

61A Lecture 26

Announcements

Programming Languages

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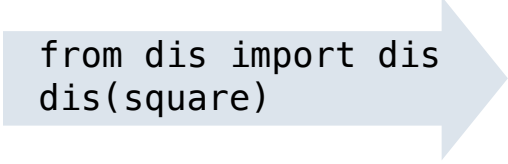
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To create a new programming language, you either need a:

- **Specification:** A document describe the precise syntax and semantics of the language
- **Canonical Implementation:** An interpreter or compiler for the language

Parsing

Reading Scheme Lists

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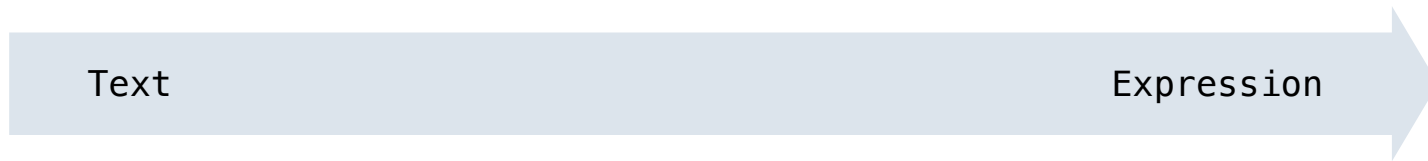
Parsing

Parsing

A Parser takes text and returns an expression

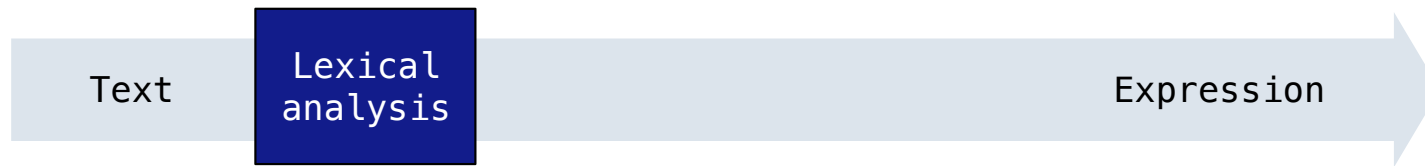
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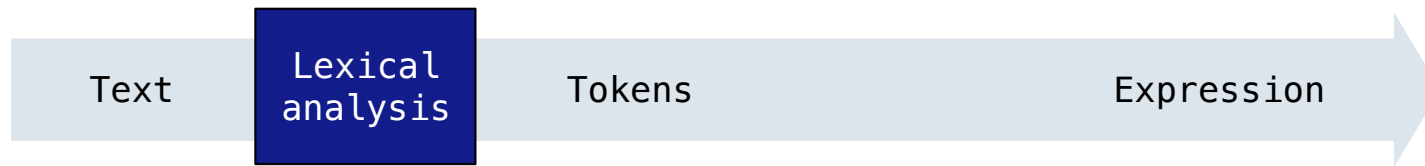
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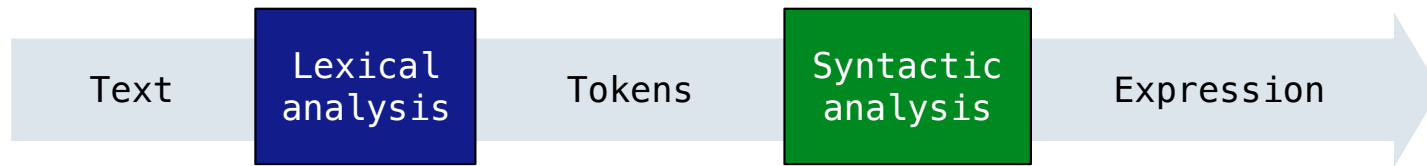
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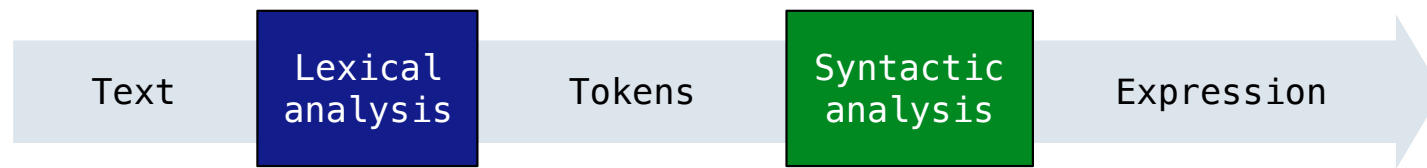
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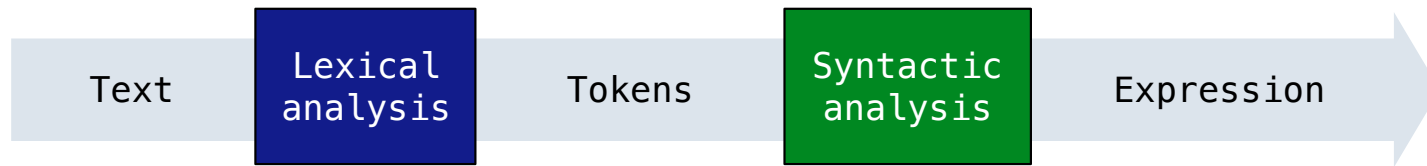


