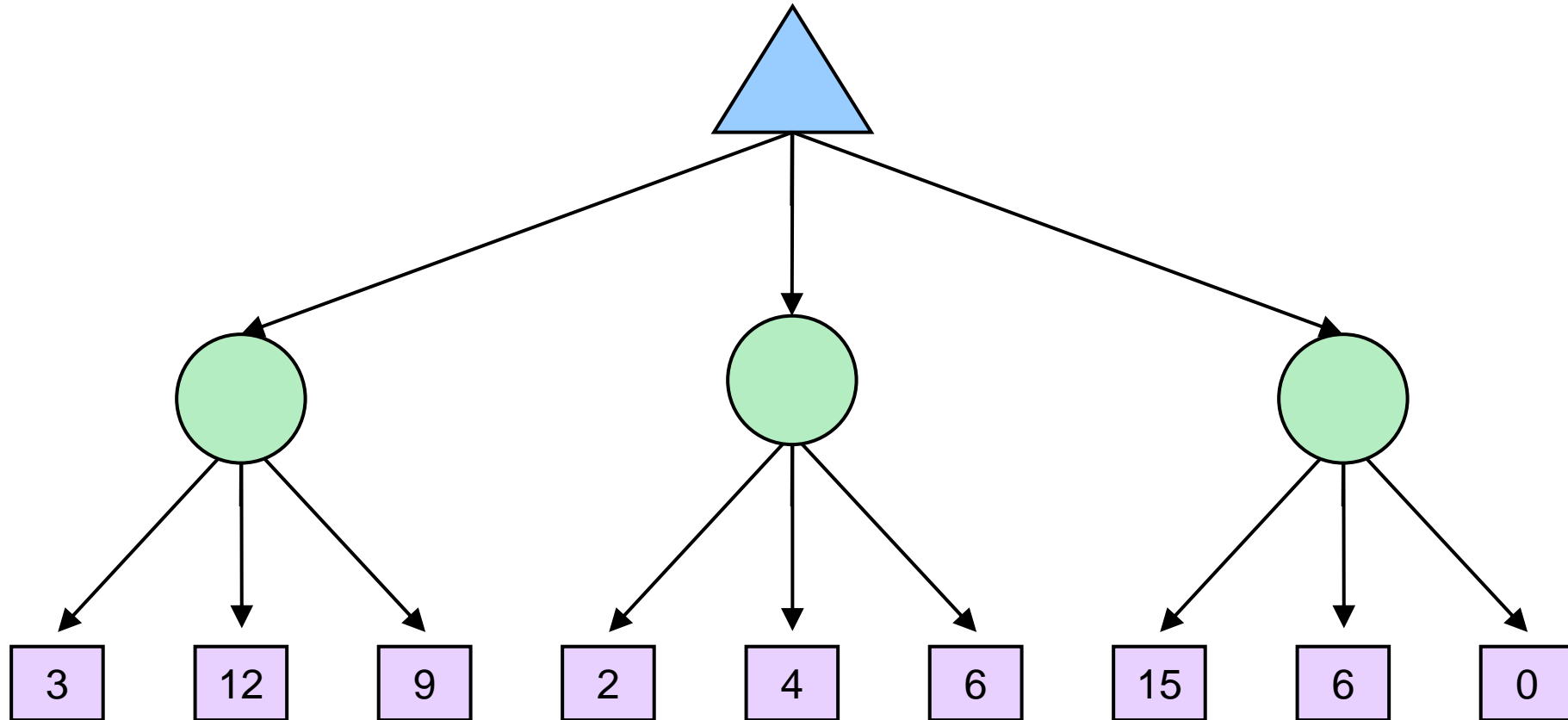
Expectimax Pseudocode

```
def exp-value(state):  
    initialize v = 0  
    for each successor of state:  
        p = probability(successor)  
        v += p * value(successor)  
    return v
```

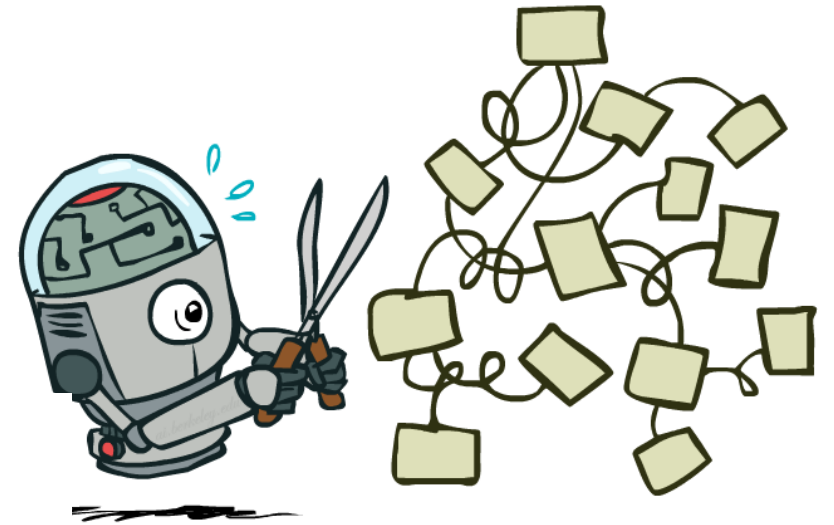
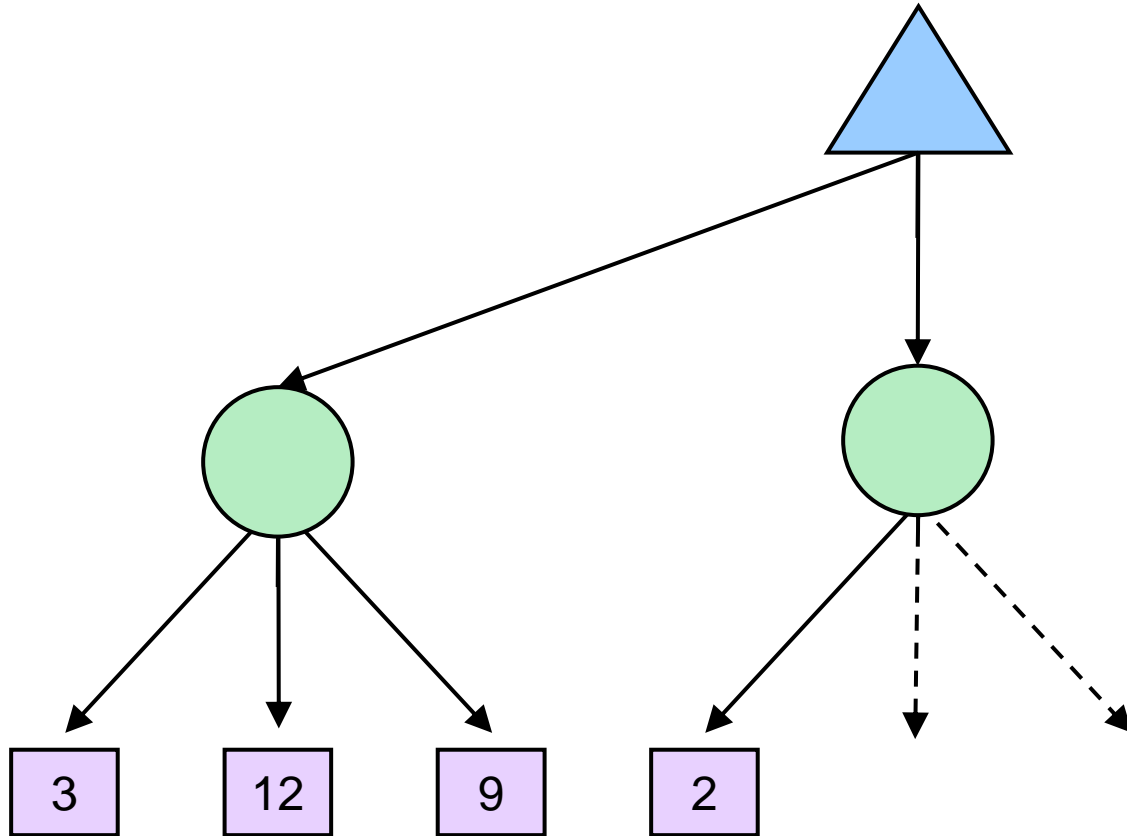


