

## 61A Lecture 27

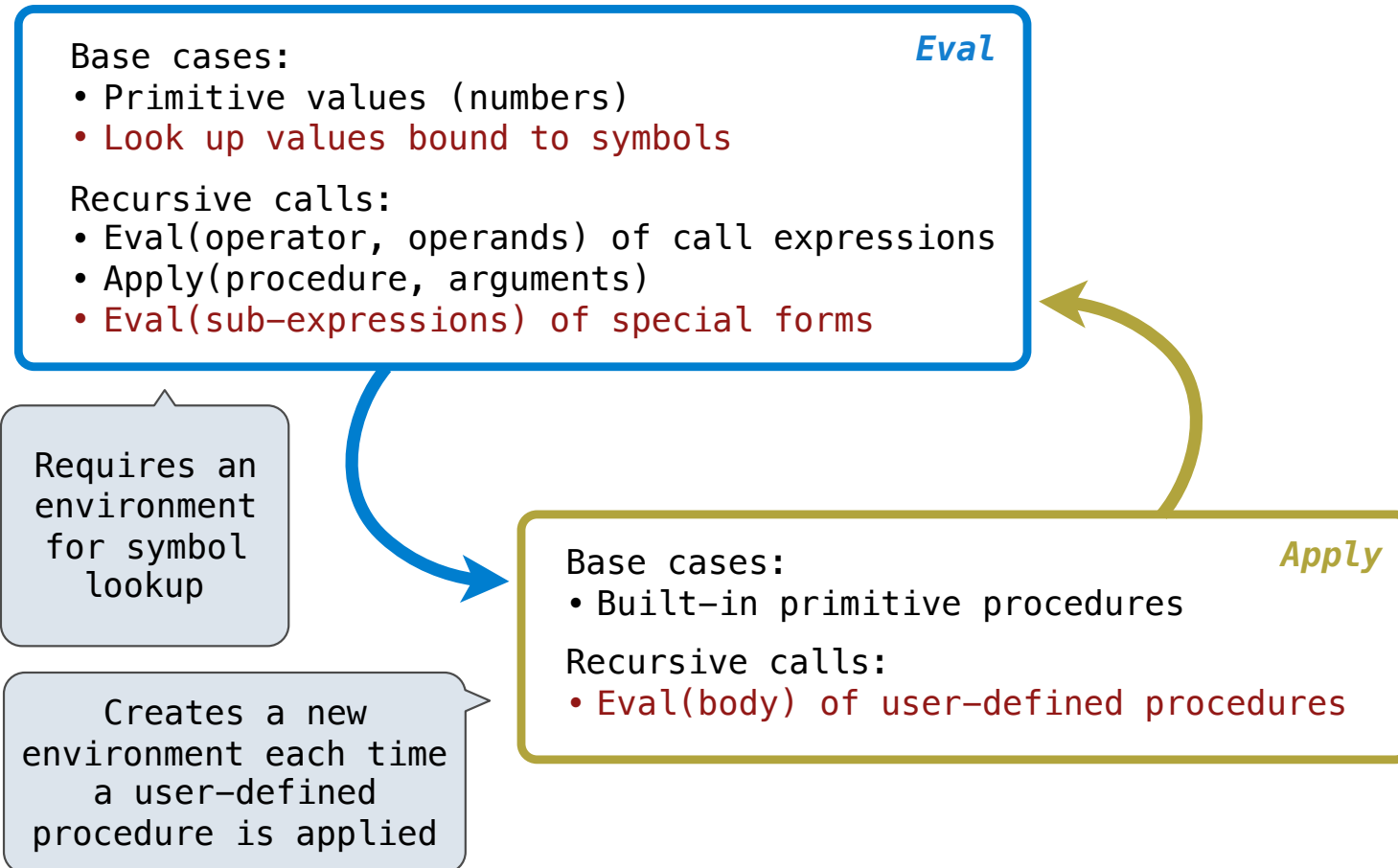
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## Announcements

## Interpreting Scheme

## The Structure of an Interpreter

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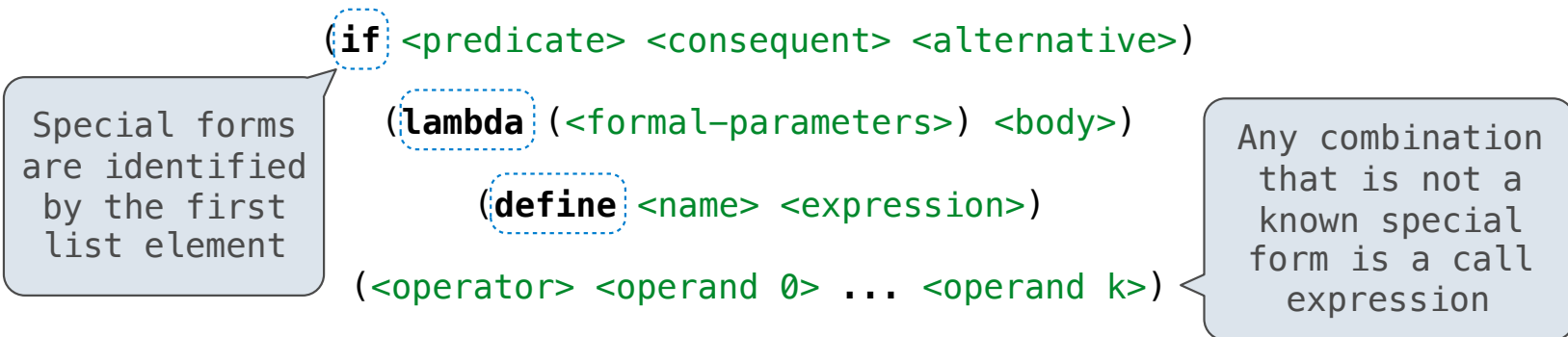
## Special Forms