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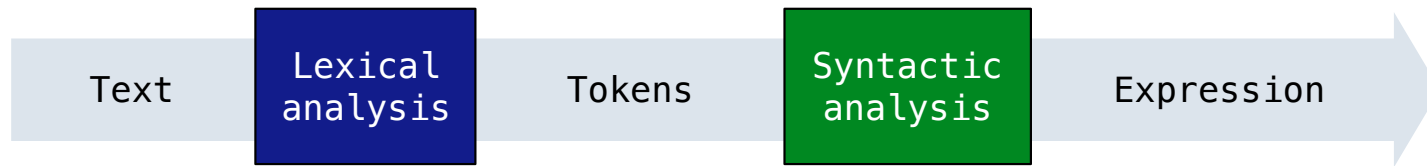
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↳

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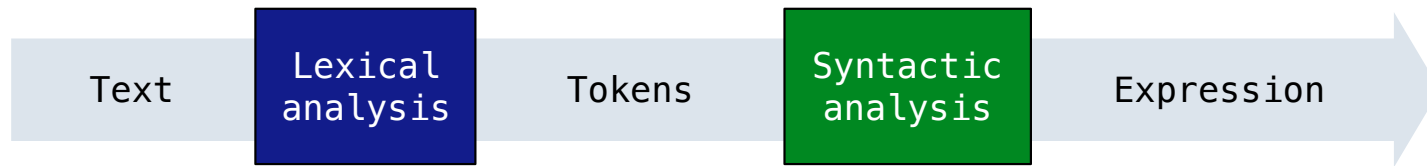
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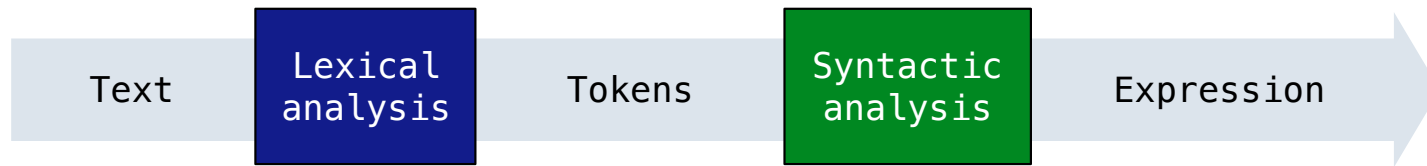
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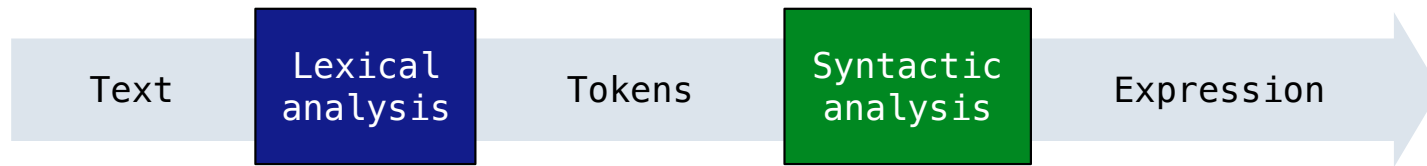
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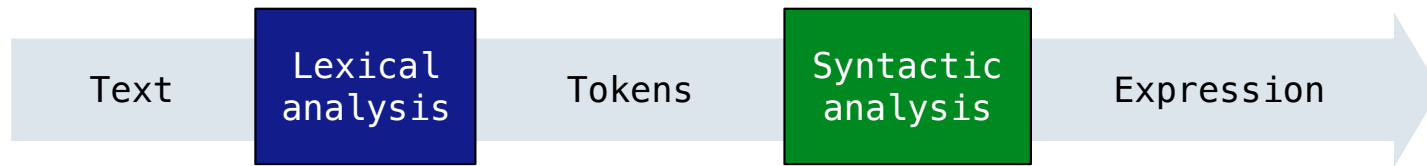
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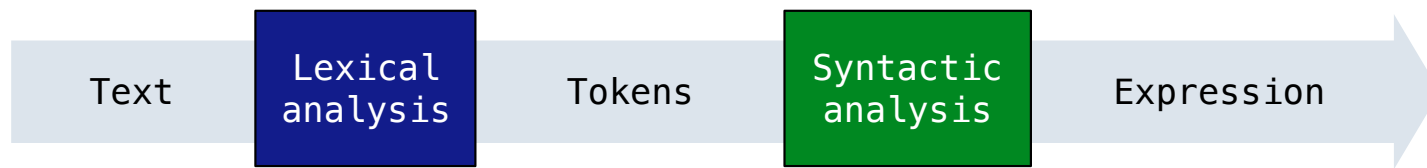
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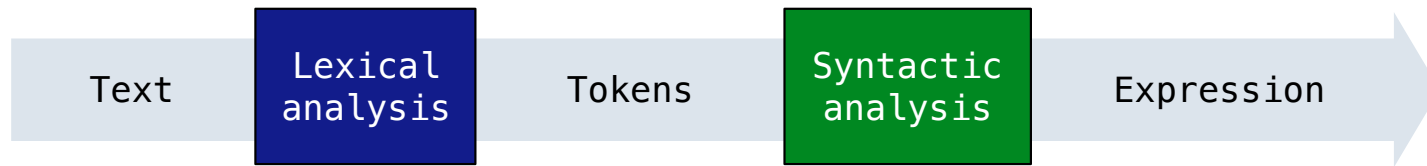
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- Iterative process