$$v = (1/2) (8) + (1/3) (24) + (1/6) (-12) = 10$$

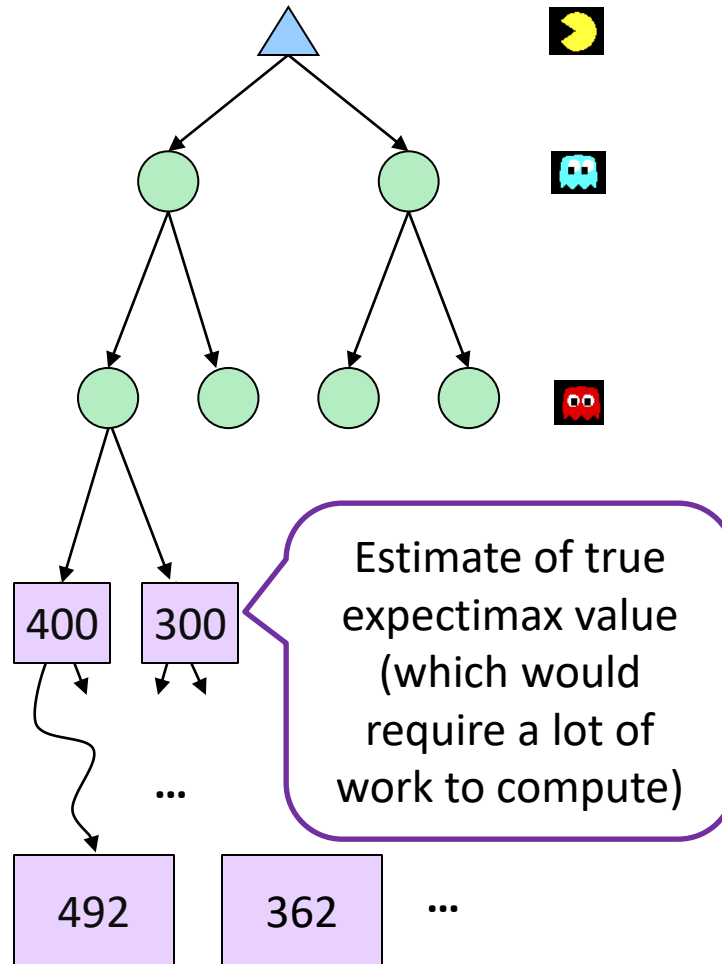
Expectimax Example



Expectimax Pruning?

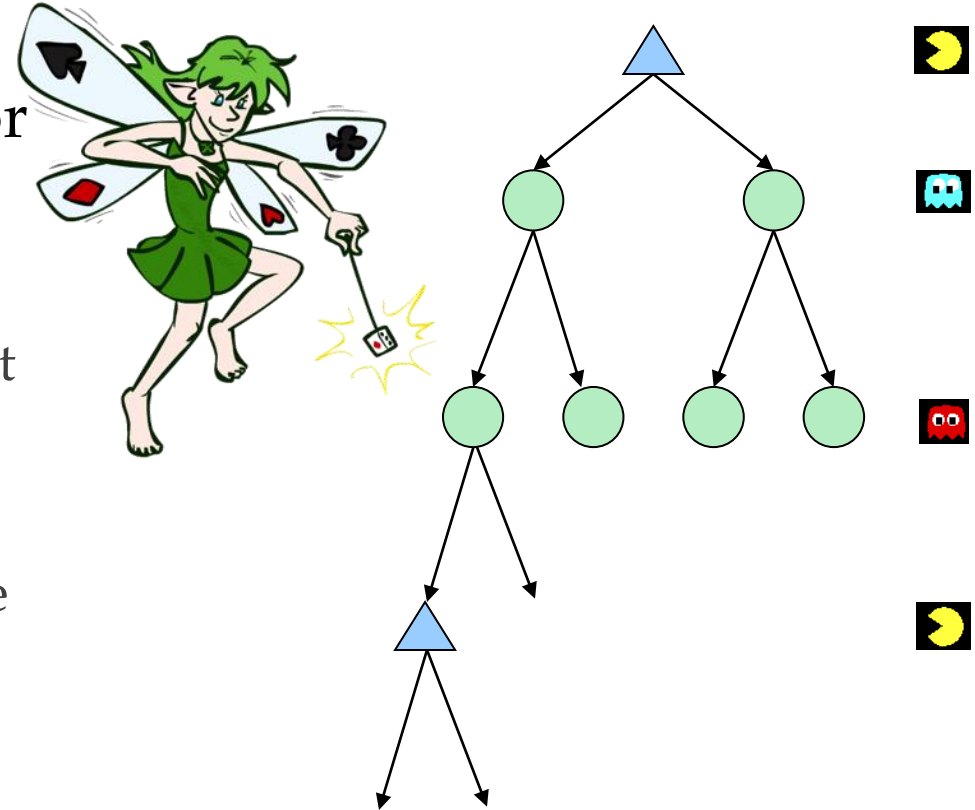


Depth-Limited Expectimax



What Probabilities to Use?

