

Instructions

Form a small group. Start on the first problem. Check off with a helper or discuss your *solution process* with another group once everyone understands *how to solve* the first problem and then repeat for the second problem ...

You may not move to the next problem until you check off or discuss with another group and *everyone understands why the solution is what it is*. You may use any course resources at your disposal: the purpose of this review session is to have everyone learning together as a group.

0.1 What would Python display?

```
>>> pikachu, charmander = 'electric', 'fire'
>>> ash = [[pikachu], [charmander], [[pikachu]]]
>>> pikachu, charmander = 2, 0
>>> ash[pikachu] = [ash, ash[pikachu][charmander]]
>>> ash
```

1 Lists & Tree Recursion

Mutative (*destructive*) operations change the state of a list by adding, removing, or otherwise modifying the list itself.

- `lst.append(element)`
- `lst.extend(lst)`
- `lst.pop(index)`
- `lst += lst` (**not** `lst = lst + lst`)
- `lst[i] = x`
- `lst[i:j] = lst`

Non-mutative (*non-destructive*) operations include the following.

- `lst + lst`
- `lst * n`
- `lst[i:j]`
- `list(lst)`

2 Midterm 2 Review

Recall: To execute assignment statements,

- Evaluate all expressions to the right of the = sign
- Bind all names to the left of the = to those resulting values

The **Golden Rule of Equals** describes how this rule behaves with composite values. *Composite values*, such as functions and lists, are connected by a pointer. When an expression evaluates to a composite value, we are returned the pointer to that value, rather than the value itself.

In an environment diagram, we can summarize this rule with,

Copy *exactly* what is in the box!

1.1 Write a list comprehension that accomplishes each of the following tasks.

- Square all the elements of a given list, `lst`.
- Compute the dot product of two lists `lst1` and `lst2`. *Hint:* The dot product is defined as $lst1[0] \cdot lst2[0] + lst1[1] \cdot lst2[1] + \dots + lst1[n] \cdot lst2[n]$. The Python **zip** function may be useful here.
- `[[0], [0, 1], [0, 1, 2], [0, 1, 2, 3], [0, 1, 2, 3, 4]]`
- Return the same list as above, except now excluding every instance of the number 2: `[[0], [0, 1], [0, 1], [0, 1, 3], [0, 1, 3, 4]]`.

1.2 Draw the environment diagram that results from running the following code.

```
pom = [16, 15, 13]
pompom = pom * 2
pompom.append(pom[:])
pom.extend(pompom)
```

1.3 Draw the environment diagram that results from running the following code.

```
bless, up = 3, 5
another = [1, 2, 3, 4]
one = another[1:]

another[bless] = up
another.append(one.remove(2))
another[another[0]] = one
one[another[0]] = another[1]
one = one + [another.pop(3)]
another[1] = one[1][1][0]
one.append([one.pop(1)])
```

1.4 **def** jerry(jerry):
 def jerome(alex):
 alex.append(jerry[1:])
 return alex
 return jerome

```
ben = ['nice', ['ice']]
jerome = jerry(ben)
alex = jerome(['cream'])
ben[1].append(alex)
ben[1][1][1] = ben
print(ben)
```

- 1.5 Implement `subset_sum`, which takes in a list of integers and a number k and returns whether there is a subset of the list that adds up to k ? *Hint*: The `in` operator can determine if an element belongs to a list.

```
def subset_sum(seq, k):
    """
    >>> subset_sum([2, 4, 7, 3], 5)      # 2 + 3 = 5
    True
    >>> subset_sum([1, 9, 5, 7, 3], 2)
    False
    """
```

2 Trees

```
def tree(label, branches=[]):
    return [label] + list(branches)
```

```
def label(tree):
    return tree[0]
```

```
def branches(tree):
    return tree[1:]
```

- 2.1 A **min-heap** is a tree with the special property that every node's value is less than or equal to the values of all of its branches.



