

## 61A Lecture 22

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Wednesday, October 22

## Announcements

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- Project 3 is due Thursday 10/23 @ 11:59pm
  - Please submit two ways: the normal way and using `python3 ok --submit!`
  - You can view your ok submission on the ok website: <http://ok.cs61a.org>
- Midterm 2 is on Monday 10/27 7pm–9pm
  - Review session on Saturday 10/25 3pm–6pm in 2050 VLSB
  - Conflict form submissions are due Wednesday 10/22!

# Sets

## Sets

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One more built-in Python container type

- Set literals are enclosed in braces
- Duplicate elements are removed on construction
- Sets are unordered, just like dictionary entries

```
>>> s = {3, 2, 1, 4, 4}
>>> s
{1, 2, 3, 4}
>>> 3 in s
True
>>> len(s)
4
>>> s.union({1, 5})
{1, 2, 3, 4, 5}
>>> s.intersection({6, 5, 4, 3})
{3, 4}
```

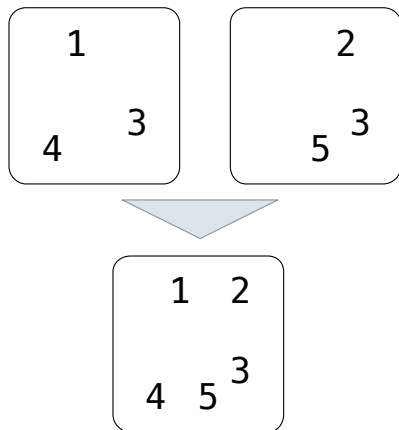
## Implementing Sets

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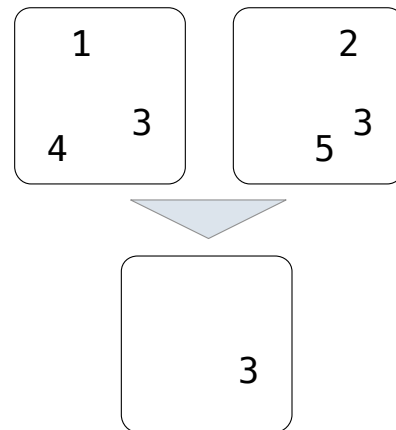
What we should be able to do with a set:

- **Membership testing:** Is a value an element of a set?
- **Union:** Return a set with all elements in set1 or set2
- **Intersection:** Return a set with any elements in set1 and set2
- **Adjoin:** Return a set with all elements in s and a value v

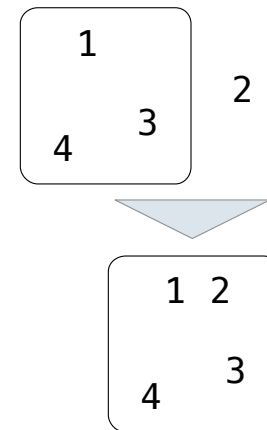
**Union**



**Intersection**



**Adjoin**



## Sets as Unordered Sequences

## Sets as Unordered Sequences

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**Proposal 1:** A set is represented by a linked list that contains no duplicate items.

```
def empty(s):
    return s is Link.empty

def set_contains(s, v):
    """Return whether set s contains value v.

    >>> s = Link(1, Link(2, Link(3)))
    >>> set_contains(s, 2)
    True
    """
    if empty(s):
        return False
    elif s.first == v:
        return True
    else:
        return set_contains(s.rest, v)
```

**Time order of growth**

$\Theta(1)$

*Time depends on whether  
& where  $v$  appears in  $s$*

$\Theta(n)$

*Assuming  $v$  either  
does not appear in  $s$   
**or**  
appears in a uniformly  
distributed random location*

(Demo)

## Sets as Unordered Sequences

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```
def adjoin_set(s, v):  
    if set_contains(s, v):  
        return s  
    else:  
        return Link(v, s)
```

```
def intersect_set(set1, set2):  
    in_set2 = lambda v: set_contains(set2, v)  
    return keep_if(set1, in_set2)
```

Need a new version defined  
for Link instances

```
def union_set(set1, set2):  
    not_in_set2 = lambda v: not set_contains(set2, v)  
    set1_not_set2 = keep_if(set1, not_in_set2)  
    return extend(set1_not_set2, set2)
```

Need a new version defined  
for Link instances

**Time order of growth**

$\Theta(n)$

The size of the set

$\Theta(n^2)$

If sets are  
the same size

$\Theta(n^2)$

(Demo)



## Sets as Ordered Sequences

## Sets as Ordered Sequences

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**Proposal 2:** A set is represented by a linked list with unique elements that is *ordered from least to greatest*

Parts of the program that...	Assume that sets are...	Using...
Use sets to contain values	Unordered collections	<code>empty, set_contains, adjoin_set, intersect_set, union_set</code>
Implement set operations	Ordered linked lists	<code>first, rest, &lt;, &gt;, ==</code>

*Different parts of a program may make different assumptions about data*

## Sets as Ordered Sequences

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**Proposal 2:** A set is represented by a linked list with unique elements that is *ordered from least to greatest*

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Link.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Link(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect_set(set1, set2.rest)
```