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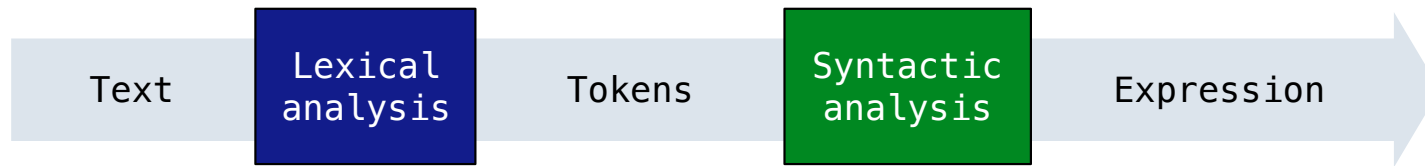
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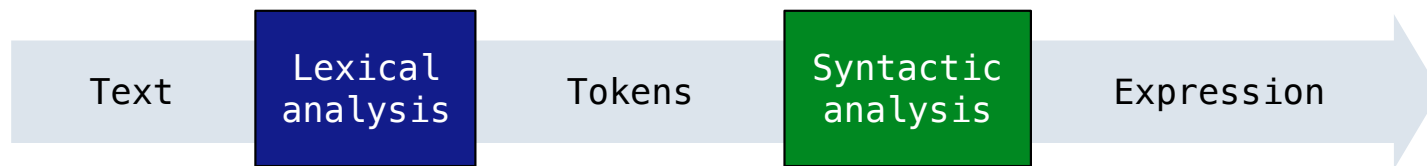
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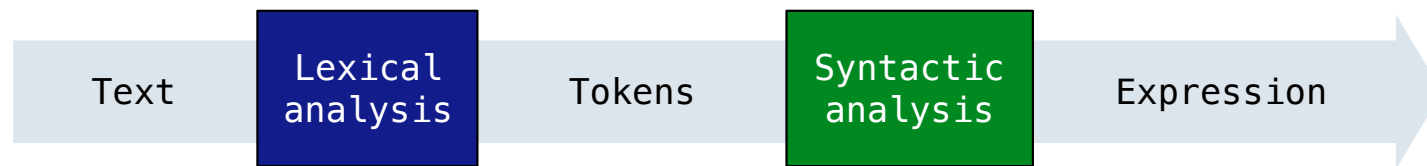
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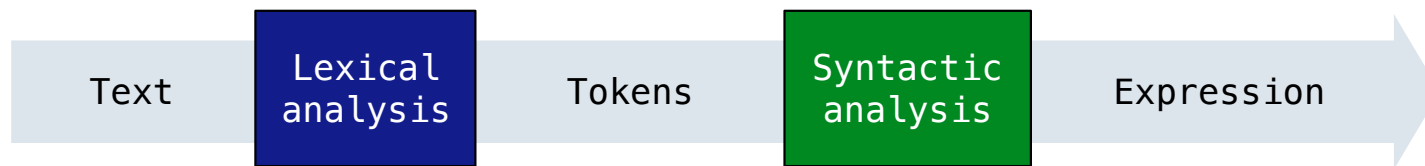
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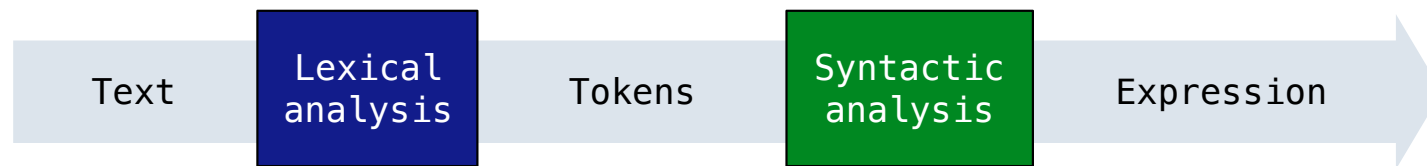
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(Demo)

