## CS61A Notes 04 - Lists (v1.0)

Pair Up!
Introducing - the only data structure you'll ever need in 61A - pairs.
A pair is a data structure that contains two things - the "things" can be atomic values or even another pair. For example, you can represent a point as ( x . y), and a date as (July . 1). Note the Scheme representation of a pair; the pair is enclosed in parentheses, and separated by a single period.

Note that there's an operator called pair? that tests whether something is a pair or not. For example, (pair? (cons 3 4)) is \#t, while (pair? 10) is \#f.

You've read about "cons", "car", and "cdr":
cons - takes in two parameters and constructs a pair out of them. So (cons 3 4) will return (3 . 4)
car - takes in a pair and returns the first part of the pair. So (car (cons 34)) will return 3.
cdr - takes in a pair and returns the second part of the pair. So (cdr (cons 34 4) ) will return 4.

These, believe it or not, will be all we'll ever need to build complex data structures in this course.
QUESTIONS: What do the following evaluate to?

```
(define u (cons 2 3)) (define w (cons 5 6)) (define x (cons u w))
(define y (cons w x)) (define z (cons 3 y))
```

1. $u$, $w, ~ x, y, z$ (write them out in Scheme's notation)
2. ( $\operatorname{car} \mathrm{y})$
3. (car (car y))
4. ( $\operatorname{cdr}(\operatorname{car}(\operatorname{cdr}(\operatorname{cdr} z))))$
5. ( $+(\operatorname{cdr}(\operatorname{car} y))(\operatorname{cdr}(\operatorname{car}(\operatorname{cdr} z))))$
6. (cons $z \mathrm{u}$ )
7. ( $\operatorname{cons}(\operatorname{car}(\operatorname{cdr} y))(\operatorname{cons}(\operatorname{car}(\operatorname{car} x))(\operatorname{car}(\operatorname{car}(\operatorname{cdr} z)))))$

## Then Came Lists

The super-cool definition of "list": a list is either an empty list, or a pair whose car is an element of the list and whose cdr is another list. Note the recursive definition - a list is a pair that contains a list! So then how does it end? Wouldn't there be an infinite number of list? Not so: an empty list, called "nil" and denoted ' () is a list containing no elements. And so it is, that every list ends with the empty list. To test whether a list is empty, you can use the null? operator on a list.

So, to make a list of elements $2,3,4$, we do this:

```
(define x (cons 2 (cons 3 (cons 4 '() ) ) ))
```

So x will be then represented as:

$$
(2 \cdot(3 .(4 . \quad()) \quad)
$$

Now, that looks a bit ugly, so Scheme, the nice friendly language that it is, sugar-coats the notation a bit so you get:

$$
\left(\begin{array}{lll}
2 & 3 & 4
\end{array}\right)
$$

It's a bit annoying to write so many cons to define x . So Scheme, the mushy-gushy language that it is, provides an operator list that takes in elements and returns them in a list. So we can also define x this way:

```
(define x (list 2 3 4))
```

Note: $(\operatorname{car} x)$ is 2 , $(\operatorname{cdr} x)$ is $(34)$, and $(\operatorname{car}(\operatorname{cdr} x))$ is 3 ! Well, it's a bit tiresome to write (car $(\operatorname{cdr} x))$ to get the second element of $x$. So Scheme, again the huggable lovable language that it is, provides a nifty short hand: ( $\operatorname{cadr} x$ ). This reads cader, and means "take the car of the cdr of". Similarly, you can use ( caddr $x$ ) - caderder - to take the car of the cdr of the cdr of $x$, which is 4. You can mix and match the ' $a$ ' and ' $d$ ' between the ' $c$ ' and ' $r$ ' to get the desired element of the list (up to a certain length).

You can also append two lists together. append takes in any number of lists and outputs a list containing those lists concatenated together. So (append (list 34) (list 5 6) ) returns (3 4 5 6).

## Don't You Mean sentence?

Oh stop grumbling. A "sentence" is actually a special kind of "list" - more specifically, a "sentence" is a flat list - a list without any sublists - whose elements can only be words or numbers. The operators of sentence - first, butfirst, se, etc. - are also much more forgiving in its domain than their list counterparts. For example, here's a list of equivalences:

```
first <==> car: (first '(1 2 3 4)) = (car '(1 1 2 3 4)) = 1
butfirst <==> cdr: (butfirst '(1 2 3 4 ) ) = (cdr '(11 2 3 4)) = (\begin{array}{lll}{2}&{3}&{4}\end{array})
empty? <==> null?: (empty? '()) = (null? '()) = #t
sentence <==> list: (se 1 2 3 4) = (list 1 2 3 4) = (1 2 3 4)
sentence <==> cons: (se 1 '(2 3 4)) = (cons 1 '(2 3 4)) = (1 2 3 4)
sentence <==> append: (se '(1 2) '(3 4)) = (append '(1 2) '(3 4)) = (1 2 3 4)
count <==> length: (count '(3 4 1)) = (length '(3 4 1)) = 3
every <==> map: (every square '(1 2)) = (map square '(1 2)) = (1 4)
keep <==> filter: (keep number? '(2 k 4)) = (filter number? '(2 k 4)) = (2 4)
```