- In expectimax search, we have a probabilistic model of how the opponent (or environment) will behave in any state
 - Model could be a simple uniform distribution (roll a die)
 - Model could be sophisticated and require a great deal of computation
 - We have a chance node for any outcome out of our control: opponent or environment
 - The model might say that adversarial actions are likely!
- For now, assume each chance node magically comes along with probabilities that specify the distribution over its outcomes



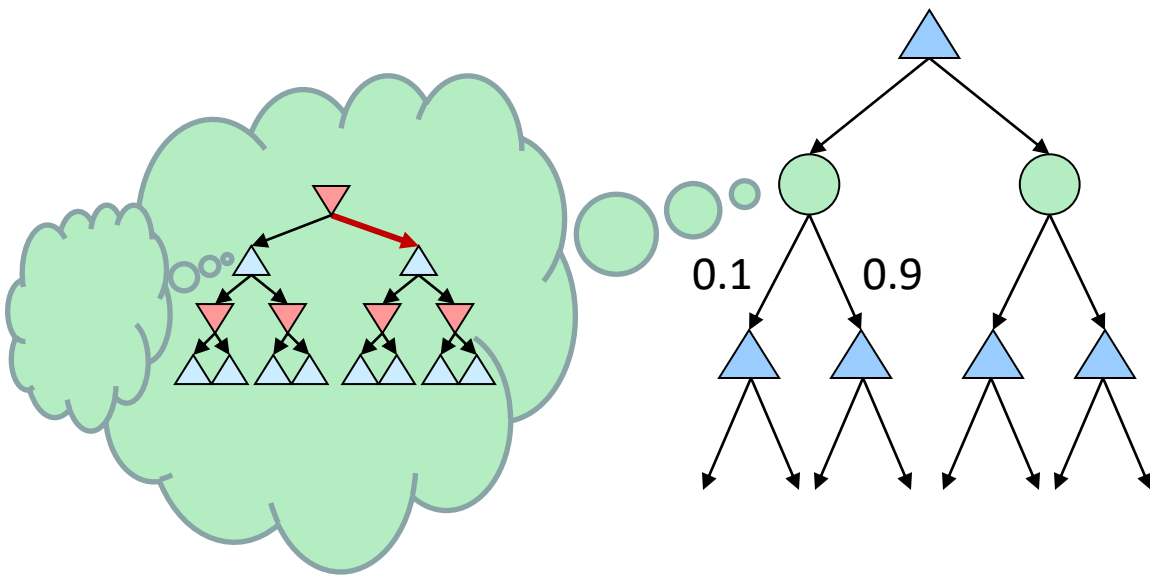
Having a probabilistic belief about another agent's action does not mean that the agent is flipping any coins!

Quiz: Informed Probabilities

- Let's say you know that your opponent is actually running a depth 2 minimax, using the result 80% of the time, and moving randomly otherwise
- Question: What tree search should you use?

- **Answer: Expectimax!**

- To figure out EACH chance node's probabilities, you have to run a simulation of your opponent
- This kind of thing gets very slow very quickly
- Even worse if you have to simulate your opponent simulating you...
- ... except for minimax and maximax, which have the nice property that it all collapses into one game tree



This is basically how you would model a human, except for their utility: their utility might be the same as yours (i.e. you try to help them, but they are depth 2 and noisy), or they might have a slightly different utility (like another person navigating in the office)

Modeling Assumptions



The Dangers of Optimism and Pessimism

Dangerous Optimism

Assuming chance when the world is adversarial

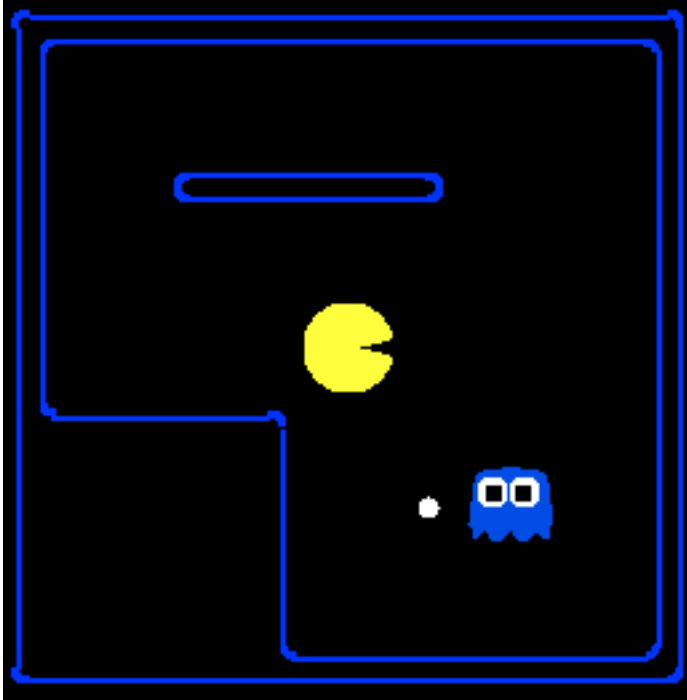


Dangerous Pessimism

Assuming the worst case when it's not likely



Assumptions vs. Reality



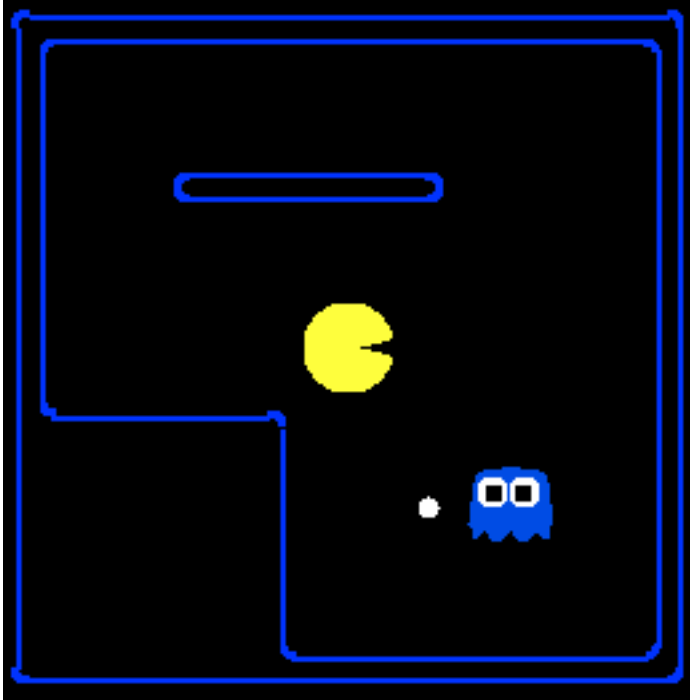
	Adversarial Ghost	Random Ghost
Minimax Pacman		
Expectimax Pacman		

Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble
Ghost used depth 2 search with an eval function that seeks Pacman

[Demos: world assumptions (L7D3,4,5,6)]

Assumptions vs. Reality



	Adversarial Ghost	Random Ghost
Minimax Pacman	Won 5/5 Avg. Score: 483	Won 5/5 Avg. Score: 493
Expectimax Pacman	Won 1/5 Avg. Score: -303	Won 5/5 Avg. Score: 503

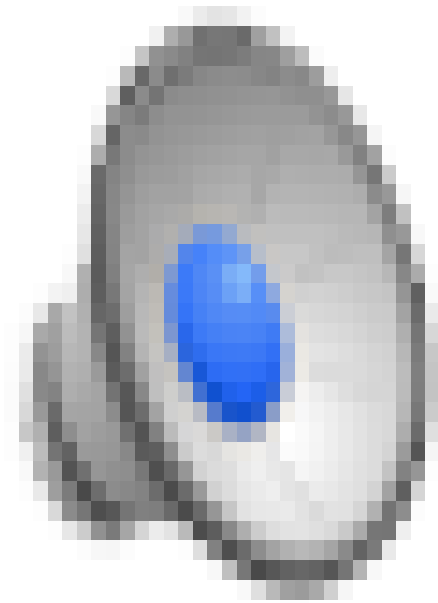
Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble
Ghost used depth 2 search with an eval function that seeks Pacman

