Calculator

Announcements

## Exceptions

## Raise Statements

Python exceptions are raised with a raise statement
raise <expression>
<expression> must evaluate to a subclass of BaseException or an instance of one

Exceptions are constructed like any other object. E.g., TypeError('Bad argument!')
TypeError -- A function was passed the wrong number/type of argument
NameError -- A name wasn't found
KeyError -- A key wasn't found in a dictionary
RecursionError -- Too many recursive calls

## Try Statements

Try statements handle exceptions

```
try:
    <try suite>
except <exception class> as <name>:
    <except suite>
```

-••

## Execution rule:

The <try suite> is executed first

If, during the course of executing the <try suite>,
an exception is raised that is not handled otherwise, and

If the class of the exception inherits from <exception class>, then

The <except suite> is executed, with <name> bound to the exception

## Example: Reduce

## Reducing a Sequence to a Value

```
def reduce(f, s, initial):
    """Combine elements of s pairwise using f, starting with initial.
    E.g., reduce(mul, [2, 4, 8], 1) is equivalent to mul(mul(mul(1, 2), 4), 8).
    >>> reduce(mul, [2, 4, 8], 1)
    6 4
    |IIII
f is ...
    a two-argument function
s is ...
    a sequence of values that can be the second argument
initial is ...
    a value that can be the first argument
```



Programming Languages

## Programming Languages

A computer typically executes programs written in many different programming languages

Machine languages: statements are interpreted by the hardware itself

- A fixed set of instructions invoke operations implemented by the circuitry of the central processing unit (CPU)
- Operations refer to specific hardware memory addresses; no abstraction mechanisms

High-level languages: statements \& expressions are interpreted by another program or compiled (translated) into another language

- Provide means of abstraction such as naming, function definition, and objects
- Abstract away system details to be independent of hardware and operating system


## Python 3 <br> def square(x): <br> return x * x <br> ```from dis import dis \\ dis(square)```

## Python 3 Byte Code

```
LOAD_FAST
0 (x)
LOAD_FAST
0 (x)

\section*{Metalinguistic Abstraction}

A powerful form of abstraction is to define a new language! E.g.,
Problem domain: The MediaWiki mark-up language was designed for generating static web pages. It has built-in elements for text formatting and cross-page linking. It is used, for example, to create Wikipedia pages
\{\{Short description|Public university in Berkeley, California\}\}
\{\{Redirect-distinguish|Berkeley University|Berkeley College|Berkeley College (Yale University)\}\} \{\{Use American English|date=February 2019\}\}
\{\{Use mdy dates|date=November 2018\}\}
\{\{Infobox university
name \(\quad=\) University of California, Berkeley
image \(\quad=\) Seal of University of California, Berkeley.svg
motto \(=\) \{\{lang|la|[[Let there be light|Fiat lux]]\}\} ([[Latin]])
mottoeng = "Let there be light"
\}\}
A programming language has:
- Syntax: The legal statements and expressions in the language
- Semantics: The execution/evaluation rule for those statements and expressions

Parsing

\section*{Reading Scheme Lists}

A Scheme list is written as elements in parentheses:
```

(<element_0> <element_1> ... <element_n>)
A Scheme list

```

Each <element> can be a combination or primitive
(+ (* 3 (+ (* 24 ) (+ 3 5))) (+ (- 10 7) 6) )
The task of parsing a language involves coercing a string representation of an expression to the expression itself

\section*{Parsing}

A Parser takes text and returns an expression


\section*{Syntactic Analysis}

Syntactic analysis identifies the hierarchical structure of an expression, which may be nested

Each call to scheme_read consumes the input tokens for exactly one expression
```

'(', '+', 1, '(', '-', 23, ')', '(', '*', 4, 5.6, ')', ')'

```

Base case: symbols and numbers
Recursive call: scheme_read sub-expressions and combine them

\title{
Scheme-Syntax Calculator
}
(Demo)

\section*{Calculator Syntax}

The Calculator language has primitive expressions and call expressions. (That's it!)
A primitive expression is a number: \(2 \quad-4 \quad 5.6\)
A call expression is a combination that begins with an operator (+, -, *, /) followed by 0 or more expressions: (+ 123 ) (/ 3 (+ 4 5))

Expressions are represented as Scheme lists (Pair instances) that encode tree structures.


\section*{Calculator Semantics}

The value of a calculator expression is defined recursively.
Primitive: A number evaluates to itself.

Call: A call expression evaluates to its argument values combined by an operator.
+: Sum of the arguments
*: Product of the arguments
-: If one argument, negate it. If more than one, subtract the rest from the first.
/: If one argument, invert it. If more than one, divide the rest from the first.

Expression
(+ 5
(* 2 3)
(* 2 5 5) )

Expression Tree


\author{
Evaluation
}

\section*{The Eval Function}

The eval function computes the value of an expression, which is always a number
It is a generic function that dispatches on the type of the expression (primitive or call)

\section*{Implementation}
def calc_eval(exp):
    if isinstance(exp, (int, float)):
        return exp
    elif isinstance(exp, Pair):
        arguments \(=\) exp.rest.map(calc_eval)
        return calc_apply:(exp.first, arguments)
    else:
        raise TypeError \(\begin{aligned} & \text { '+', '-', } \\ & \text { ' } * \text { ', '/' }\end{aligned}\)

A Scheme list of numbers

\section*{Language Semantics}

A number evaluates...
to itself
A call expression evaluates...
to its argument values
combined by an operator

\section*{Applying Built-in Operators}

The apply function applies some operation to a (Scheme) list of argument values
In calculator, all operations are named by built-in operators: +, -, *, /

\section*{Implementation}
```

def calc_apply(operator, args):
if operator == '+':
return reduce(add, args, 0)
elif operator == '-':
elif operator == '*':
" !
elif operator == '/':
*"
else:
raise TypeError

```

Interactive Interpreters

\section*{Read-Eval-Print Loop}

The user interface for many programming languages is an interactive interpreter
1. Print a prompt
2. Read text input from the user
3. Parse the text input into an expression
4. Evaluate the expression
5. If any errors occur, report those errors, otherwise
6. Print the value of the expression and repeat

\section*{Raising Exceptions}

Exceptions are raised within lexical analysis, syntactic analysis, eval, and apply

Example exceptions
-Lexical analysis: The token 2.3.4 raises ValueError("invalid numeral")
-Syntactic analysis: An extra ) raises SyntaxError("unexpected token")
Eval: An empty combination raises TypeError("() is not a number or call expression")
Apply: No arguments to - raises TypeError("- requires at least 1 argument")

\section*{Handling Exceptions}

An interactive interpreter prints information about each error
A well-designed interactive interpreter should not halt completely on an error, so that the user has an opportunity to try again in the current environment```

