## Lecture 28: Computer Security

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Many slides are adapted from CS 161 (Computer Security)

#### Announcements

- Final Exam on Friday (8/12) from 5-8pm in 155 Dwinelle
- Scheme Recursive Art submissions due today (8/9)!
- Potluck II tomorrow (8/10)! 5-8pm in Wozniak Lounge
- Homework 10 is due today (8/9)
  - AutoStyle EC portion due 8/10, last part due 8/11
- Homework 11 and 12 will be due 8/10 and 8/12
  - Last two of the three extra credit surveys
  - Vote for your favorite Recursive Art submissions!
- Check your grades! Details on Piazza, regrades close 8/10

#### Roadmap

Introduction

Functions

Data Mutability

Objects

- $\boldsymbol{\cdot}$  This week (Applications), the goals are:
  - To go beyond CS 61A and see examples of what comes next
  - To wrap up CS 61A!

Interpretation

Paradigms

Applications

#### Computer Security

# Computer Security

- A subfield of computer science with two main goals:
  - Allow intended use of computer systems
  - $\boldsymbol{\cdot}$  Prevent unwanted use that may cause harm
- · Why should you care?
  - ${\boldsymbol{\cdot}}$  The Internet has a lot of information about you...
- $\boldsymbol{\cdot}$  Today, we'll look at two problems:
  - Cryptography: secure communication over insecure communication channels
  - Injection Attacks

# Today's Special Guests! Alice Bob The Adversary (Eve or Mallory)

## Cryptography

#### Cryptography

- Three main goals: confidentiality, integrity, authenticity
- $\ensuremath{^{\bullet}}$  Today, we'll focus on confidentiality
- Confidentiality: prevent adversaries from reading private
  - Can Alice and Bob communicate in a way that even an eavesdropper Eve can't understand what they're saying?







#### The Caesar Cipher

(demo)

- One of the first attempts to encrypt a message
  - Was used by Roman dictator Julius Caesar
- ${}^{\bullet}$  Alice and Bob agree on a secret number (key) between 0 and 25 to shift the alphabet
  - For example, if the number is 2 then 'A' becomes 'C', 'B' becomes 'D', ..., 'Y' becomes 'A', 'Z' becomes 'B'

http://www.cryptoclub.org/tools/caesar\_cipher.ph

## Breaking the Caesar Cipher

(demo)

vgg ocz rjmgy'n v novbz ,
viy vgg ocz hzi viy rjhzi hzmzgt kgvtzmn :
oczt cvqz oczdm zsdon viy oczdm ziomvixzn ;
viy jiz hvi di cdn odhz kgvtn hvit kvmon ,

- Observation: There are only 26 possible keys
- $\bullet \ {\tt Observation:} \ {\tt Computers} \ {\tt are} \ {\tt fast}$
- ${\boldsymbol{\cdot}}$  Observation: Letters don't appear in English with the exact same frequency
  - ${\boldsymbol{\cdot}}$  For example, 'E' appears more often than 'Z'

# The Enigma Machine



- ${\boldsymbol{\cdot}}$  Used by the German military in World War II
- First broken by Polish mathematicians in 1932
- Information gained by the Allied forces is estimated to have shortened fighting by two years
- Implemented a progressive substitution cipher (e.g. different shift for each letter of the message)

# Better Cryptography

- This will require a bit of math, but the detailed steps aren't particularly important
- $\bullet$  From here onward, we'll represent a message with a number  $\mbox{\it m},$  rather than a string of characters
- Main idea: It is feasible to find three large numbers  $e, \\ d,$  and n such that  $(m^e)^d$  = m  $(mod\ n)$

#### The RSA Algorithm

- RSA is an example of public-key cryptography
  - The public key is known to everyone and is used to encrypt messages for the user
  - ${}^{\raisebox{3.5pt}{\text{\circle*{1.5}}}}$  The private key is only known by the user and is the only way to decrypt a message
  - This is also known as asymmetric cryptography: the message sender and recipient have two different keys
- Main idea: It is feasible to find three large numbers e, d, and n such that  $(m^e)^d = m \pmod n$
- Public key: **e** and **n** ("modulus")
- Private key: d