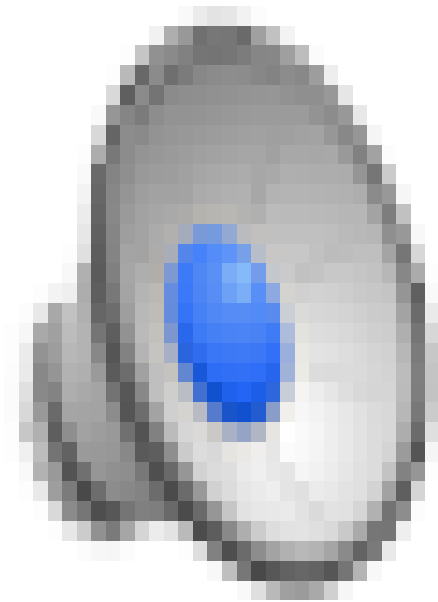
[Demos: world assumptions (L7D3,4,5,6)]

Video of Demo World Assumptions

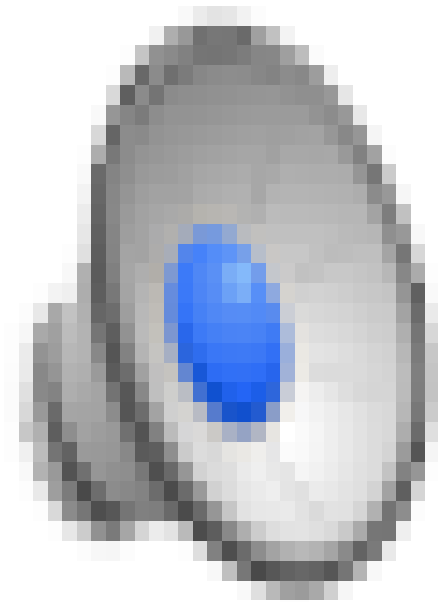
Random Ghost – Expectimax Pacman



Video of Demo World Assumptions Adversarial Ghost – Minimax Pacman

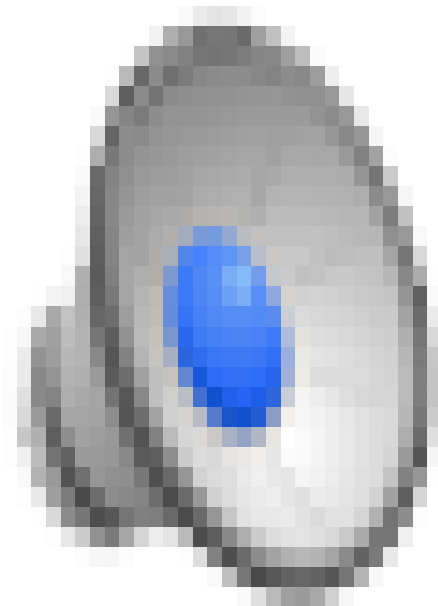


Video of Demo World Assumptions Adversarial Ghost – Expectimax Pacman



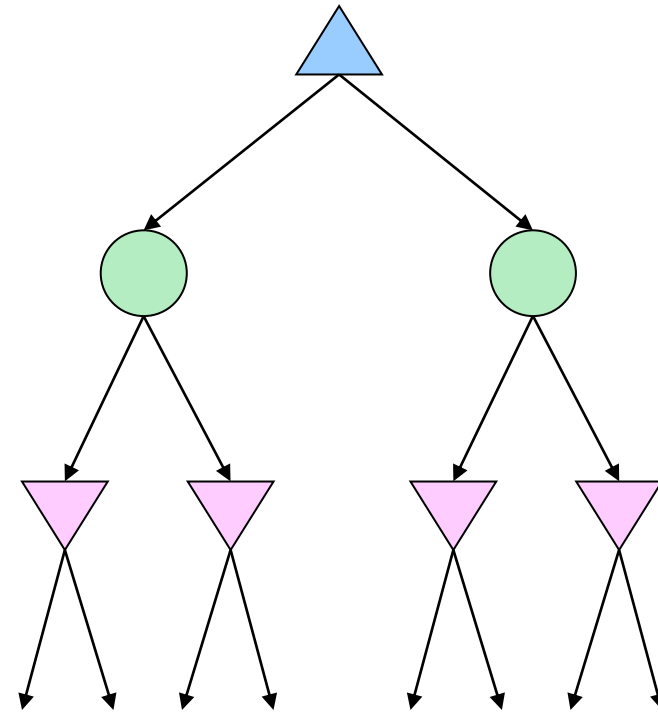
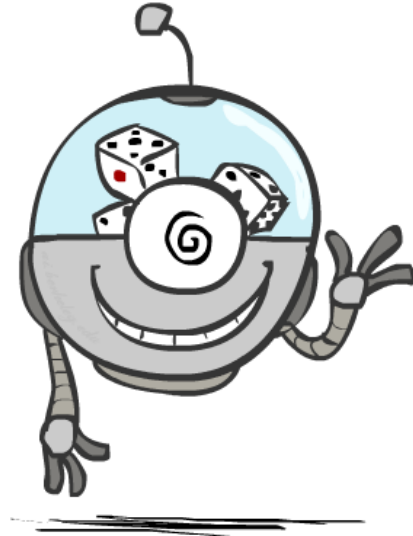
Video of Demo World Assumptions

Random Ghost – Minimax Pacman



Mixed Layer Types

- E.g. Backgammon
- Expectiminimax
 - Environment is an extra “random agent” player that moves after each min/max agent
 - Each node computes the appropriate combination of its children

