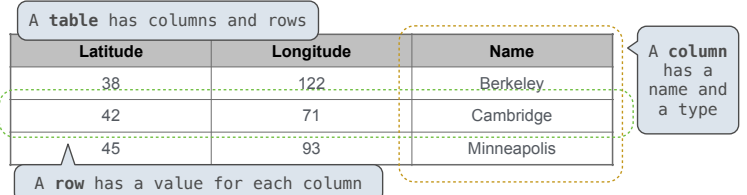


```
(car (cons 1 nil)) => 1
(cdr (cons 1 nil)) => ()
(cdr (cons 1 (cons 2 nil))) => (2)
```

```
(define size 5) ; => size
(* 2 size) ; => 10
(if (> size 0) size (- size)) ; => 5
(cond ((> size 0) size) ((= size 0) 0) (else (- size))) ; => 5
(lambda (x y) (+ x y size)) size (+ 1 2) ; => 13
(let ((a size) (b (+ 1 2))) (* 2 a b)) ; => 30
(map (lambda (x) (+ x size)) (quote (2 3 4))) ; => (7 8 9)
(filter odd? (quote (2 3 4))) ; => (3)
(list (cons 1 nil) size 'size) ; => ((1) 5 size)
(list (equal? 1 2) (null? nil) (= 3 4) (eq? 5 5)) ; => (#f #t #f #t)
(list (or #f #t) (or) (or 1 2)) ; => (#t #f 1)
(list (and #f #t) (and) (and 1 2)) ; => (#f #t 2)
(append '(1 2) '(3 4)) ; => (1 2 3 4)
(not (> 1 2)) ; => #t
(begin (define x (+ size 1)) (* x 2)) ; => 12
(+ size (- ,size) ,(* 3 4)) ; => (+ size (- 5) 12)
```

```
;; Return a copy of s reversed.
(define (reverse s)
  (define (iter s r)
    (if (null? s) r
        (iter (cdr s) (cons (car s) r))))
  (iter s nil))

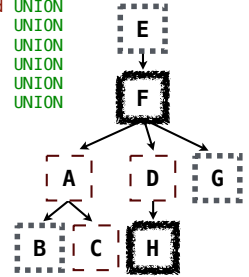
;; Apply fn to each element of s.
(define (map fn s)
  (define (map-reverse s m)
    (if (null? s) m
        (map-reverse (cdr s) (cons (fn (car s)) m))))
  (reverse (map-reverse s nil)))
```



```
SELECT [expression] AS [name], [expression] AS [name], ... ;
SELECT [columns] FROM [table] WHERE [condition] ORDER BY [order];
```

```
CREATE TABLE parents AS
SELECT "abraham" AS parent, "barack" AS child UNION
SELECT "abraham" , "clinton" UNION
SELECT "delano" , "herbert" UNION
SELECT "fillmore" , "abraham" UNION
SELECT "fillmore" , "delano" UNION
SELECT "fillmore" , "grover" UNION
SELECT "eisenhower" , "fillmore";
```

```
CREATE TABLE dogs AS
SELECT "abraham" AS name, "long" AS fur UNION
SELECT "barack" , "short" UNION
SELECT "clinton" , "long" UNION
SELECT "delano" , "long" UNION
SELECT "eisenhower" , "short" UNION
SELECT "fillmore" , "curly" UNION
SELECT "grover" , "short" UNION
SELECT "herbert" , "curly";
```



```
SELECT a.child AS first, b.child AS second
FROM parents AS a, parents AS b
WHERE a.parent = b.parent AND a.child < b.child;
```

First	Second
barack	clinton
abraham	delano
abraham	grover
delano	grover

Optional content removed

The number of groups is the number of unique values of an expression
 A having clause filters the set of groups that are aggregated

```
select weight/legs, count(*) from animals
group by weight/legs
having count(*)>1;
```

weight/legs	count(*)
5	2
2	2

- weight/legs=5
- weight/legs=2
- weight/legs=2
- weight/legs=3
- weight/legs=5
- weight/legs=6000

kind	legs	weight
dog	4	20
cat	4	10
ferret	4	10
parrot	2	6
penguin	2	10
t-rex	2	12000

Optional content removed

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2 3.3, true, +, quotient, ...
- Combinations: (quotient 10 2), (not true), ...

Numbers are self-evaluating; *symbols* are bound to values. Call expressions have an operator and 0 or more operands.

A combination that is not a call expression is a *special form*:

- If expression: (if <predicate> <consequent> <alternative>)
- Binding names: (define <name> <expression>)
- New procedures: (define (<name> <formal parameters>) <body>)

```

> (define pi 3.14)      > (define (abs x)
> (* pi 2)              >   (- x 0)
6.28                    >   (if (< x 0)
                          >     (- x)
                          >     x))
                          > (abs -3)
                          3
    
```

Lambda expressions evaluate to anonymous procedures.

```
(lambda (<formal-parameters>) <body>)
```

Two equivalent expressions:

```
(define (plus4 x) (+ x 4))
(define plus4 (lambda (x) (+ x 4)))
```

An operator can be a combination too:

```
((lambda (x y z) (+ x y (square z))) 1 2 3)
```



In the late 1950s, computer scientists used confusing names.