Implement `is_min_heap` which takes in a tree data abstraction and returns whether the tree satisfies the min-heap property or not.

3 Growth

3.1 Give a tight asymptotic runtime bound for the following functions in $\Theta(\cdot)$ notation, or “Infinite” if the program does not terminate.

(a) **def** one(n):

```

    while n > 0:
        n = n // 2

```

(b) **def** two(n):

```

    for i in range(n):
        for j in range(i):
            print(str(i), str(j))

```

(c) **def** three(n):

```

    i = 1
    while i <= n:
        for j in range(i):
            print(j)
        i *= 2

```

4 Nonlocals & OOP

4.1 Draw the environment diagram that results from running the code.

```

def campa(nile):
    def ding(ding):
        nonlocal nile
        def nile(ring):
            return ding
        return nile(ding(1914)) + nile(1917)

```

```
ring = campa(lambda nile: 103)
```

4.2 Implement the classes so that the code to the right runs.

```
class Plant:
```

```
    def __init__(self):
```

```
        def absorb(self):
```

```
        def grow(self):
```

```
class Leaf:
```

```
    def __init__(self, plant):
```

```
        def absorb(self):
```

```
        def __repr__(self):
            return 'Leaf'
```

```
class Sugar:
```

```
    def __init__(self, leaf, plant):
```

```
        def activate(self):
```

```
        def __repr__(self):
            return 'Sugar'
```

```
>>> p = Plant()
>>> p.height
1
>>> p.materials
[]
>>> p.absorb()
>>> p.materials
[Sugar]
>>> Sugar.sugars_created
1
>>> p.leaf.sugars_used
0
>>> p.grow()
>>> p.materials
[]
>>> p.height
2
>>> p.leaf.sugars_used
1
```

5 Exam Preparation *Extra Practice*

- 5.1 Implement `slice_reverse` which takes a linked list `s` and mutatively reverses the elements on the interval, $[i, j)$ (including i but excluding j). Assume `s` is zero-indexed, $i > 0$, $i < j$, and that `s` has at least j elements.

```
def slice_reverse(s, i, j):
    """
    >>> s = Link(1, Link(2, Link(3)))
    >>> slice_reverse(s, 1, 2)
    >>> s
    Link(1, Link(2, Link(3)))
    >>> s = Link(1, Link(2, Link(3, Link(4, Link(5)))))
    >>> slice_reverse(s, 2, 4)
    >>> s
    Link(1, Link(2, Link(4, Link(3, Link(5)))))
    """
    start = _____

    for _____:

        start = _____

    reverse = Link.empty

    current = _____

    for _____:

        _____

        current.rest = _____

        reverse = _____

        current = _____

    _____

    _____

    _____
```

- 5.2 A **Binary Search Tree** is a tree where each node contains either 0, 1, or 2 nodes and where the left branch (if present) contains values *strictly less than* ($<$) the root value, and the right branch (if present) contains values *strictly greater than* ($>$) the root value. The definition is recursive: both the left and right branches must also be BSTs for the entire tree to be a BST.

Implement `is_binary` which that takes in a `Tree t`, and returns `True` if `t` is a Binary Search Tree and `False` otherwise. Trees can contain any number of branches, but if a tree contains only one branch, interpret it as a left branch.

```
def is_binary(t):
    def binary(t, lo, hi):

        if _____:

            if t.is_leaf():
                return True

            elif _____:

                return _____

            elif _____:

                return _____

        return False
    return binary(t, float('-inf'), float('inf'))
```

- 5.3 Give a tight asymptotic runtime bound for the following scenarios in $\Theta(\cdot)$ notation, or “Infinite” if the program does not terminate. Assume the implementation of `is_binary` is optimal.

- (a) `is_binary` on a well-formed binary search tree with n nodes.
- (b) `is_binary` on a tree where each node contains 3 branches and the overall height of the tree is n .