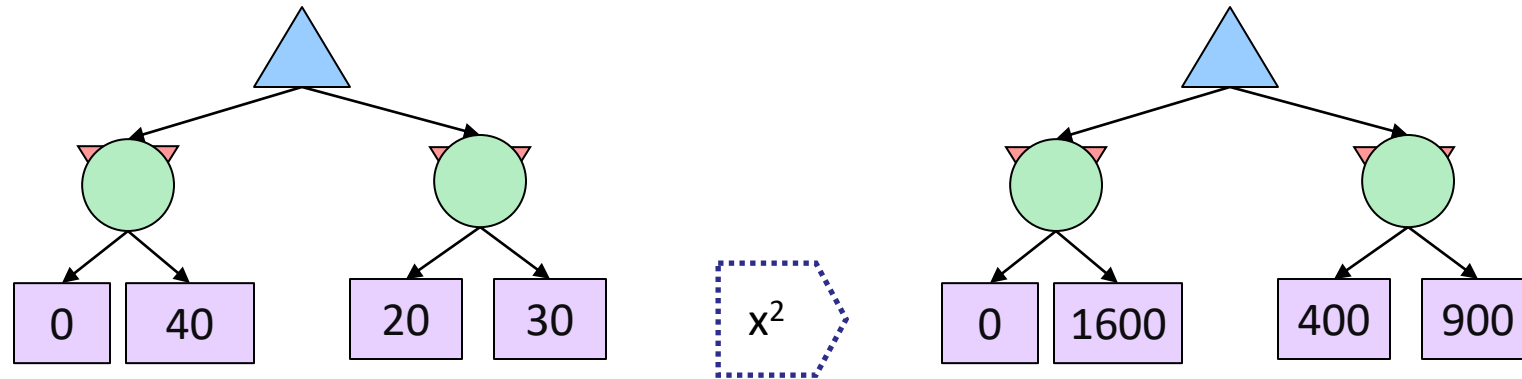


Example: Backgammon

- Dice rolls increase b : 21 possible rolls with 2 dice
 - Backgammon ≈ 20 legal moves
 - Depth 2 = $20 \times (21 \times 20)^3 = 1.2 \times 10^9$
- As depth increases, probability of reaching a given search node shrinks
 - So usefulness of search is diminished
 - So limiting depth is less damaging
 - But pruning is trickier...
- Historic AI: TDGammon uses depth-2 search + very good evaluation function + reinforcement learning:
world-champion level play
- 1st AI world champion in any game!



What Utility Values to Use?



$$x > y \Rightarrow f(x) > f(y)$$

$$f(x) = Ax + B \text{ where } A > 0$$

- For worst-case minimax reasoning, evaluation function scale doesn't matter
 - We just want better states to have higher evaluations (get the ordering right)
 - Minimax decisions are *invariant with respect to monotonic transformations on values*
- Expectiminimax decisions are *invariant with respect to positive affine transformations*
- Expectiminimax evaluation functions have to be aligned with actual win probabilities!



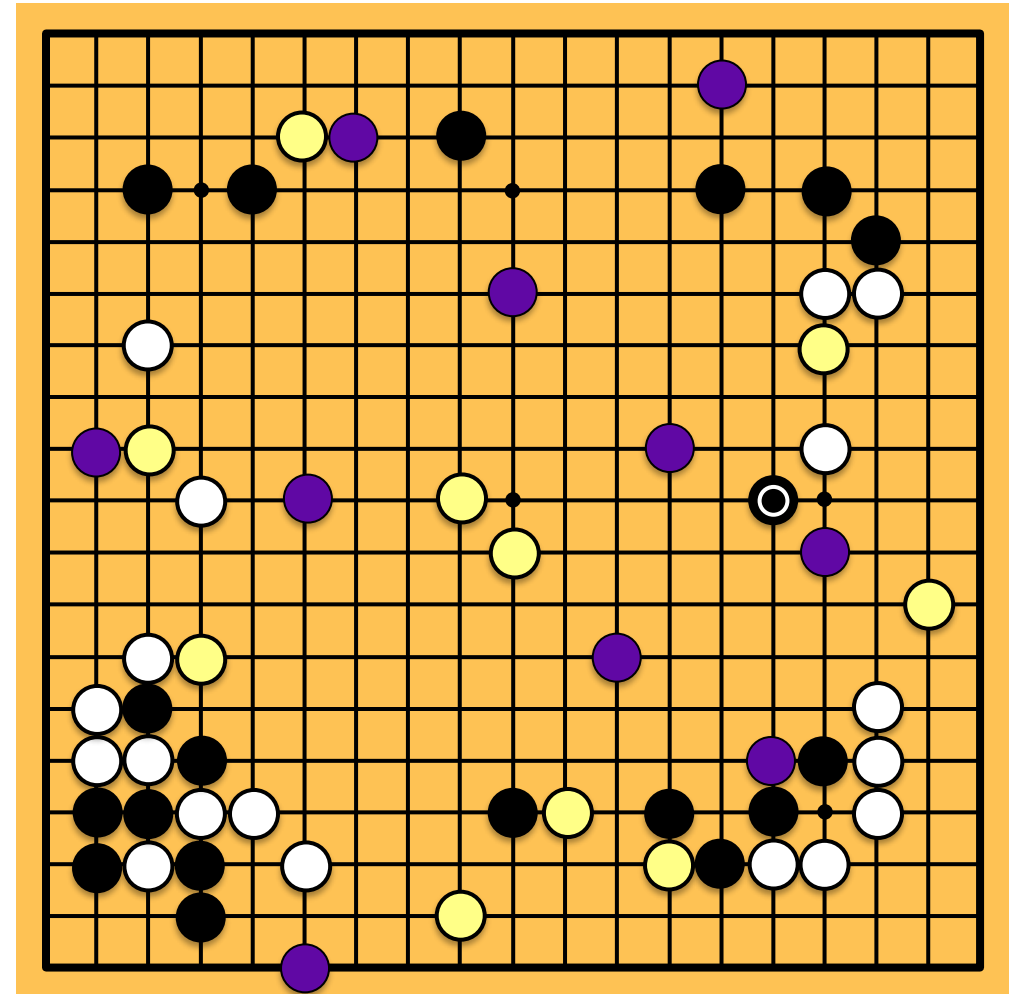
Monte Carlo Tree Search

- Methods based on alpha-beta search assume a fixed horizon
 - Pretty hopeless for Go, with $b > 300$
- MCTS combines two important ideas:
 - *Evaluation by rollouts* – play multiple games to termination from a state s (using a simple, fast rollout policy) and count wins and losses
 - *Selective search* – explore parts of the tree that will help improve the decision at the root, regardless of depth