- **cons**: Two-argument procedure that **creates a pair**
 - **car**: Procedure that returns the **first element** of a pair
 - **cdr**: Procedure that returns the **second element** of a pair
 - **nil**: The empty list
- They also used a non-obvious notation for linked lists.
- A (linked) Scheme list is a pair in which the second element is nil or a Scheme list.
 - Scheme lists are written as space-separated combinations.
 - A dotted list has an arbitrary value for the second element of the last pair. Dotted lists may not be well-formed lists.

```

> (define x (cons 1 nil))
> x
(1)
> (car x)
1
> (cdr x)
()
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))
(1 2 3 4)
    
```

Symbols normally refer to values; how do we refer to symbols?

```

> (define a 1)
> (define b 2)
> (list a b)
(1 2)
    
```

No sign of "a" and "b" in the resulting value

Quotation is used to refer to symbols directly in Lisp.

```

> (list 'a 'b)
(a b)
> (list 'a b)
(a 2)
    
```

Symbols are now values

Quotation can also be applied to combinations to form lists.

```

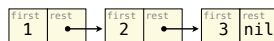
> (car '(a b c))
a
> (cdr '(a b c))
(b c)
    
```

```

class Pair:
    """A pair has two instance attributes:
        first and rest.
    """
    rest must be a Pair or nil.
    """
    def __init__(self, first, rest):
        self.first = first
        self.rest = rest
    
```

```

>>> s = Pair(1, Pair(2, Pair(3, nil)))
>>> s
Pair(1, Pair(2, Pair(3, nil)))
>>> print(s)
(1 2 3)
    
```

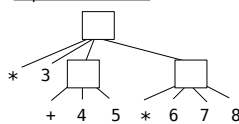


The Calculator language has primitive expressions and call expressions

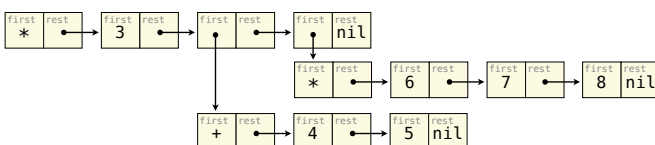
Calculator Expression

```
(* 3 (+ 4 5) (* 6 7 8))
```

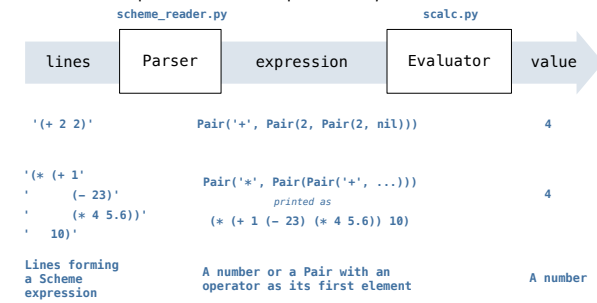
Expression Tree



Representation as Pairs



A basic interpreter has two parts: a *parser* and an *evaluator*.



A Scheme list is written as elements in parentheses:

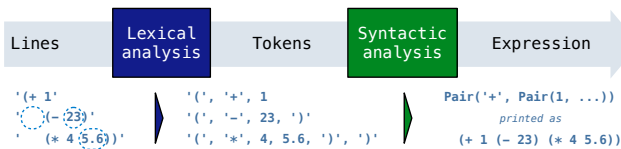
```
(<element> <element> ... <element>)
```

A Scheme list

Each <element> can be a combination or atom (primitive).
(+ (* 3 (+ (* 2 4) (+ 3 5))) (+ (- 10 7) 6))

The task of *parsing* a language involves coercing a string representation of an expression to the expression itself. Parsers must validate that expressions are well-formed.

A Parser takes a sequence of lines and returns an expression.



- Iterative process
 - Checks for malformed tokens
 - Determines types of tokens
 - Processes one line at a time
- Tree-recursive process
 - Balances parentheses
 - Returns tree structure
 - Processes multiple lines

Syntactic analysis identifies the hierarchical structure of an expression, which may be nested.

Each call to scheme_read consumes the input tokens for exactly one expression.

Base case: symbols and numbers
Recursive call: scheme_read sub-expressions and combine them

Base cases:

- Primitive values (numbers)
- Look up values bound to symbols

Recursive calls:

- Eval(operator, operands) of call expressions
- Apply(procedure, arguments)
- Eval(sub-expressions) of special forms

Eval

The structure of the Scheme interpreter

Creates a new environment each time a user-defined procedure is applied

Apply

Base cases:

- Built-in primitive procedures

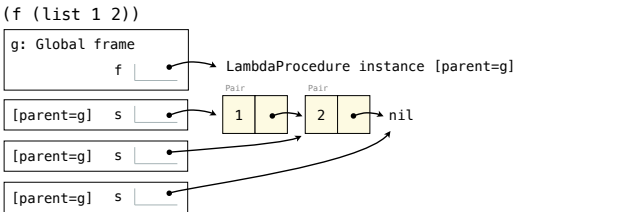
Recursive calls:

- Eval(body) of user-defined procedures

Requires an environment for name lookup

To apply a user-defined procedure, create a new frame in which formal parameters are bound to argument values, whose parent is the env of the procedure, then evaluate the body of the procedure in the environment that starts with this new frame.

```
(define (f s) (if (null? s) '(3) (cons (car s) (f (cdr s)))))
```



How to Design Functions:

- 1) Identify the information that must be represented and how it is represented. Illustrate with examples.
- 2) State what kind of data the desired function consumes and produces. Formulate a concise answer to the question *what* the function computes.
- 3) Work through examples that illustrate the function's purpose.
- 4) Outline the function as a template.
- 5) Fill in the gaps in the function template. Exploit the purpose statement and the examples.
- 6) Convert examples into tests and ensure that the function passes them.