Pair Up!

QUESTIONS: What do the following evaluate to?

```
(define u (cons 2 3))
                       (define w (cons 5 6)) (define x (cons u w))
(define y (cons w x))
                       (define z (cons 3 y))
1. u, w, x, y, z (write them out in Scheme's notation)
   u: (2 . 3)
   w: (5 . 6)
   x: ((2 . 3) 5 . 6)
   y: ((5 . 6) (2 . 3) 5 . 6)
   z: (3 (5 . 6) (2 . 3) 5 . 6)
   Note: when you type this into STk, y and z look weird; I'm not sure why
   right now, but I'll look into it. As data structures, they still work
   fine.
2. (car y)
   (5.6)
3. (car (car y))
   5
4. (cdr (car (cdr (cdr z))))
   З
5. (+ (cdr (car y)) (cdr (car (cdr z))))
   12
6. (cons z u)
   ((3 (5 . 6) (2 . 3) 5 . 6) 2 . 3)
7. (cons (car (cdr y)) (cons (car (car x)) (car (cdr z)))))
   ((2.3) 2.5)
```

Then Came Lists

QUESTIONS:

1. Define a procedure list-4 that takes in 4 elements and outputs a list equivalent to one created by calling list.

```
(define (list-4 e1 e2 e3 e4)
  (cons e1 (cons e2 (cons e3 (cons e4 `())))))
```

2. Define a procedure length that takes in a list and returns the number of elements within the list. (define (length ls)

```
(if (null? ls)
0
(+ 1 (length (cdr ls)))))
```

3. Define a procedure list? that takes in something and returns #t if it's a list, #f otherwise.

Where's the base case?!

4. Define append for two lists.

5. Suppose we have x bound to a mysterious element. All we know is this: (list? x) ==> #t

```
(pair? x) ==> #f
What is x?
The only thing that's a list but not a pair is `(), the null list.
```

6. Add in procedure calls to get the desired results. The blanks don't need to have anything:

```
(cons 'a '(b c d e))
==> (a b c d e)
(append '(cs61a is) (list 'cool ))
==> (cs61a is cool)
(cons '(back to) '(save the universe))
==> ((back to) save the universe)
(cons '(I keep the wolf) (car '((from the door)) ) )
==> ((I keep the wolf) from the door)
7. Define a procedure (insert-after item mark 1s) which inserts item after mark in 1s.
(define (insert-after item mark 1s)
(cond ((null? 1s) '())
((equal? (car 1s) mark)
(cons (car 1s) (cons item (cdr 1s))))
(else (cons (car 1s) (insert-after item mark (cdr 1s))))))
```

(Slightly) Harder Lists

You probably need a while to convince yourself this is right. Why add 1 to the depth of car but not to the depth of cdr?

2. Define a procedure (remove item ls) that takes in a list and returns a new list with item removed from ls.

```
(define (remove item ls)
    (cond ((null? ls) `())
               ((equal? item (car ls)) (remove item (cdr ls)))
                (else (cons (car ls) (remove item (cdr ls))))))
```

- 3. Define a procedure (unique-elements 1s) that takes in a list and returns a new list without duplicates. You've already done this with remove-dups, and it used to do this: (remove-dups `(3 5 6 3 3 5 9 8)) ==> (6 3 5 9 8) where the *last* occurrence of an element is kept. We'd like to keep the *first* occurrences: (unique-elements `(3 5 6 3 3 5 9 8)) ==> (3 5 6 9 8) Try doing it without using member?. You might want to use remove above. (define (unique-elements 1s) (if (null? 1s) `() (cons (car 1s) (unique-elements (remove (car 1s) (cdr 1s)))))
- 4. Define a procedure (count-of item 1s) that returns how many times a given item occurs in a given list; it could also be in a sublist. So,

- 5. Define a procedure (interleave ls1 ls2) that takes in two lists and returns one list with elements from both lists interleaved. So, (interleave '(a b c d) (1 2 3 4 5 6 7)) ==> (a 1 b 2 c 3 d 4 5 6 7) (define (interleave ls1 ls2) (cond ((null? ls1) ls2) ((null? ls2) ls1) (else (cons (car ls1) (interleave ls2 (cdr ls1))))))
- 6. Write a procedure (apply-procs procs args) that takes in a list of single-argument procedures and a list of arguments. It then applies each procedure in procs to each element in args in order. It returns a list of results. For example, (apply-procs (list square double +1) '(1 2 3 4)) ==> (3 9 19 33) (define (apply-procs procs args) (if (null? procs) args (apply-procs (cdr procs) (map (car procs) args))))

Expression Lists

QUESTIONS

1. Define a procedure (eval-plus exp) that takes in a valid Scheme expression consisting only of + and numbers, and evaluates it to the correct value. Assume that + always only gets two arguments. For example,

```
(eval-plus 3) ==> 3
(eval-plus `(+ 3 4)) ==> 7
(eval-plus `(+ 10 (+ 3 2)) ==> 15
(define (eval-plus exp)
        (cond ((atom? exp) exp)
                    (else (+ (eval-plus (cadr exp)) (eval-plus (caddr exp))))))
```

2. (HARD!) Define (eval-plus exp) again, but let + take any number of arguments.

```
We're going to do "mutual recursion":
(define (eval-plus exp)
    (if (atom? exp)
        exp
        (add-expressions (cdr exp))))
(define (add-expressions exps)
    (if (null? exps)
        0
        (+ (eval-plus (car exps)) (add-expressions (cdr exps)))))
Note that eval-plus calls add-expressions to add up a list of
expressions, and add-expressions calls eval-plus to find out the value
of each expression it is given!
```

3. We'd like some easy way of creating a lambda expression. Write (make-lambda args body) that takes in the argument list and the body of a procedure, and produces the corresponding lambda expression. For example,

```
(make-lambda '(x y) '(+ x (* y x))) = > (lambda (x y) (* y x))
```

```
(define (make-lambda args body)
   (list 'lambda args body))
```

4. Recall that there are two ways of defining procedures: the "real" way, and the sugar-coated way. Write a procedure (unsugar def) that takes in a procedure definition in sugar-coated syntax, and returns the same definition without using the syntactic sugar. For example, (unsugar `(define (square x) (* x x)))

```
==> (define square (lambda (x) (* x x)))
```

```
Let's define a few helpers to help us:
(define (def-name def) (caadr def))
(define (def-args def) (cdadr def))
(define (def-body def) (caddr def))
(define (unsugar def)
```

```
(list 'define (def-name def)
      (make-lambda (def-args def) (def-body def))))
```

5. Recall that a let expression is actually just a lambda expression. Write a procedure (let->lambda exp) that takes in a let expression and returns the corresponding lambda expression. For example,

```
(let->lambda `(let ((x 3) (y 10)) (+ x y)))
==> ((lambda (x y) (+ x y)) 3 10)
```

Again, we'll use some helpers to make our code more readable: (define (let-vars exp) (map car (cadr exp))) (define (let-vals exp) (map cadr (cadr exp))) (define (let-body exp) (caddr exp)) (define (let->lambda exp) (cons (make-lambda (let-vars exp) (let-body exp)) (let-vals exp)))