Rollouts

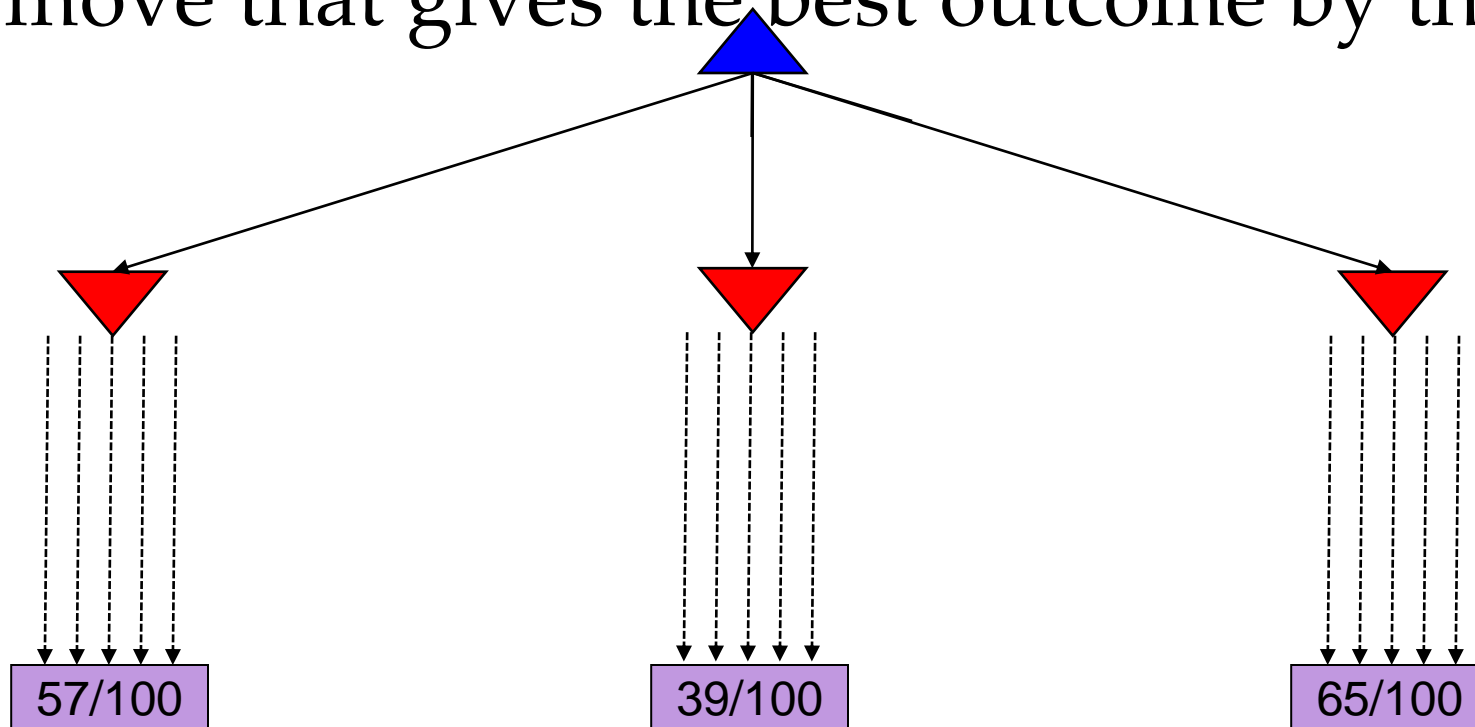
- For each rollout:
 - Repeat until terminal:
 - Play a move according to a fixed, fast rollout policy
 - Record the result
- Fraction of wins correlates with the true value of the position!
- Having a “better” rollout policy helps

“Move 37”



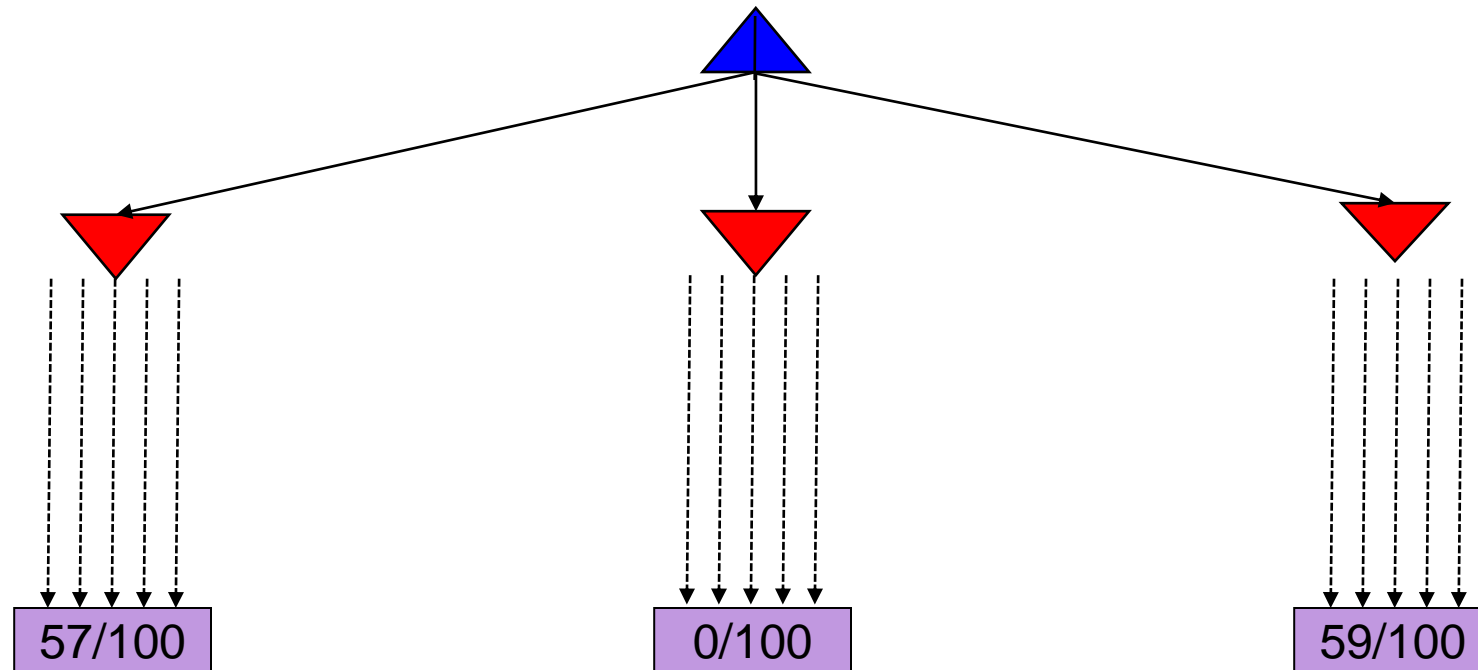
MCTS Version 0

- Do N rollouts from each child of the root, record fraction of wins
- Pick the move that gives the best outcome by this metric



