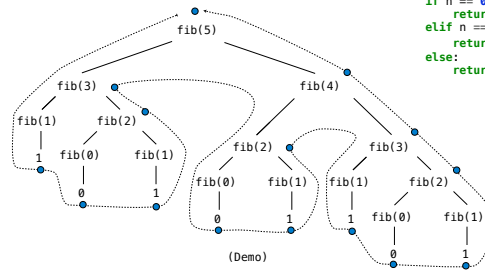# Efficiency

Announcements

Measuring Efficiency

Recursive Computation of the Fibonacci Sequence

Our first example of tree recursion:

```
def fib(n):
    if n == 0:
        return 0
    elif n == 1:
        return 1
    else:
        return fib(n-2) + fib(n-1)
```

fib(5)

fib(3) fib(4)

fib(1) fib(2)   fib(2)   fib(3)

1 fib(0) fib(1) fib(0) fib(1) fib(1) fib(2)

0   1   0   1   1 fib(0) fib(1)

0   1

(Demo)

Memoization

Memoization

Idea: Remember the results that have been computed before
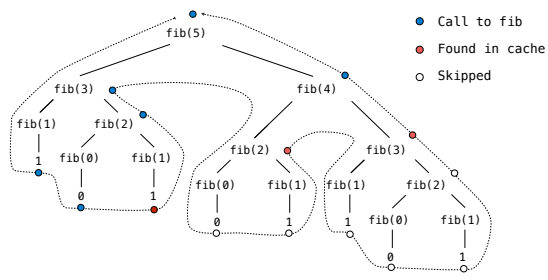
```
def memo(f):
    cache = {}
    def memoized(n):
        if n not in cache:
            cache[n] = f(n)
        return cache[n]
    return memoized
```

Keys are arguments that map to return values

Same behavior as f, if f is a pure function

(Demo)

Memoized Tree Recursion

fib(5)

fib(3)   fib(4)

fib(1) fib(2)   fib(2)   fib(3)

1 fib(0) fib(1) fib(0) fib(1) fib(1) fib(2)

0   1   0   1   1 fib(0) fib(1)

0   1

- ● Call to fib
- ● Found in cache
- ○ Skipped

Exponentiation

Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
    if n == 0:
        return 1
    else:
        return b * exp(b, n-1)
```

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ b \cdot b^{n-1} & \text{otherwise} \end{cases}$$

```
def exp_fast(b, n):
    if n == 0:
        return 1
    elif n % 2 == 0:
        return square(exp_fast(b, n//2))
    else:
        return b * exp_fast(b, n-1)
```

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ (b^{\frac{1}{2}n})^2 & \text{if } n \text{ is even} \\ b \cdot b^{n-1} & \text{if } n \text{ is odd} \end{cases}$$

```
def square(x):
    return x * x
```

(Demo)

Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
    if n == 0:
        return 1
    else:
        return b * exp(b, n-1)
```

Linear time:
• Doubling the input **doubles** the time
• 1024x the input takes 1024x as much time

```
def exp_fast(b, n):
    if n == 0:
        return 1
    elif n % 2 == 0:
        return square(exp_fast(b, n//2))
    else:
        return b * exp_fast(b, n-1)
```

Logarithmic time:
• Doubling the input **increases** the time by a constant C
• 1024x the input increases the time by only 10 times C

```
def square(x):
    return x * x
```

Orders of Growth

Quadratic Time

Functions that process all pairs of values in a sequence of length n take quadratic time

```
def overlap(a, b):
    count = 0
    for item in a:
        for other in b:
            if item == other:
                count += 1
    return count

overlap([3, 5, 7, 6], [4, 5, 6, 5])
```
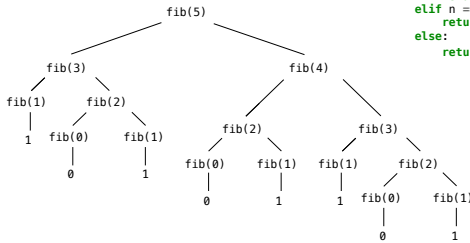
	3	5	7	6
4	0	0	0	0
5	0	1	0	0
6	0	0	0	1
5	0	1	0	0

(Demo)

Exponential Time

Tree-recursive functions can take exponential time

```
def fib(n):
    if n == 0:
        return 0
    elif n == 1:
        return 1
    else:
        return fib(n-2) + fib(n-1)
```


fib(5)
fib(3)   fib(4)
fib(1) fib(2)   fib(2)   fib(3)
1   fib(0) fib(1) fib(0) fib(1) fib(1) fib(2)
    0   1   0   1   1   fib(0) fib(1)
                                       0   1

http://en.wikipedia.org/wiki/File:Fibonacci.jpg

Common Orders of Growth

Time for n+n	Time for input n+1	Time for input n

Exponential growth. E.g., recursive fib
Incrementing n multiplies *time* by a constant

$$a \cdot b^{n+1} = (a \cdot b^n) \cdot b$$

Quadratic growth. E.g., overlap
Incrementing n increases *time* by n times a constant

$$a \cdot (n+1)^2 = (a \cdot n^2) + a \cdot (2n+1)$$

Linear growth. E.g., slow exp
Incrementing n increases *time* by a constant

$$a \cdot (n+1) = (a \cdot n) + a$$

Logarithmic growth. E.g., exp_fast
Doubling n only increments *time* by a constant

$$a \cdot \ln(2 \cdot n) = (a \cdot \ln n) + a \cdot \ln 2$$

Constant growth. Increasing n doesn't affect time

Space

Space and Environments

Which environment frames do we need to keep during evaluation?

At any moment there is a set of active environments

Values and frames in active environments consume memory

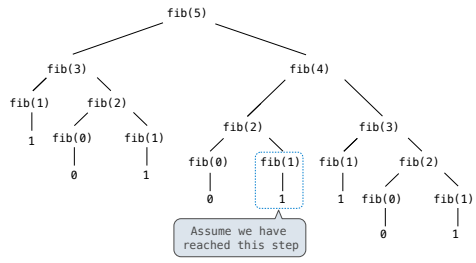Memory that is used for other values and frames can be recycled

Active environments:

•Environments for any function calls currently being evaluated

•Parent environments of functions named in active environments

(Demo)

fib(5)

fib(3)   fib(4)

fib(1)   fib(2)   fib(2)   fib(3)

1   fib(0)   fib(1)   fib(0)   fib(1)   fib(1)   fib(2)

0   1   0   1   1   fib(0)   fib(1)

0   1

Assume we have
reached this step

Has an active environment
Can be reclaimed
Hasn't yet been created

fib(5)

fib(3)   fib(4)

fib(1)   fib(2)   fib(2)   fib(3)

1   fib(0)   fib(1)   fib(0)   fib(1)   fib(1)   fib(2)

0   1   0   1   1   fib(0)   fib(1)

0   1

Assume we have
reached this step