Interpreters

Announcements

Interpreting Scheme

The Structure of an Interpreter

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    self.env = env ........................................................... A Frame instance
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f1: [parent=g]

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| :--- | :--- |
|  | 4 |

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$$
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& \begin{array}{l|l}
y & 3 \\
z & 5
\end{array} \\
& \text { f1: [parent=g] } \\
& \begin{array}{l|l}
x & 2 \\
z & 4
\end{array} \\
& \text { (Demo) }
\end{aligned}
$$

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## Eval/Apply in Lisp 1.5

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```
\(\operatorname{apply}[\mathrm{fn} ; \mathbf{x} ; \mathrm{a}]=\)
    [atom \([\mathrm{fn}] \rightarrow[\mathrm{eq}[\mathrm{fn} ; \mathrm{CAR}] \rightarrow \operatorname{caar}[\mathrm{x}] ;\)
        eq[fn;CDR] \(\rightarrow\) cdar[x];
            eq[fn;CONS] \(\rightarrow \operatorname{cons[car[x];cadr[x]];~}\)
            eq[fn;ATOM] \(\rightarrow\) atom[car[x]];
            eq[fn;EQ] \(\rightarrow\) eq[car[x];cadr[x]];
            \(T \rightarrow \operatorname{apply}[e v a l[f n ; a] ; x ; a]] ;\)
    eq[car[fn];LAMBDA] \(\rightarrow\) eval[caddr[fn];pairlis[cadr[fn];x;a]];
    eq[car[fn];LABEL] \(\rightarrow\) apply[caddr[fn];x;cons[cons[cadr[fn];
                        caddr[fn]];a]]]
\(\operatorname{eval}[e ; a]=[\) atom \([e] \rightarrow \operatorname{cdr}[\operatorname{assoc}[e ; a]] ;\)
    atom[car[e]] \(\rightarrow\)
            [eq[car[e],QUOTE] \(\rightarrow\) cadr[e];
            eq[car[e];COND] \(\rightarrow\) evcon[cdr[e];a];
            T \(\rightarrow\) apply[car[e];evlis[cdr[e];a];a]];
    \(\mathrm{T} \rightarrow\) apply[car[e];evlis[cdr[e];a];a]]
```