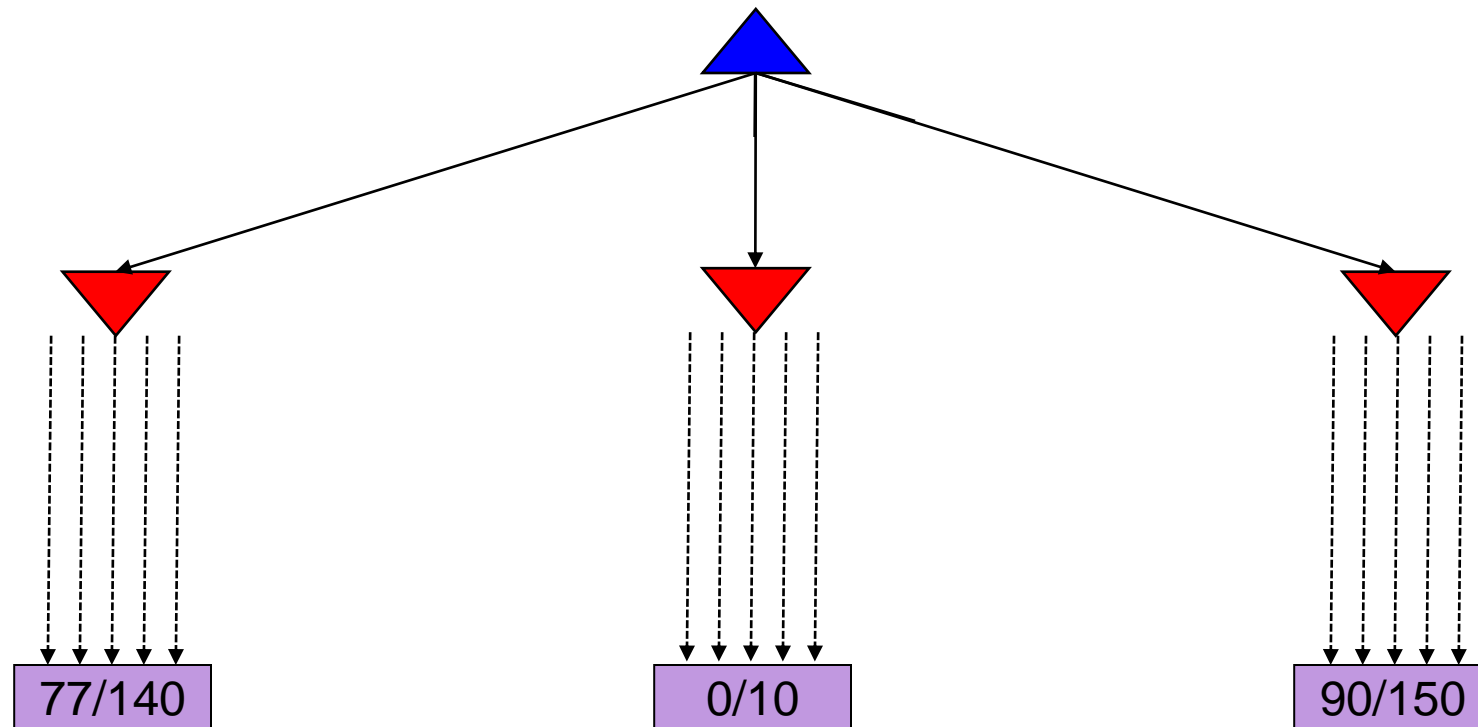
MCTS Simple Version

- Do N rollouts from each child of the root, record fraction of wins
- Pick the move that gives the best outcome by this metric