## Scheme Evaluation

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The `scheme_eval` function choose behavior based on expression form:

- Symbols are looked up in the current environment
- Self-evaluating expressions are returned as values
- All other legal expressions are represented as Scheme lists, called combinations



```
(define (demo s) (if (null? s) '(3) (cons (car s) (demo (cdr s)))))
```

```
(demo (list 1 2))
```

# Logical Forms

## Logical Special Forms

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Logical forms may only evaluate some sub-expressions

- **If** expression:        `(if <predicate> <consequent> <alternative>)`
- **And** and **or**:        `(and <e1> ... <en>),     (or <e1> ... <en>)`
- **Cond** expression:     `(cond (<p1> <e1>) ... (<pn> <en>) (else <e>))`

The value of an if expression is the value of a sub-expression:

- Evaluate the predicate
- Choose a sub-expression: <consequent> or <alternative>
- Evaluate that sub-expression to get the value of the whole expression

do\_if\_form

(Demo)



Quotation

## Quotation


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The quote special form evaluates to the quoted expression, which is not evaluated

`(quote <expression>)`

`(quote (+ 1 2))`

evaluates to the  
three-element Scheme list



`(+ 1 2)`

The `<expression>` itself is the value of the whole quote expression

`'<expression>` is shorthand for `(quote <expression>)`

`(quote (1 2))`

is equivalent to

`'(1 2)`

The `scheme_read` parser converts shorthand `'` to a combination that starts with quote

(Demo)

# Lambda Expressions

## Lambda Expressions

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Lambda expressions evaluate to user-defined procedures

```
(lambda (<formal-parameters>) <body>)
```

```
(lambda (x) (* x x))
```

```
class LambdaProcedure:
```

```
    def __init__(self, formals, body, env):
```

```
        self.formals = formals ..... A scheme list of symbols
```

```
        self.body = body ..... A scheme list of expressions
```

```
        self.env = env ..... A Frame instance
```

## Frames and Environments

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A frame represents an environment by having a parent frame

Frames are Python instances with methods **lookup** and **define**

In Project 4, Frames do not hold return values

g: Global frame

y	3
z	5

f1: [parent=g]

x	2
z	4

(Demo)

## Define Expressions

## Define Expressions

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Define binds a symbol to a value in the first frame of the current environment.

```
(define <name> <expression>)
```

1. Evaluate the <expression>
2. Bind <name> to its value in the current frame

```
(define x (+ 1 2))
```

Procedure definition is shorthand of define with a lambda expression

```
(define (<name> <formal parameters>) <body>)
```

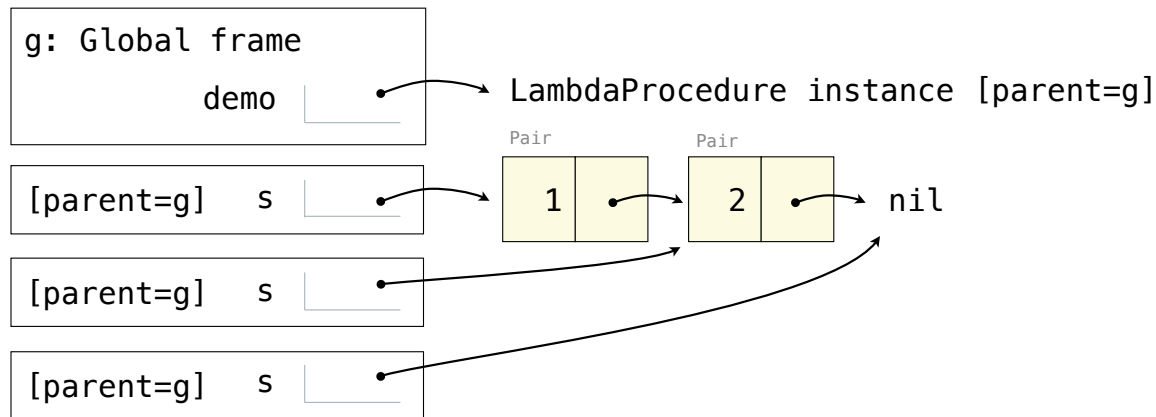
```
(define <name> (lambda (<formal parameters>) <body>))
```

## Applying User-Defined Procedures

To apply a user-defined procedure, create a new frame in which formal parameters are bound to argument values, whose parent is the **env** attribute of the procedure

Evaluate the body of the procedure in the environment that starts with this new frame

```
(define (demo s) (if (null? s) '(3) (cons (car s) (demo (cdr s)))))  
  
      (demo (list 1 2))
```





## Eval/Apply in Lisp 1.5

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```
apply[fn;x;a] =
  [atom[fn] → [eq[fn;CAR] → caar[x];
               eq[fn;CDR] → cdar[x];
               eq[fn;CONS] → cons[car[x];cadr[x]];
               eq[fn;ATOM] → atom[car[x]];
               eq[fn;EQ] → eq[car[x];cadr[x]];
               T → apply[eval[fn;a];x;a]];
  eq[car[fn];LAMBDA] → eval[caddr[fn];pairlis[cadr[fn];x;a]];
  eq[car[fn];LABEL] → apply[caddr[fn];x;cons[cons[cadr[fn];
                                                    caddr[fn]];a]]]

eval[e;a] = [atom[e] → cdr[assoc[e;a]];
            atom[car[e]] →
              [eq[car[e],QUOTE] → cadr[e];
               eq[car[e];COND] → evcon[cdr[e];a];
               T → apply[car[e];evlis[cdr[e];a];a]];
            T → apply[car[e];evlis[cdr[e];a];a]]
```