

CS61A Lecture 23

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Announcements



- Ants project due Monday

- HW8 due next Wednesday at 7pm

- Midterm 2 next Thursday at 7pm
 - Review session Sat. 3/16 at 2pm in 2050 VLSB
 - Office hours Sun. 3/17 12-4pm in 310 Soda
 - HKN review session Sun. 3/17 at 4pm in 145 Dwinelle
 - See course website for more information

The Independence of Data Types



Data abstraction and class definitions keep types separate

Some operations need to cross type boundaries

*How do we add a complex number
and a rational number together?*

— `add_rational mul_rational` —

*Rational numbers as
numerators & denominators*

`add_complex mul_complex` —

*Complex numbers as
two-dimensional vectors*

There are many different techniques for doing this!

Type Dispatching



Define a different function for each possible combination of types for which an operation (e.g., addition) is valid

```
def iscomplex(z):
    return type(z) in (ComplexRI, ComplexMA)
def isrational(z):
    return type(z) is Rational
def add_complex_and_rational(z, r):
    return ComplexRI(z.real + r.numerator / r.denominator,
                    z.imag)
def add_by_type_dispatching(z1, z2):
    """Add z1 and z2, which may be complex or rational."""
    if iscomplex(z1) and iscomplex(z2):
        return add_complex(z1, z2)
    elif iscomplex(z1) and isrational(z2):
        return add_complex_and_rational(z1, z2)
    elif isrational(z1) and iscomplex(z2):
        return add_complex_and_rational(z2, z1)
    else:
        add_rational(z1, z2)
```

Converted to a real number (float)

Tag-Based Type Dispatching



Idea: Use dictionaries to dispatch on type (like we did for message passing)

```
def type_tag(x):  
    return type_tags[type(x)]
```

```
type_tags = {ComplexRI: 'com',  
             ComplexMA: 'com',  
             Rational:  'rat'}
```

Declares that ComplexRI and ComplexMA should be treated uniformly

```
def add(z1, z2):  
    types = (type_tag(z1), type_tag(z2))  
    return add_implementations[types](z1, z2)
```

```
add_implementations = {}  
add_implementations[('com', 'com')] = add_complex  
add_implementations[('rat', 'rat')] = add_rational  
add_implementations[('com', 'rat')] = add_complex_and_rational  
add_implementations[('rat', 'com')] = add_rational_and_complex
```

```
lambda r, z: add_complex_and_rational(z, r)
```

Type Dispatching Analysis



Minimal violation of abstraction barriers: we define cross-type functions as necessary, but use abstract data types

Extensible: Any new numeric type can "install" itself into the existing system by adding new entries to various dictionaries

```
def add(z1, z2):  
    types = (type_tag(z1), type_tag(z2))  
    return add_implementations[types](z1, z2)
```

Question: How many cross-type implementations are required to support m types and n operations?