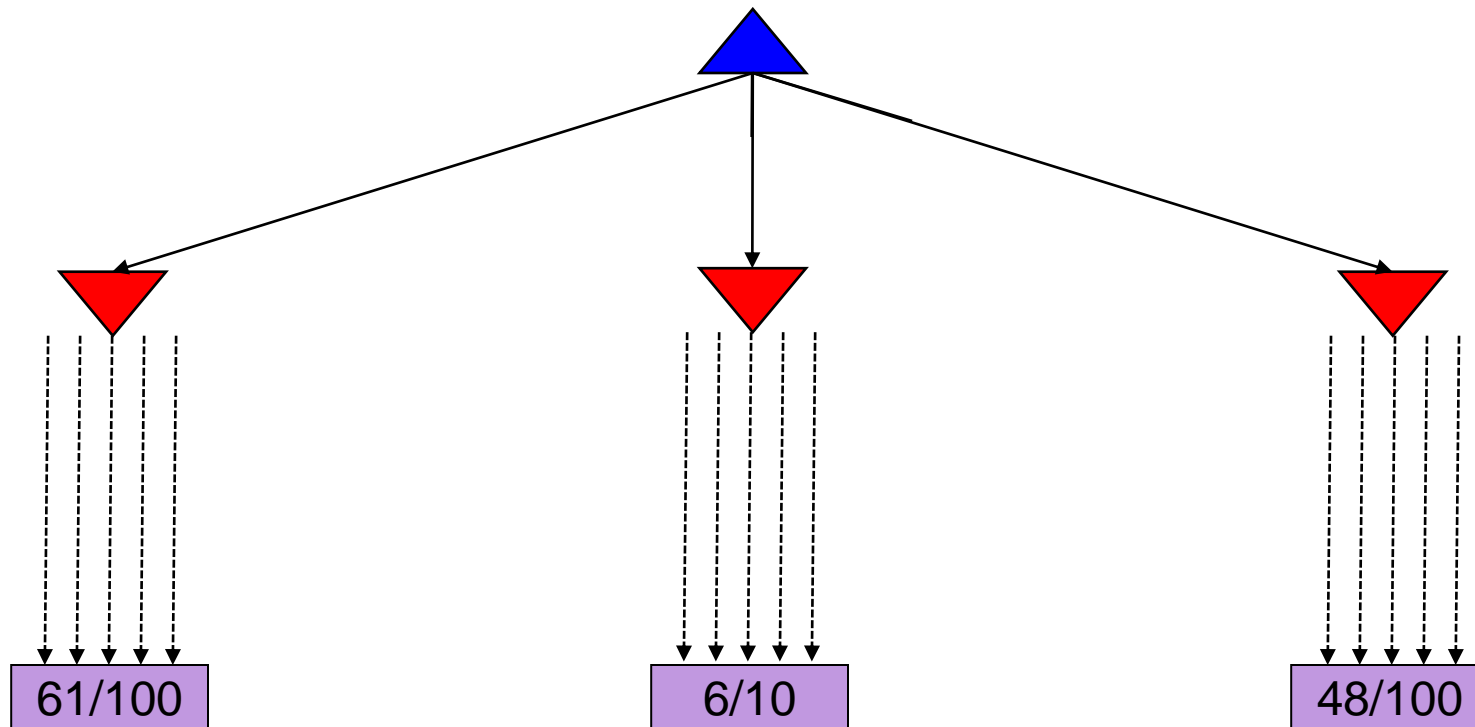
MCTS

- Allocate rollouts to more promising nodes



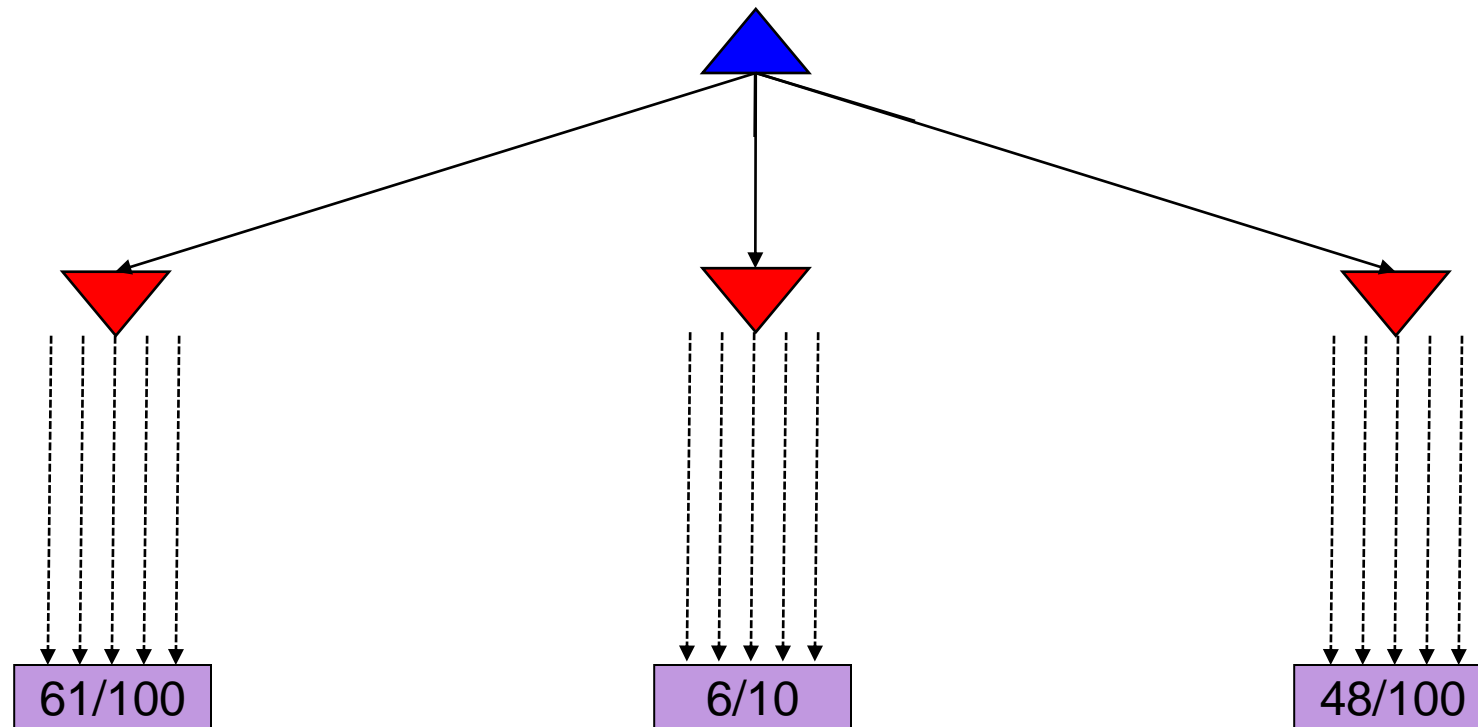
MCTS

- Allocate rollouts to more promising nodes



MCTS Version 1

- Allocate rollouts to more promising nodes
- Allocate rollouts to more uncertain nodes



UCB heuristics

- UCB1 formula combines “promising” and “uncertain”:

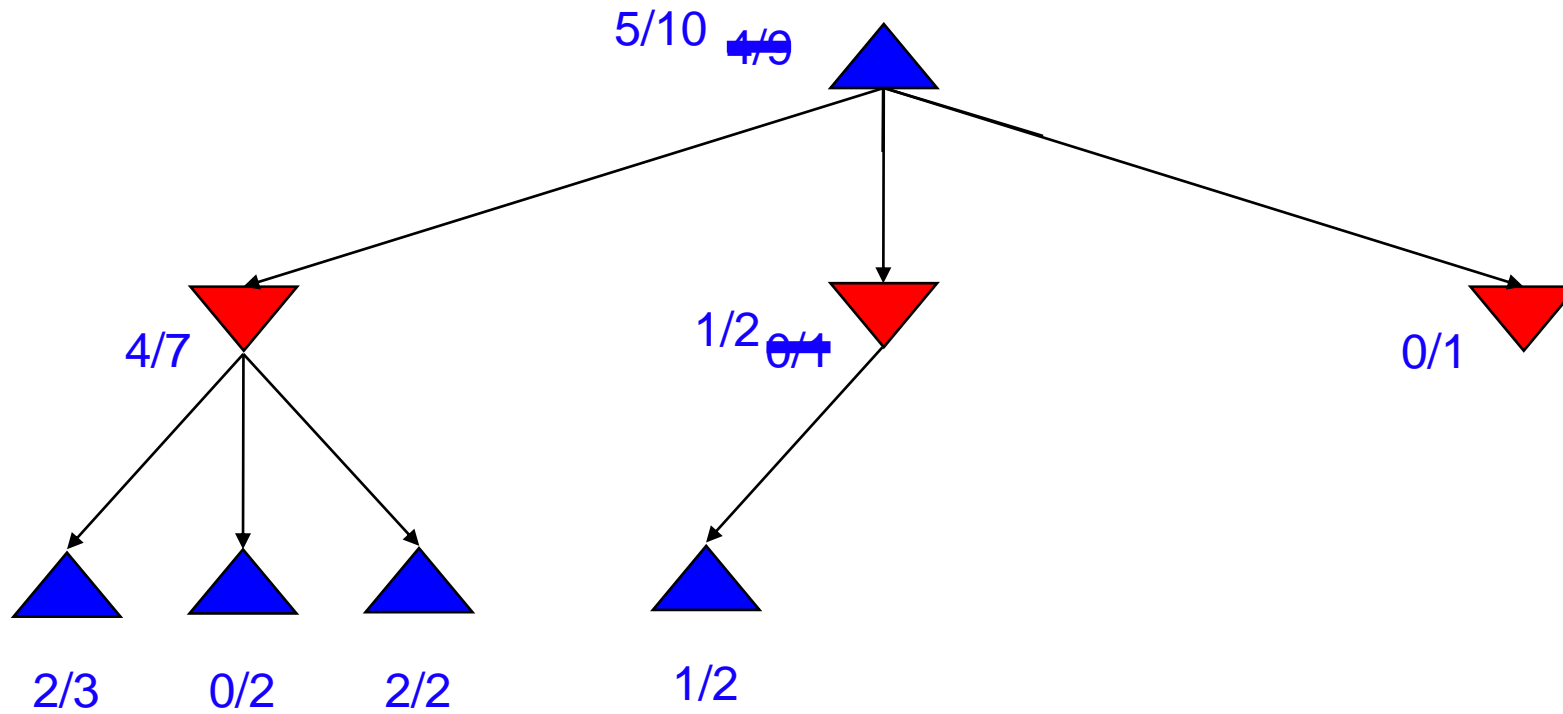
$$UCB1(n) = \frac{U(n)}{N(n)} + C \times \sqrt{\frac{\log N(\text{PARENT}(n))}{N(n)}}$$

- $N(n)$ = number of rollouts from node n
- $U(n)$ = total utility of rollouts (e.g., # wins) for $\text{Player}(\text{Parent}(n))$

MCTS Version 2: UCT

- Repeat until out of time:
 - Given the current search tree, recursively apply UCB to choose a path down to a leaf (not fully expanded) node n
 - Add a new child c to n and run a rollout from c
 - Update the win counts from c back up to the root
- Choose the action leading to the child with highest N

UCT Example



Why is there no min or max?????

- “Value” of a node, $U(n)/N(n)$, is a weighted *sum* of child values!
- Idea: as $N \rightarrow \infty$, the vast majority of rollouts are concentrated in the best children, so weighted average \rightarrow **max/min**
- Theorem: as $N \rightarrow \infty$ UCT selects the minimax move
 - (but N never approaches infinity!)

Summary

- Games require decisions when optimality is impossible
 - Bounded-depth search and approximate evaluation functions
- Games force efficient use of computation
 - Alpha-beta pruning, MCTS
- Game playing has produced important research ideas
 - Reinforcement learning (checkers)
 - Iterative deepening (chess)
 - Rational metareasoning (Othello)
 - Monte Carlo tree search (chess, Go)
 - Solution methods for partial-information games in economics (poker)
- Video games present much greater challenges – lots to do!
 - $b = 10^{500}$, $|S| = 10^{4000}$, $m = 10,000$, partially observable, often > 2 players