

Integer Types and Literals

Signed?	Literals
Yes	Cast from int: (byte) 3
Yes	None. Cast from int: (short) 4096
No	'a' // (char) 97
	'\n' // newline ((char) 10)
	'\t' // tab ((char) 8)
	'\' ' // backslash
	'A', '\101', '\u0041' // == (char) 65
Yes	123
	0100 // Octal for 64
	0x3f, 0xffffffff // Hexadecimal 63, -1 (!)
Yes	123L, 01000L, 0x3fL 1234567891011L

Literals are just negated (positive) literals.

Means that there are 2^N integers in the domain of the type:

range of values is $-2^{N-1} .. 2^{N-1} - 1$.

And, only non-negative numbers, and range is $0..2^N - 1$.

5:44:05 2017

CS61B: Lecture #14 2

CS61B Lecture #14: Integers

5:44:05 2017

CS61B: Lecture #14 1

Modular Arithmetic: Examples

8) yields 0, since $512 - 0 = 2 \times 2^8$.

2) and (byte) (127+1) yield -128, since $128 - (-128) =$

*99) yields 15, since $9999 - 15 = 39 \times 2^8$.

*13) yields 122, since $-390 - 122 = -2 \times 2^8$.

yields $2^{16} - 1$, since $-1 - (2^{16} - 1) = -1 \times 2^{16}$.

5:44:05 2017

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Modular Arithmetic

How do we handle overflow, such as occurs in $10000 * 10000 * 10000$?

Some languages throw an exception (Ada), some give undefined re-

sult as the result of any arithmetic operation or conversion wraps to "wrap around"—*modular arithmetic*.

"next number" after the largest in an integer type is (like "clock arithmetic").

$128 == (\text{byte}) (127+1) == (\text{byte}) -128$

Result of some arithmetic subexpression is supposed to be in the range of T , an n -bit integer type,

compute the real (mathematical) value, x ,

and x' a number, x' , that is in the range of T , and that is congruent to x modulo 2^n .

Means that $x - x'$ is a multiple of 2^n .

5:44:05 2017

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Negative numbers

Representation for -1?

$$\begin{array}{r|l} 1 & 00000001_2 \\ + -1 & 11111111_2 \\ \hline = 0 & 100000000_2 \end{array}$$

In a byte, so bit 8 falls off, leaving 0.

Next bit is in the 2^8 place, so throwing it away gives an error modulo 2^8 . All bits to the left of it are also divisible

by 2^8 . Arithmetic is the same, but we choose to only use non-negative numbers modulo 2^{16} :

$$\begin{array}{r|l} 1 & 0000000000000001_2 \\ + 2^{16} - 1 & 1111111111111111_2 \\ = 2^{16} + 0 & 1000000000000000_2 \end{array}$$

5:44:05 2017

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Modular Arithmetic and Bits

How do we handle overflow?

Representation is the natural one for a machine that uses binary

Consider bytes (8 bits):

Decimal	Binary
101	1100101
$\times 99$	1100011
9999	100111 00001111
- 9984	100111 00000000
15	00001111

Bit n , counting from 0 at the right, corresponds to 2^n .

Bits to the left of the vertical bars therefore represent multiples of 256.

Throwing them away is the same as arithmetic modulo 256.

5:44:05 2017

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Promotion

operations (+, *, ...) promote operands as needed.

just implicit conversion.

operations,

operand is long, promote both to long.

promote both to int.

```
int aByte = 3;
long aLong = 3;
int i = (int) aByte + 3; // Type int
long l = aLong + (long) 3; // Type long
int j = (int) 'A' + 2; // Type int
char c = 'A' + 1; // ILLEGAL (why?)
```

ely,

```
1; // Defined as aByte = (byte) (aByte+1)
```

mples:

```
char aChar = 'A';
char lowerCaseChar = (char) ('a' + aChar - 'A'); // why cast?
```

5:44:05 2017

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Bit twiddling

C++ allow for handling integer types as sequences of "version to bits" needed: they already are.

and their uses:

Set	Flip	Flip all
00101100	00101100	
10100111	~ 10100111	~ 10100111
10101111	10001011	01011000

	Arithmetic Right	Logical Right
1 << 3	10101101 >> 3	10101100 >>> 3
0	11110101	00010101

```
1) >>> 29? = 7.
<< n?
>> n?
>>> 3) & ((1<<5)-1)?
```

5:44:05 2017

CS61B: Lecture #14 10

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1) >>> 29? = 7.
<< n? = x * 2^n.
>> n? = [x/2^n] (i.e., rounded down).
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Conversion

Java will silently convert from one type to another if this and no information is lost from value.

cast explicitly, as in (byte) x.

```
byte aByte; char aChar; short aShort; int anInt; long aLong;
```

```
aByte; anInt = aByte; anInt = aShort;
aChar; aLong = anInt;
```

(, might lose information:

```
long; aByte = anInt; aChar = anInt; aShort = anInt;
aChar; aChar = aShort; aChar = aByte;
```

special dispensation:

```
3; // 13 is compile-time constant
2+100 // 112 is compile-time constant
```

5:44:05 2017

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1) >>> 29? $= 7.$
 << n? $= x \cdot 2^n.$
 >> n? $= \lfloor x/2^n \rfloor$ (i.e., rounded down).
 >>> 3) & ((1<<5)-1)? **5-bit integer, bits 3-7 of x.**