Note that while se can be used for any combination of single-elements and sentences to make another sentence, list, cons and append are a bit more subtle, and what you pass in as arguments really matters.. For example,

```
(se '(1 2) '(3 4)) = (1 2 3 4) != (list '(1 2) '(3 4 4) )=((1 2) (3 4))
(se '(1 2) 3) = (1 2 3) != (cons '(1 2) 3) = ((\begin{array}{lll}{1}&{2}\end{array}). 3)
    != (list '(1 2) 3) = ((1 2) 3) != (append '(1 2) 3) = Error: 3 not a list!
```

And so on. You must be more careful with what you pass into the list operators! What do you get for all that trouble? Power - you can put anything into lists, not just words and numbers. You will now be able to construct deep lists - lists that contain sublists, which allows you to represent all sorts of cool things. You can also store exotic things like procedures in a list. The possibilities are endless!

## QUESTIONS:

1. Define a procedure list-4 that takes in 4 elements and outputs a list equivalent to one created by calling list.
2. Define a procedure length that takes in a list and returns the number of elements within the list.
3. Define a procedure list? that takes in something and returns \#t if it's a list, \#f otherwise.
4. Define append for two lists.
5. Suppose we have $x$ bound to a mysterious element. All we know is this:
(list? x) ==> \#t
(pair? $x$ ) ==> \#f
What is $x$ ?
6. Add in procedure calls to get the desired results. The blanks don't need to have anything:

7. Define a procedure (insert-after item mark ls) which inserts item after mark in ls.

## (Slightly) Harder Lists

1. Define a procedure (depth 1 s ) that calculates how maximum levels of sublists there are in 1 s . For example, (depth '(1) $\left.2 \begin{array}{lll}1 & 3 & 4\end{array}\right)$ ) $==>1$ (depth '(1 2 (3 4) 5)) ==> 2 (depth '(1 $2(345(67) 8) 9(1011) 12))==>3$
Remember that there's a procedure called max that takes in two numbers and returns the greater of the two.
2. Define a procedure (remove item ls) that takes in a list and returns a new list with item removed from ls.
3. Define a procedure (unique-elements 1 s ) that takes in a list and returns a new list without duplicates. You've already done this with remove-dups, and it used to do this:
(remove-dups '( $\left.\begin{array}{llllllll}3 & 5 & 6 & 3 & 3 & 5 & 9 & 8\end{array}\right)==\left(\begin{array}{lllll}6 & 3 & 5 & 9 & 8\end{array}\right)$
where the last occurrence of an element is kept. We'd like to keep the first occurrences:

Try doing it without using member?. You might want to use remove above.
4. Define a procedure (count-of item ls) that returns how many times a given item occurs in a given list; it could also be in a sublist. So, (count-of 'a '(abca a (b dac(a e) a) b (a))) ==> 7
5. Define a procedure (interleave ls1 ls2) that takes in two lists and returns one list with elements from both lists interleaved. So,

6. Write a procedure (apply-procs procs args) that takes in a list of single-argument procedures and a list of arguments. It then applies each procedure in procs to each element in args in order. It returns a list of results. For example,
(apply-procs (list square double +1) '(1 2234$)$ )
$==>\left(\begin{array}{llll}3 & 9 & 19 & 33\end{array}\right)$

## Expression Lists

Here's something interesting: a Scheme expression is really just a list! For example, ( +23 ) is just a list whose first element is the symbol + , second element the number 2, and third element the number 3. Similarly, (define x (* 45 )) is just a list whose first element is the symbol "define", second element the symbol " $x$ ", and third element the sublist ( $\begin{array}{lll} & 4 & 5\end{array}$ ). Given that insight, let's start manipulating more lists!

## QUESTIONS

1. Define a procedure (eval-plus exp) that takes in a valid Scheme expression consisting only of + and numbers, and evaluates it to the correct value. Assume that + always only gets two arguments. For example,
(eval-plus 3) ==> 3
(eval-plus '(+ 3 4)) ==> 7
(eval-plus '(+ 10 (+ 3 2)) ==> 15
2. (HARD!) Define (eval-plus exp) again, but let + take any number of arguments.
3. We'd like some easy way of creating a lambda expression. Write (make-lambda args body) that takes in the argument list and the body of a procedure, and produces the corresponding lambda expression. For example,
```
(make-lambda '(x y) '(+ x (* y x))) ==> (lambda (x y) (* y x))
```

4. Recall that there are two ways of defining procedures: the "real" way, and the sugar-coated way. Write a procedure (unsugar def) that takes in a procedure definition in sugar-coated syntax, and returns the same definition without using the syntactic sugar. For example,
```
(unsugar '(define (square x) (* x x)))
    ==> (define square (lambda (x) (* x x)))
```

5. Recall that a let expression is actually just a lambda expression. Write a procedure (let->lambda exp) that takes in a let expression and returns the corresponding lambda expression. For example,
(let->lambda '(let ((x 3) (y 10)) (+ x y)))
$==>$ ((lambda (x y) (+ x y)) 3 10)