Calculator

(Demo)

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    def __init__(self, first, second):
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The Pair Class

The Pair class represents Scheme pairs and lists. A list is a pair whose second element is either a list or nil.

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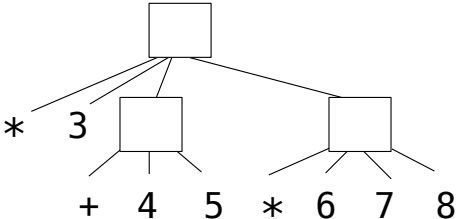
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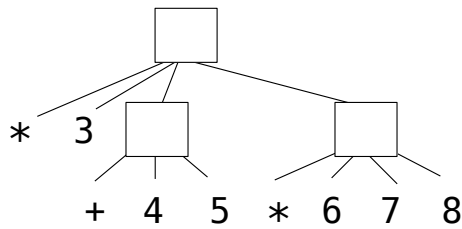
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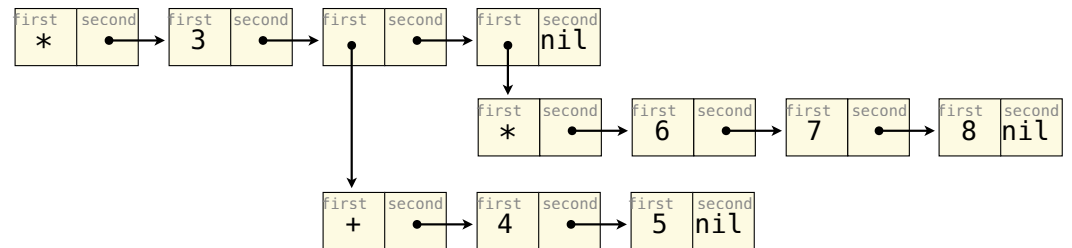
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Representation as Pairs



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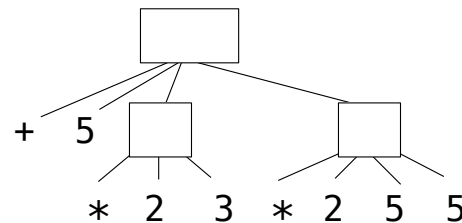
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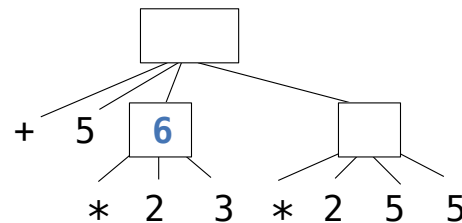
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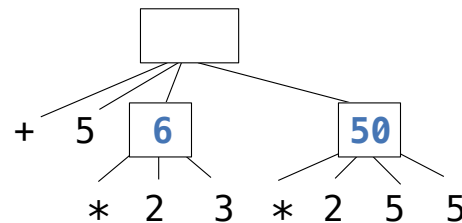
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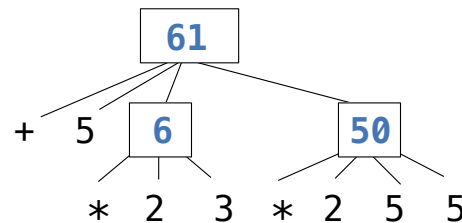
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Evaluation

The Eval Function

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```
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        ...
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-:
    ...
...
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(Demo)

Interactive Interpreters

Read-Eval-Print Loop

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