## CS61A Notes 15 - Logic Is What Logic Declares Logic To Be [Solutions v1.0]

Lists Again (and again, and again, and again, and again... Someone changed the Matrix!)

## QUESTIONS

1. Write a rule for car of list. For example, ( $\operatorname{car}\left(\begin{array}{llll}1 & 2 & 3 & 4\end{array}\right)$ ?x) would have ?x bound to 1. (rule (car (?car . ?cdr) ?car))
2. Write a rule for $c d r$ of list. For example, ( $\operatorname{cdr}\left(\begin{array}{lll}1 & 2 & 3\end{array}\right)$ ? y$)$ would have ?y bound to (2 3). (rule (cdr (?car . ?cdr) ?cdr))
3. Using the above, write a query that would bind ? $x$ to the car of my-list. Write another query that would bind ?y to the cdr of my-list.
The temptation is to write (car my-list ?x) or (cdr my-list ?y). This doesn't work! There is no entry in the database whose first element is "car" and whose second element is the word "my-list". If you did that, you're thinking in the old Scheme way - that some "evaluator" will see my-list as a symbol and substitute in (1 234 ). This will not happen, since my-list isn't a variable! What you have to do is this:
(and (my-list ?ls) (car ?ls ?x))
First, we match ?ls to be (1 2 3 4), and then match ?x to be 1.
4. Define our old friend, member, so that (member 4 (1 $\left.2 \begin{array}{llll}1 & 3 & 4 & 5\end{array}\right)$ ) would be satisfied, and (member 3 (4 5 6)) would not, and (member 3 (1 2 (3 4) 5)) would not.
(rule (member ?item (?item . ?cdr)))
(rule (member ?item (?car . ?cdr)) (member ?item ?cdr))
5. Define its cousin, deep-member, so that (deep-member 3 (1 2 (3 4) 5) ) would be satisfied as well.
```
(rule (deep-member ?item (?item . ?cdr)))
(rule (deep-member ?item (?car . ?cdr)) (deep-member ?item ?car))
(rule (deep-member ?item (?car . ?cdr)) (deep-member ?item ?cdr))
Note how ?item can either be in ?car or ?cdr, so we need three rules.
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6. Define another old friend, reverse, so that (reverse (lllll $\left.\begin{array}{lll}1 & 2 & 3\end{array}\right)\left(\begin{array}{ll}3 & 2\end{array}\right)$ ) would be satisfied.
(rule (reverse () ()))
(rule (reverse (?car . ?cdr) ?reversed-ls)
(and (reverse ?cdr ?r-cdr)
(append ?r-cdr (?car) ?reversed-ls)))
7. (HARD!) Define its cousin, deep-reverse, so that
(deep-reverse (1 $2\left(\begin{array}{ll}3 & 4\end{array}\right)$ 5) (5 (4 3) 2 1) ) would be satisfied.
(rule (deep-reverse ?item ?item) (lisp-value atom? ?item))
(rule (deep-reverse () ()))
(rule (deep-reverse (?car . ?cdr) ?dr-ls)
(and (deep-reverse ?car ?r-car)
(deep-reverse ?cdr ?r-cdr)
(append ?r-cdr (?r-car) ?dr-ls)))
```
We need the first rule because recall that a "deep-list" could be an
atom, and that the third rule does not check if the ?car is an atom or
not when it recurses on it.
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8. Write the rule remove so that (remove $\begin{array}{lllllll}3 & \left(\begin{array}{lllll}1 & 2 & 3 & 4 & 3\end{array} \text { 2) ?what) binds ?what to (1 } 24\right. & 4\end{array}$ 2) - the list with 3 removed.
```
(rule (remove ?item () ()))
(rule (remove ?item (?item . ?cdr) ?result)
    (remove ?item ?cdr ?result))
(rule (remove ?item (?car . ?cdr) (?car . ?r-cdr))
    (and (not (same ?item ?car))
        (remove ?item ?cdr ?r-cdr)))
```

 what to (1 a 2 b 3 c d).
(rule (interleave ?ls () ?ls))
(rule (interleave () ?ls ?ls))
(rule (interleave (?car . ?cdr) ?ls2 (?car . ?r-cdr))
(interleave ?ls2 ?cdr ?r-cdr))
10. Consider this, not very interesting rule: ! (listify ? x (? x )) . So if we do (listify 3 ? what), ? what would be bound to (3).

Define a rule map with syntax (map procedure list result), so that (map listify (1) $\left.\begin{array}{l}1 \\ 2\end{array}\right)$ ((1) (2) (3))) would be satisfied, as would (map reverse ((1) $\left.\begin{array}{l}1\end{array}\right)\left(\begin{array}{ll}3 & 4 \\ 5\end{array}\right)$ ) ( (2 1) (5 4 3) )). In fact, we should be able to do something cool like (map ?what (1 2 3) ( (1)
(2) (3)) ) and have ?what bound to the word "listify". Assume the "procedure" we pass into map are of the form (procedure-name argument result).

```
(rule (map ?proc () ()))
(rule (map ?proc (?car . ?cdr) (?new-car . ?new-cdr))
    (and (?proc ?car ?new-car)
    (map ?proc ?cdr ?new-cdr)))
```

11. We can let predicates have the form (predicate-name argument true/false). Define a rule even so that (even 3) is not be satisfied, and (even 4) is.
(rule (even ?x) (lisp-value even? ?x))
12. The above is a way to make predicates. And once we have predicates, we can - and will, of course write a filter rule with the syntax (filter predicate list result) so that (filter even (1 $\left.\begin{array}{llllll}2 & 3 & 5 & 6\end{array}\right)\left(\begin{array}{lll}2 & 4 & 6\end{array}\right)$ ) would be satisfied, and querying (filter ?what (10 1112 13) (10 12)) would bind ?what to the word "even".
```
(rule (filter ?pred () ()))
(rule (filter ?pred (?car . ?cdr) (?car . ?new-cdr))
    (and (?pred ?car)
    (filter ?pred ?cdr ?new-cdr)))
(rule (filter ?pred (?car . ?cdr) ?new-ls)
    (and (not (?pred ?car))
        (filter ?pred ?cdr ?new-ls)))
```