#### RSA Encryption and Decryption

- $\cdot$  Suppose that Bob wants to send a message  $\mathbf{m}$  to Alice
- He can encrypt a message by computing  $c = m^e \pmod{n}$ 
  - $\boldsymbol{\cdot}$  Everyone knows that Alice's public key is  $\boldsymbol{e}$  and  $\boldsymbol{n}$
- She can decrypt his message by computing  $c^d = (m^e)^d = m \pmod{n}$ 
  - Only Alice knows her private key  $\boldsymbol{d}$





#### Breaking RSA

- Eve needs to compute **d** to decrypt the message
- $\cdot$  e, d, and n aren't just three arbitrarily chosen numbers!
  - $\mathbf{n}$  =  $\mathbf{pq}$ , where  $\mathbf{p}$  and  $\mathbf{q}$  are two very large primes (~2<sup>1024</sup>)
  - For RSA encryption and decryption to work, ed = 1 (mod (p-1)\*(q-1)) (Euler's totient theorem)
- $\boldsymbol{\cdot}$  As far as we know, computing  $\boldsymbol{d}$  means that we have to
  - 1. Factor **n** into **p** and **q**
  - 2. Solve ed = 1 (mod (p-1)\*(q-1)) for d
- It turns out that Step 2 is easy and Step 1 is hard!
- $\boldsymbol{\cdot}$  The security of RSA relies on factoring being difficult

#### Factoring is (Maybe) Hard

- · Quick! Factor 561!
- There is no known efficient factoring algorithm
- Researchers spent 2007-2009 on factoring a 768-bit modulus (232 digits)
  - It took the equivalent of almost 2000 years of computing
  - Factoring a 1024-bit RSA modulus would be 1000x harder, but could happen in the next decade (2019 is coming up!)

# Factoring Complexity

- When people talk about factoring complexity, they typically describe runtime with respect to the bits that it takes to represent the number n (i.e.  $\log_2 n$ )
- Factoring is in NP: the answer can be verified by multiplying, which takes polynomial time
- We don't know if factoring is in P: the best algorithms for factoring are better than exponential but worse than polynomial
- ${\bf \cdot}$  Quantum computers can factor large numbers in polynomial time with Shor's algorithm
  - But their most recent breakthrough was factoring 21, so...

# Applications of RSA

- $\cdot$  For now (and for many years to come), RSA is secure
- ${\boldsymbol{\cdot}}$  Many protocols rely on RSA today
  - $\ensuremath{^{\bullet}}\xspace$  SSH (how to connect securely to the lab computers)
  - $\bullet$  SSL/TLS (the "S" in "HTTPS", how to connect securely to Facebook, etc.)

Break!	Injection Attacks

## Compromising Web Servers

- What could you do if you controlled one of Facebook's servers?
- $\boldsymbol{\cdot}$  Steal sensitive data (e.g. data from many users)
- Change server data (e.g. affect users)
- Gateway to enabling attacks on users
- $\boldsymbol{\cdot}$  Impersonation (of users to servers, or vice versa)

# Code Injection Attacks

(demo)

- $\boldsymbol{\cdot}$  Injection attacks are one way to compromise web servers
- People first started talking about this back in 1998, with hundreds of proposed fixes and solutions
- · General attack structure:
  - Attacker user provides some bad input
  - ${\boldsymbol{\cdot}}$  Web server does not check input format
  - Enables attacker to execute arbitrary code on the server

# ${\tt Summary}$

- Computer security studies how we can allow for the intended use of computer systems while preventing unwanted use that may cause harm
- ${\boldsymbol{\cdot}}$  Cryptography studies how we can communicate with others securely
- ${}^{\bullet}$  As programmers, we must be mindful of security best practices when developing applications
  - Even then, it might not be enough!
- $\boldsymbol{\cdot}$  CS 161 (Computer Security) goes into much more depth
- CS 261 and CS 276 are the graduate-level security and cryptography classes