Order of growth?  $\Theta(n)$

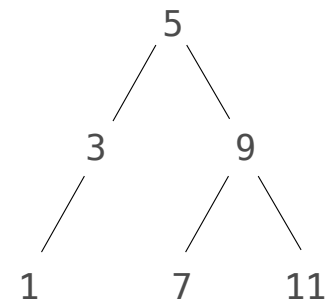
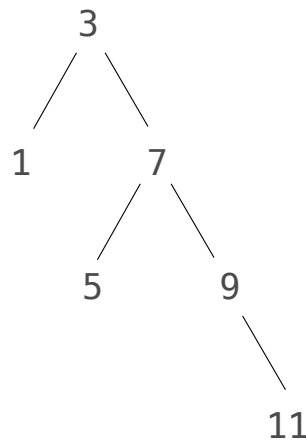
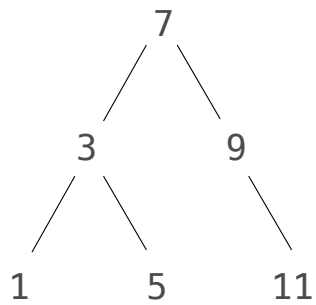
## Sets as Binary Search Trees

## Binary Search Trees

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**Proposal 3:** A set is represented as a Tree with two branches. Each entry is:

- Larger than all entries in its left branch and
- Smaller than all entries in its right branch

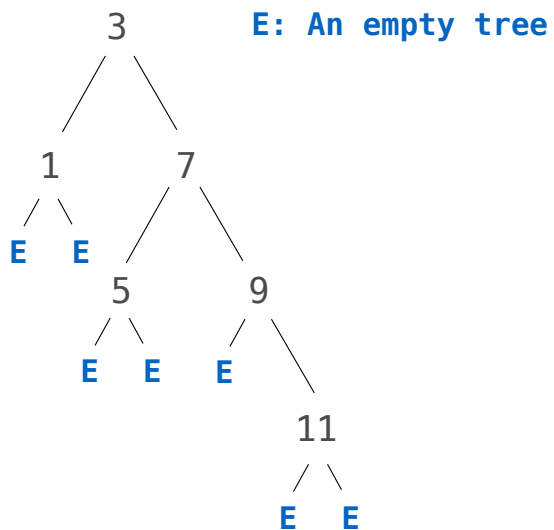


## Binary Tree Class

A binary tree is a tree that has a left branch and a right branch

**Idea:** Fill the place of a missing left branch with an empty tree

**Idea 2:** An instance of BinaryTree always has exactly two branches



```
class BinaryTree(Tree):
    empty = Tree(None)
    empty.is_empty = True

    def __init__(self, entry, left=empty, right=empty):
        Tree.__init__(self, entry, (left, right))
        self.is_empty = False
```

```
@property
def left(self):
    return self.branches[0]
```

```
@property
def right(self):
    return self.branches[1]
```

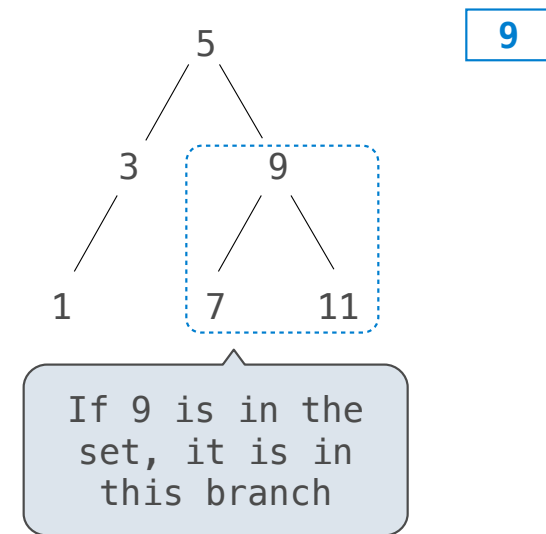
```
Bin = BinaryTree
t = Bin(3, Bin(1),
        Bin(7, Bin(5),
            Bin(9, Bin.empty,
                Bin(11))))
```

## Membership in Binary Search Trees

`set_contains` traverses the tree

- If the element is not the entry, it can only be in either the left or right branch
- By focusing on one branch, we reduce the set by about half with each recursive call

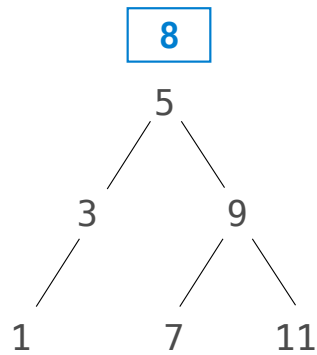
```
def set_contains(s, v):  
    if s.is_empty:  
        return False  
    elif s.entry == v:  
        return True  
    elif s.entry < v:  
        return set_contains(s.right, v)  
    elif s.entry > v:  
        return set_contains(s.left, v)
```



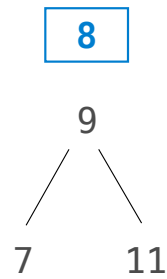
Order of growth?  $\Theta(h)$  on average

$\Theta(\log n)$  on average for a balanced tree

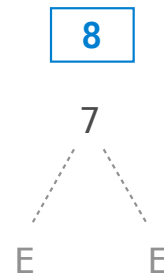
# Adjoining to a Tree Set



*Right!*



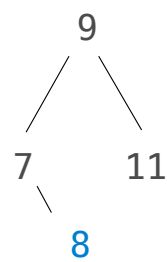
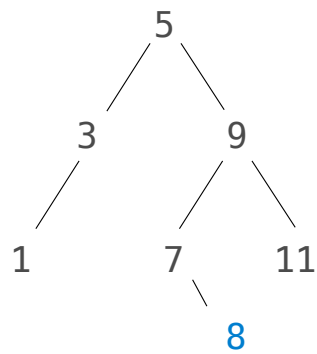
*Left!*



*Right!*



*Stop!*



(Demo)



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