

CS61A Lecture 12

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(calc) review

(scheme-1) has
lambda **but NOT**
define



Remember calc-apply?

```
STk> (calc-apply '+ '(1 2 3))
6
STk> (calc-apply '* '(2 4 3))
24
STk> (calc-apply '/ '(10 2))
5
STk> (calc-apply '- '(9 2 3 1))
3
```



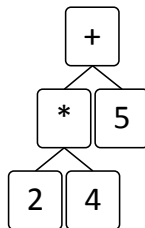
(calc) read-eval-print loop

```
(define (calc)
  (display "calc: ")
  (flush)
  (print (calc-eval (read)))
  (calc))
```



calc-eval

```
STk> (calc)
calc: (+ (* 2 4) 5)
13
(define (calc-eval exp)
  (cond
    ((number? exp) exp)
    ((list? exp)
     (calc-apply
      (car exp)
      (map calc-eval (cdr exp))))
    (else (error "Calc: bad exp"))))
```



(scheme-1)

DOES NOT HAVE DEFINE!

```
STk> (scheme-1)
Scheme-1: (lambda (x) (* x x))
(lambda (x) (* x x))
Scheme-1: ((lambda (x) (* x x)) 3)
9
```

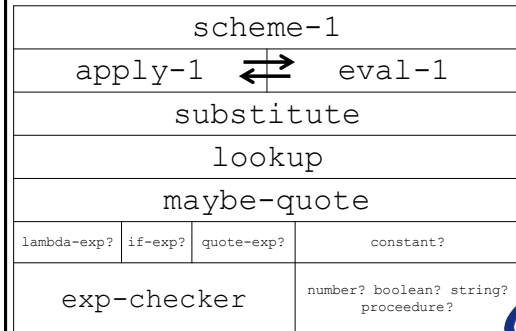


Working with large programs!

- Start with small functions
- Understanding code
 - Read it
 - Recursively figure out any functions it calls
 - Try to call the function to see what it does in different cases
 - Trace the function and try to call it from the functions that call it.

Cal

Approximate hierarchy of calls



Cal

lambda-exp?

```
STk> (lambda-exp? '(lambda (x) (+ x 2)))
#t
STk> (lambda-exp? '(+ 3 4))
#f
STk> (lambda-exp? '+)
#f
STk> (lambda-exp? '(lambda))
#t
```

It isn't as picky as we might hope...

Cal

Write lambda-exp? in terms of exp-checker

```
(define (exp-checker type)
  (lambda (exp)
    (and
      (pair? exp)
      (eq? (car exp) type))))

(lambda-exp? '(lambda (x) (+ x 2)))
```

Did you write it with syntactic sugar? A) Yes B)No

Cal

Some Helpers

```
(define quote-exp?
  (exp-checker 'quote))
(define if-exp?
  (exp-checker 'if))
(define (constant? exp)
  (or (number? exp)
      (boolean? exp)
      (string? exp)
      (procedure? exp)))
```

Cal

What is string?

```
STk> (string? "hello")
#t
STk> (string? 123)
#f
STk> (string? 'hello)
#f
```

Cal

(lookup name params args)

```
STk> (lookup 'x '(x) '(3))
3
STk> (lookup 'y '(x y) '(2 3))
3
STk> (lookup 'y '(x) '(3))
y
STk> (lookup '* '(x) '(3))
*
```

Just returns it if it
isn't in there

Cal

lookup full functionality

```
STk> (lookup 'fn
         '(x fn)
         '(3 (lambda (y) (* y y))))
(lambda (y) (* y y))
```

This already works

```
STk> (lookup 'x '(x) '(cat))
(quote cat)
```

cat was already a word, but
we want to tell other people this
thing IS ACTUALLY a word

Cal

Full lookup

```
(define (lookup name params args)
  (cond
    ((null? params) name)
    ((eq? name (car params))
     (maybe-quote (car args)))
    (else
     (lookup name
              (cdr params)
              (cdr args)))))
```

Cal

maybe-quote

```
(define (maybe-quote value)
  (cond
    ((lambda-exp? value) value)
    ((constant? value) value)
    ((procedure? value) value)
    (else (list 'quote value))))
```

Cal

Substitution using substitute**(substitute exp params args bound)**

```
STk> ((lambda (x) (* x x)) 3)
9
STk> (substitute '(* x x) '(x) '(3) '())
(* 3 3)
STk> ((lambda (x y) (+ x y)) 3 4)
7
STk> (substitute [ ] [ ] [ ] '())
(+ 3 4)
```

Cal

(substitute exp params args bound)

```
STk> ((lambda (x)
         (lambda (y) (* x y)) )
       3)
```

What does this return?

STk> (substitute

'()

Cal

Scheme substitution review

```
STk> ((lambda (x)
      (lambda (x)
        (* x x) )
      4)
#[closure arglist=(x) 7ff1a1f8]
STk> ((lambda (x)
      (lambda (x)
        (* x x))
      4)
      3)
```

What does this return? A) 9 B) 16 C)??



