#### CS61B Lecture #31

ed search structures (DS(IJ), Chapter 9

bm Numbers (DS(IJ), Chapter 11)

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#### The Trie: Example

abash, abate, abbas, axolotl, axe, fabric, facet} show paths followed for "abash" and "fabric" node corresponds to a possible prefix. n path to node = that prefix.



#### A Side-Trip: Scrunching

bvious implementation for internal nodes is array inaracter.

erformance, L length of search key. independent of N, number of keys. Is there a depen-

ays are sparsely populated by non-null values—waste of

arrays on top of each other!

mpty) entries of one array to hold non-null elements of

arkers to tell which entries belong to which array.

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#### Public Service Announcement

erimental Social Science Laboratory (Xlab) invites cipate in social science studies! Experiments conab (located in Hearst Gym, Suite 2) are computern-making studies such as tasks, surveys, and games. sionally offer remote online and mobile studies that leted anywhere. Participants earn \$15/hour on avtime they participate. For more information, visit y.edu. To sign up, visit berkeley.sona-systems.com."

## Ily Efficient Use of Keys: the Trie

lent about cost of comparisons. worst case is length of string. hould throw extra factor of key length, L, into costs: parisons really means  $\Theta(ML)$  operations. for key X, keep looking at same chars of X M times. tter? Can we get search cost to be O(L)?

multi-way decision tree, with one decision per character

#### Adding Item to a Trie

ding bat and faceplate.

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faceplate□ facet□

icked.

#### robabilistic Balancing: Skip Lists

n be thought of as a kind of n-ary search tree in which put the keys at "random" heights.

hought of as an ordered list in which one can skip large



# 

tart at top layer on left, search until next step would nen go down one layer and repeat.

, we search for 125 and 127. Gray nodes are looked at; nodes are overshoots.

he nodes were chosen randomly so that there are about nodes that are > k high as there are that are k high.

hes fast with high probability.

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#### Scrunching Example

(unrelated to Tries on preceding slides)

rrays, each indexed 0..9



them, but keep track of original index of each item:



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Summary	tructures that Implement	Abstractions		
arch trees allows us to realize $\Theta(\lg N)$ performance. -black trees: N) performance for searches, insertions, deletions. ood for external storage. Large nodes minimize # of ations performance for searches, insertions, and deletions, s length of key being processed. to manage space efficiently. idea: scrunched arrays share space. able $\Theta(\lg N)$ performace for searches, insertions, dele- nplement. for interesting ideas: probabilistic balance, random- structures. 8:19 2017 C5618: Lecture #31 20	linked lists, circular buffers et Queue: heaps Set: binary search trees, rea arrays or linked lists d Set: hash table Nap: hash table p: red-black trees, B-trees, sort	d-black trees, B-trees, ed arrays or linked lists 22		
Example: Adding and deleting m initial list: r, we add 126 and 127 (choosing random heights for emove 20 and 40: the second	Multiset Contains, iterator Priority Queue Contains, iterator Priority Queue Contains, iterator Priority Queue Contains, iterator get Ordered Set first Contend Set Sorted Set Sorted Set Contains, iterator get Ordered	rractions ava has corresponding interface Java has no corresponding interface	Corresponding Classes iction) ist, LinkedList, Stack, Array iet Queue: PriorityQueue Set (SortedSet): TreeSet d Set: HashSet Nap: HashMap p (SortedMap): TreeMap	<b>s in Java</b> yBlockingQueue,
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