integer, rational, real,
complex

$$m \cdot (m - 1) \cdot n$$

add, subtract, multiply,
divide

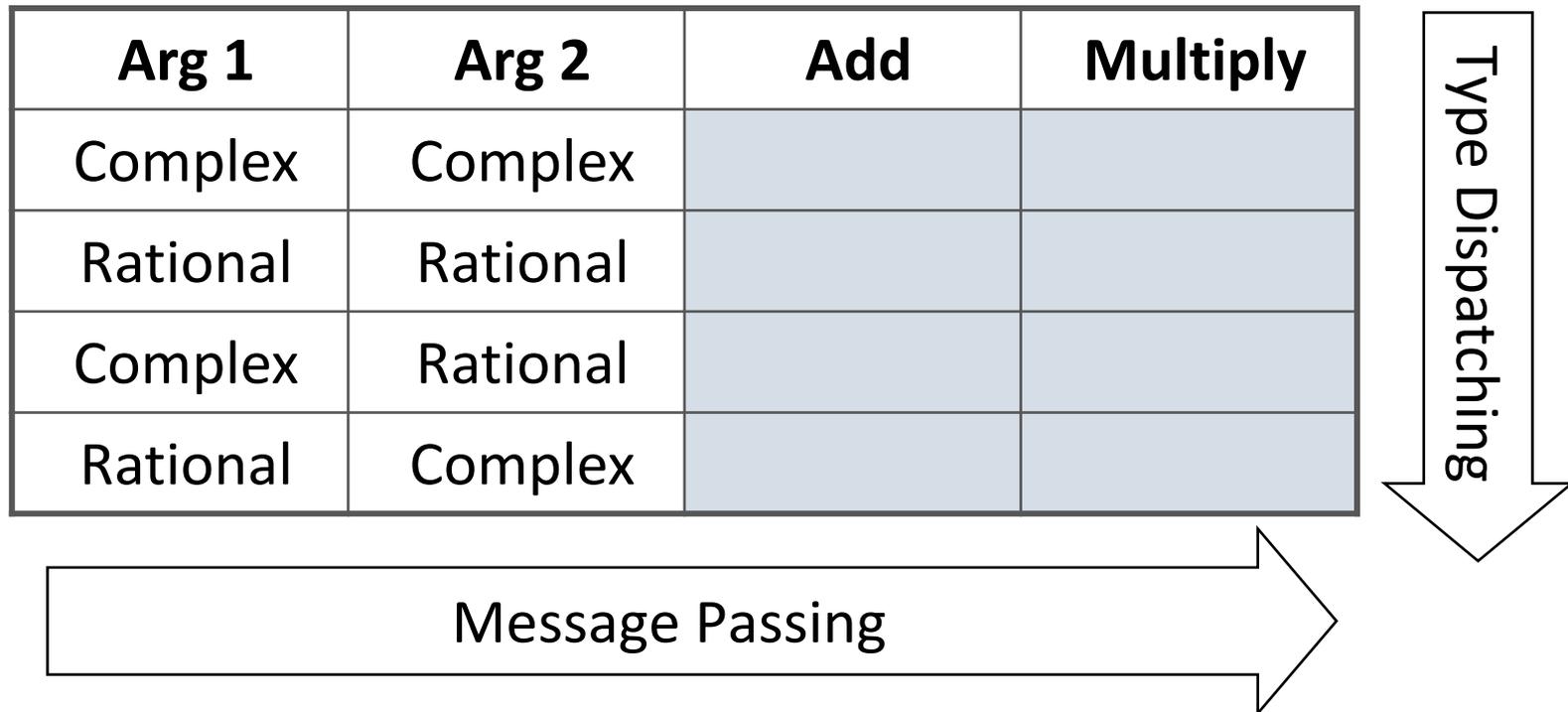
$$4 \cdot (4 - 1) \cdot 4 = 48$$

Type Dispatching Analysis



Minimal violation of abstraction barriers: we define cross-type functions as necessary, but use abstract data types

Extensible: Any new numeric type can "install" itself into the existing system by adding new entries to various dictionaries



Data-Directed Programming



There's nothing addition-specific about `add`

Idea: One dispatch function for (operator, types) pairs

```
def apply(operator_name, x, y):
    tags = (type_tag(x), type_tag(y))
    key = (operator_name, tags)
    return apply_implementations[key](x, y)
```

```
apply_implementations = {
    ('add', ('com', 'com')): add_complex,
    ('add', ('rat', 'rat')): add_rational,
    ('add', ('com', 'rat')): add_complex_and_rational,
    ('add', ('rat', 'com')): add_rational_and_complex,
    ('mul', ('com', 'com')): mul_complex,
    ('mul', ('rat', 'rat')): mul_rational,
    ('mul', ('com', 'rat')): mul_complex_and_rational,
    ('mul', ('rat', 'com')): mul_rational_and_complex
}
```

Coercion



Idea: Some types can be converted into other types

Takes advantage of structure in the type system

```
def rational_to_complex(x):  
    return ComplexRI(x.numerator / x.denominator, 0)  
  
coercions = {('rat', 'com'): rational_to_complex}
```

Question: Can any numeric type be coerced into any other?

