Exceptions are raised with a raise statement.
raise <expr>
<expr> must evaluate to a subclass of BaseException or an instance of one. try

| ```try: <try suite> except <exception class> as <name>: <except suite>``` | ```>>> try: x = 1/0 except ZeroDivisionError as e: print('handling a', type(e)) x = 0``` |
| :---: | :---: |
| The <try suite> is executed first. <br> If, during the course of executing the <try suite>, an exception is raised that is not handled otherwise, and | ```handling a <class 'ZeroDivisionError'> >>> x 0``` | that is not handled otherwise, and 0



If the class of the exception inherits from <exception class>, then The <except suite> is executed, with <name> bound to the exception.
(append $\mathbf{s} \mathbf{t}$ ): list the elements of $s$ and $t$; append can be called on more than 2 lists
(map $f$ s): call a procedure $f$ on each element of a list $s$ and list the results
(filter $f$ s): call a procedure $f$ on each element of a list $s$ and list the elements for which a true value is the result
(apply f s): call a procedure $f$ with the elements of a list as its arguments

```
(define size 5) ; => size
(* 2 size) ; => 10
(if (> size 0) size (- size)) ; => 5
(cond ((> size 0) size) ((= size 0) 0) (else (- size))) ; => 5
((lambda (x y) (+ x y size)) size (+ 1 2)) ; => 13
(let ((a size) (b (+1 2))) (* 2 a b)); => 30
(map (lambda (x) (+ x size)) (quote (2 3 4))) ; => (7 8 9)
(filter odd? (quote (2 3 4))) ; => (3)
(list (cons 1 nil) size 'size)'; => ((1) 5 size)
(list (equal? 1 2) (null? nil) (= 3 4) (eq? 5 5)) ; => (#f #t #f #t)
(list (or #f #t) (or) (or 1 2)) ; => (#t #f 1)
(list (and #f #t) (and) (and 1 2)) ; => (#f #t 2)
(list 'a 2) ; => (a 2)
(append '(1 2) '((3 4)) ; => (1 1 2 3 4)
(not (> 1 2)) ; => #t
(begin (define x (+ size 1)) (* x 2)) ; => }1
```

(define (factorial $n$ )
(if (= n 0) 1
$(* \mathrm{n}($ factorial $(-\mathrm{n} 1))))$
(define (fib n)
( cond
$(1=n$ 0) o)
$\left(=\begin{array}{ll}11) \\ 1)\end{array}\right.$
(else $(+(f i b(-n 2))(f i b(-n 1)))))$
(define (nines num)
(if (= num 0)
0
(if (= (modulo num 10) 9)
(+ 1 (nines (floor (/ num 10))))
(nines (floor (/ num 10))))))
The way in which names are looked up in Scheme and Python is
called lexical scope (or static scope).
Lexical scope: The parent of a frame is the environment in
which a procedure was defined. (lambda ...)
Dynamic scope: The parent of a frame is the environment in
which a procedure was called. (mu ...)
$>$ (define $f(\operatorname{mu}(x)(+x y)))$
$>$ (define g (lambda $(x y)(f(+x x))))$
$>\left(\begin{array}{ll}g & 3\end{array}\right)$
$>13$
There are two ways to quote an expression
$\left.\begin{array}{lll}\text { Quote: } & \text { ' (a b) } & => \\ \text { Quasiquote: } & (\mathrm{a} & \mathrm{a}\end{array} \mathrm{b}\right)$
Parts of a quasiquoted expression can be unquoted with ,
Quote: $\quad\left(\begin{array}{ll}\text { (define } b & 4) \\ (a,(+b 1))\end{array} \quad \Rightarrow \quad(a\right.$ (unquote $(+b 1))$
Quasiquote: ' $(\mathrm{a},(+\mathrm{b} 1))=>\quad$ ( a 5 )
Quasiquotation is convenient for generating Scheme expressions:
(define (make-add-lambda n) '(lambda (d) (+d, n)))
$\begin{array}{ll} & (\text { define (make-add-lambda } n)(l a m b d a \\ & (\text { make-add-lambda } 2)(+d, n)))\end{array}$
ept ZeroDivisionError as e:
print('handling a', type(e))
print
$\mathrm{x}=0$
except <exception class> as <name>:
<except suite>


| A table has columns and rows | Longitude | Name |
| :---: | :---: | :---: |
| Latitude | 122 | Berkeley |
| 38 | 71 | Cambridge |
| 42 | 93 | Minneapolis |
|  | 45 |  |
| A row has a value for each column |  |  |

A column
has a
name and
a type

SELECT [expression] AS [name], [expression] AS [name], ... ;
SELECT [columns] FROM [table] WHERE [condition] ORDER BY [order];
CREATE TABLE parents AS

create table lift as


String values can be combined to form longer strings

```
- sqlite> SELECT "hello," || " world";
    hello, world
```

Basic string manipulation is built into SQL, but differs from Python
sqlite> CREATE TABLE phrase AS SELECT "hello, world" AS s; sqlite> SELECT substr(s, 4, 2) \|| substr(s, instr(s, " ")+1, 1) FROM phrase;
low