(substitute exp params args bound)

```
STk> ((lambda (x)
      (lambda (x)
        (* x x) )
      4)
#[closure arglist=(x) 7ff1a1f8]
STk> (substitute
      [redacted]
      [redacted]
      [redacted]
      [redacted])
```

[redacted]

[redacted]



(apply-1 proc args)

```
STk> (apply-1 + '(3 4))
7
```

Unlike calc-apply
apply-1 can be called
with REAL scheme functions

```
STk> (apply-1
      '(lambda (x) (* x x))
      '(3))
9
```

Or lists representing functions
Remember lambda-exp??



apply-1

```
(define (apply-1 proc args)
  (cond ((lambda (x) (* x x))
        ((procedure? proc)
         (apply proc args))
        ((lambda-exp? proc)
         (eval-1 (substitute
                  [redacted]
                  [redacted]
                  [redacted]
                  [redacted]))
         (else (error "bad proc:" proc))))
  (substitute exp params args bound))
```

scheme-1

```
(define (scheme-1)
  (display "Scheme-1: ")
  (flush)
  (print (eval-1 (read)))
  (scheme-1))
```



eval-1

```
STk> (eval-1 5)
5
(cond ((constant? exp) exp))

STk> (eval-1 '+)
#[closure arglist=args 7ff53de8]
(cond ((symbol? exp) (eval exp)))
```



eval-1

```
STk> (eval-1 '(if (> 3 4) 5 7))
7
```

```
(cond
  ((if-exp? exp)
   (if (eval-1 (cadr exp))
       (eval-1 (caddr exp))
       (eval-1 (caddr exp))))
```

Cal

eval-1

```
STk> (eval-1 'x)
*** Error:
      unbound variable: x
Current eval stack:
-----
0    x
1    (eval exp)
```

Things like + are quoted: '+
when they are passed to
eval-1 so this assumes x
will be a variable not a word.

eval-1 with words

```
STk> (eval-1 '(quote x))
x
```

```
STk> (eval-1 (quote (quote x)))
x (cond
  ((quote-exp? exp) (cadr exp)))
```

```
STk> (eval-1 ''x)
x
```

These are all equivalent!

Cal

eval-1

```
STk> (eval-1 '(lambda (x) (* x x)))
(lambda (x) (* x x))
```

```
(cond
  ((lambda-exp? exp) exp)
```

Cal

```
(define (eval-1 exp)
  (cond
    ((constant? exp) exp)
    ((symbol? exp) (eval exp))
    ((quote-exp? exp) (cadr exp))
    ((if-exp? exp)
     (if (eval-1 (cadr exp))
         (eval-1 (caddr exp))
         (eval-1 (caddr exp))))
    ((lambda-exp? exp) exp)
    ((pair? exp) _____)
    (else (error "?!?" exp))))
```

Cal

(cond ((pair? exp) _____))

```
STk> (eval-1 '(+ 2 3))
5
```

```
STk> (eval-1 '(+ (- 3 1) 5))
7
```

```
STk> (eval-1 '((lambda (x) (* x x)) 3))
9
```

Cal

(cond ((pair? exp) ____)

```
(cond
  ((pair? exp)
   (apply-1
    (eval-1 (car exp))
    (map eval-1 (cdr exp)))))
```

```
STk> (eval-1 '(+ 2 3))
STk> (eval-1 '(+ (- 3 1) 5))
STk> (eval-1 '((lambda (x) (* x x)) 3))
```

Cal

lambda-exp? Solution

```
(define (exp-checker type)
  (lambda (exp)
    (and
     (pair? exp)
     (eq? (car exp) type))))

(define (lambda-exp? exp)
  ((exp-checker 'lambda) exp))

(define lambda-exp?
  (exp-checker 'lambda))
```

**Write lookup
(some functionality missing)**

```
(define (lookup name params args)
  (cond
   ((null? params) name)
   ((eq? name (car params))
    (car args))
   (else
    (lookup name
             (cdr params)
             (cdr args)))))
```

Cal

Substitution using substitute**(substitute exp params args bound)**

```
STk> ((lambda (x) (* x x)) 3)
9
STk> (substitute '(* x x) '(x) '(3) '())
(* 3 3)
STk> ((lambda (x y) (+ x y)) 3 4)
7
STk> (substitute '(+ x y) '(x y) '(3 4) '())
(+ 3 4)
```

Cal

(substitute exp params args bound)

```
STk> ((lambda (x)
        (lambda (y) (* x y)) )
      3)

What does this return?
#[closure arglist=(y) 7ff1cc48]
A procedure that takes argument y and adds 3 to it
STk> (substitute
      '(lambda (y) (* x y))
      '(x)
      '(3)
      '())
(lambda (y) (* 3 y))
```

Cal

(substitute exp params args bound)

```
STk> ((lambda (x)
        (lambda (x)
          (* x x)) )
      4)

#[closure arglist=(x) 7ff1a1f8]
STk> (substitute
      '(lambda (x) (* x x))
      '(x)
      '(4)
      '(x))
(lambda (x) (* x x))
```

A recursive call
will be made
where bound
will be '(x)

Cal