Question: Have we been repeating ourselves with data-directed programming?

Applying Operators with Coercion



1. Attempt to coerce arguments into values of the same type
2. Apply type-specific (not cross-type) operations

```
def coerce_apply(operator_name, x, y):
    tx, ty = type_tag(x), type_tag(y)
    if tx != ty:
        if (tx, ty) in coercions:
            tx, x = ty, coercions[(tx, ty)](x)
        elif (ty, tx) in coercions:
            ty, y = tx, coercions[(ty, tx)](y)
        else:
            return 'No coercion possible.'
    assert tx == ty
    key = (operator_name, tx)
    return coerce_implementations[key](x, y)
```

Coercion Analysis



Minimal violation of abstraction barriers: we define cross-type coercion as necessary, but use abstract data types

Requires that all types can be coerced into a common type

More sharing: All operators use the same coercion scheme

Arg 1	Arg 2	Add	Multiply
Complex	Complex		
Rational	Rational		
Complex	Rational		
Rational	Complex		



From	To	Coerce
Complex	Rational	
Rational	Complex	

Type	Add	Multiply
Complex		
Rational		

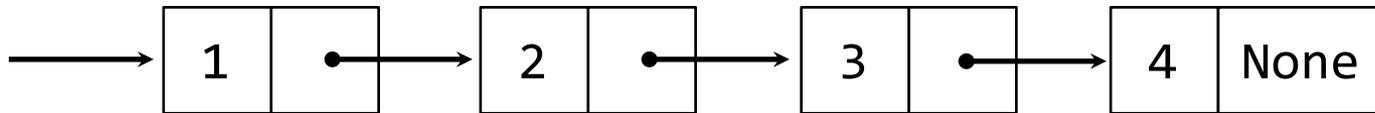
Closure Property of Data



A tuple can contain another tuple as an element.

Pairs are sufficient to represent sequences.

Recursive list representation of the sequence 1, 2, 3, 4:



Recursive lists are recursive: the rest of the list is a list.

Nested pairs (old): `(1, (2, (3, (4, None))))`

Rlist class (new): `Rlist(1, Rlist(2, Rlist(3, Rlist(4))))`

Recursive List Class



Methods can be recursive as well!

```
class Rlist(object):  
    class EmptyList(object):  
        def __len__(self):  
            return 0  
    empty = EmptyList()  
    def __init__(self, first, rest=empty):  
        self.first = first  
        self.rest = rest  
    def __len__(self):  
        return 1 + len(self.rest)  
    def __getitem__(self, i):  
        if i == 0:  
            return self.first  
        return self.rest[i - 1]
```

There's the
base case!

Yes, this call is
recursive

Recursive Operations on Rlists



Recursive list processing almost always involves a recursive call on the rest of the list.

```
>>> s = Rlist(1, Rlist(2, Rlist(3)))
```

```
>>> s.rest  
Rlist(2, Rlist(3))
```

```
>>> extend_rlist(s.rest, s)  
Rlist(2, Rlist(3, Rlist(1, Rlist(2, Rlist(3)))))
```

```
def extend_rlist(s1, s2):  
    if s1 is Rlist.empty:  
        return s2  
    return Rlist(s1.first, extend_rlist(s1.rest, s2))
```

Map and Filter on Rlists



We want operations on a whole list, not an element at a time.

```
def map_rlist(s, fn):
    if s is Rlist.empty:
        return s
    return Rlist(fn(s.first), map_rlist(s.rest, fn))
```

```
def filter_rlist(s, fn):
    if s is Rlist.empty:
        return s
    rest = filter_rlist(s.rest, fn)
    if fn(s.first):
        return Rlist(s.first, rest)
    return rest
```