The number of groups is the number of unique values of an expression A having clause filters the set of groups that are aggregated select weight/legs, count(*) from animals

|  |  | group by weight/legs having count(*)>1; | kind | legs | weight |
| :---: | :---: | :---: | :---: | :---: | :---: |
| weight/ legs | count(*) | weight/legs=5 | dog | 4 | 20 |
|  |  |  | cat | 4 | 10 |
| 5 | 2 | weight/legs=2 | ferret | 4 | 10 |
| 2 | 2 | weight/legs=3 | parrot | 2 | 6 |
|  |  | weight/legs=5 | penguin | 2 | 10 |
|  |  | weight/legs=6000 | t-rex | 2 | 12000 |



A macro is an operation performed on the source code of a program before evaluation
Macros exist in many languages, but are easiest to define correctly in a language like Lisp
Scheme has a define-macro special form that defines a source code transformation

```
(define-macro (twice expr)
    (list 'begin expr expr))
> (twice (print 2))}>\mathrm{ (begin (print 2) (print 2))
2
```

Evaluation procedure of a macro call expression:

- Evaluate the operator sub-expression, which evaluates to a macro
- Call the macro procedure on the operand expressions without evaluating them first
- Evaluate the expression returned from the macro procedure

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, + quotient
- Combinations: (quotient 10 2), (not true)

Numbers are self-evaluating; symbols are bound to values.
Call expressions have an operator and 0 or more operands.
A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- Binding names: (define <name> <expression>)
- New procedures: (define (<name> <formal parameters>) <body>)

| $\begin{aligned} & >(\text { define pi 3.14) } \\ & >(* \text { pi } 2) \\ & 6.28 \end{aligned}$ | > (define (abs x) |
| :---: | :---: |
|  | (if ( $<x$ ) |
|  | ( -x ) |
|  | ( x) ) |
|  | $>(a b s-3)$ |
|  | 3 |

Lambda expressions evaluate to anonymous procedures.
(lambda (<formal-parameters>) <body>)
Two equivalent expressions:
(define (plus4 x) (+ x 4))

(define plus4 (lambda (x) (+ x 4)))
An operator can be a combination too:
((lambda (x y z) (+ x y (square z))) 12 3)

> In the late 1950s, computer scientists used confusing names.
> - cons: Two-argument procedure that creates a pair
> - car: Procedure that returns the first element of a pair
> - cdr: Procedure that returns the second element of a pair
> - nil: The empty list
> They also used a non-obvious notation for linked lists
> - A (linked) Scheme list is a pair in which the second element is
> nil or a Scheme list.
> - Scheme lists are written as space-separated combinations.
> - A dotted list has an arbitrary value for the second element of the last pair. Dotted lists may not be well-formed lists.
> $>$ (define $x$ (cons 1 nil))
> $\begin{aligned} & > \\ & \text { (1) }\end{aligned}$
> $>(\operatorname{car} \mathrm{x})$
> $>(\operatorname{cdr} x)$
> $>($ cons $1($ cons $2($ cons 3 (cons 4 nil))))
> (1 23 4)
> Symbols normally refer to values; how do we refer to symbols?
> Quotation is used to refer to symbols directly in Lisp.
> Quotation can also be applied to combinations to form lists. $>\left(\operatorname{car}^{\prime}(\mathrm{a} b \mathrm{c})\right)$
> $>\left(c d r{ }^{\prime}(a \quad b c)\right)$
> (b c)
$\begin{array}{ll}(c a r & (c o n s \\ 1 & \text { nil) }) \\ \text { (cdr }(\text { cons } 1 \text { nil) }) & \text {-> () }\end{array}$
(cdr (cons 1 (cons 2 nil))) $\rightarrow$ (2)

## class Pair:

""A pair has two instance attributes:
first and rest.
rest must be a Pair or nil.
def _init_(self, first, rest): self.first $=$ first self.rest $=$ rest
$\ggg \mathbf{s}=\operatorname{Pair}(1, \operatorname{Pair}(2, \operatorname{Pair}(3, \operatorname{nil})))$
>>> s
Pair(1, Pair(2, Pair(3, nil)))
>>> print(s)
(1 2 3)
The Calculator language has primitive expressions and call expressions
Calculator Expression


$$
\begin{aligned}
& (* 3 \\
& \quad(+45) \\
& \quad(* 678))
\end{aligned}
$$



Representation as Pairs


A basic interpreter has two parts: a parser and an evaluator.

expression
operator as its first element
A number
A Scheme list is written as elements in parentheses:


Each <element> can be a combination or atom (primitive).
(+ (* 3 (+ (* 24 ) (+ 3 5))) (+ (- 10 7) 6))
The task of parsing a language involves coercing a string representation of an expression to the expression itself. Parsers must validate that expressions are well-formed.
A Parser takes a sequence of lines and returns an expression.


## - Iterative process <br> - Checks for malformed tokens <br> - Determines types of tokens <br> - Processes one line at a time

- Tree-recursive process
- Balances parentheses
- Returns tree structure
- Processes multiple lines

Syntactic analysis identifies the hierarchical structure of an expression, which may be nested.
Each call to scheme_read consumes the input tokens for exactly one expression.

Base case: symbols and numbers
Recursive call: scheme_read sub-expressions and combine them


To apply a user-defined procedure, create a new frame in which formal parameters are bound to argument values, whose parent is the env of the procedure, then evaluate the body of the procedure in the environment that starts with this new frame. (define (f s) (if (null? s) '(3) (cons (car s) (f (cdr s))))) (f (list 12 ))


