Public-Service Announcement I

an engineer, designer, or entrepreneur? Come check
evelopers of Berkeley. We create end-to-end solu-
diverse breath of skills. Our members hone their
tend, backend, design, and engineering while devel-
erstanding of topics such as software architecture,
perating systems, and networking. We build apps
wide range of industries while leveraging the latest
in IOT, AI, fintech, ML, AR, and more. Join us for
e tonight, Wednesday 9/6. Our members will show
and what it took to make them. Free food, fidget
wag will also be there.
s://www.mobiledevsberkeley.org/ for further infor-

Public-Service Announcement II

Political Computer Science @ Berkeley! PCS har-
nts' intellectual capabilities and potential in CS and
affairs to address current issues within the United
and political systems.
dy have an Amazon Alexa Skill called Political Pun-
bout to be published; it is able to give responses to
ut politics/government.
he many other projects that we have in the works
o gerrymandering, campaign finance, Alexa, and more),
site: pcsberkeley.wixsite.com/pcsberkeley
w accepting applications for the 2017-2018 school
takes 5 minutes to complete)! Apply on our website!"

Recreation

n of the coefficients of

$$(1 - 3x + 3x^2)^{743}(1 + 3x - 3x^2)^{744}$$

and collecting terms?

CS61B Lecture #6: Arrays

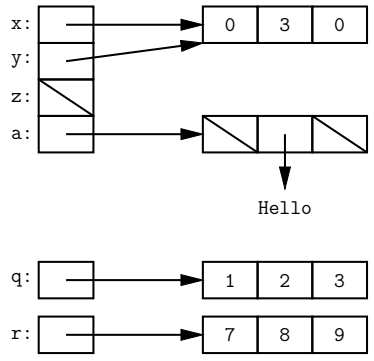
structured container whose components are
fixed integer.
e of **length** simple containers of the same type, num-
m 0.
eld usually implicit in diagrams.)
nonymous, like other structured containers.
rred to with pointers.
inted to by A,
A.length
d component i is A[i] (i is the *index*)
t feature: index can be *any integer expression*.

A Few Samples

Java

```
;
[3];
';
{ 1, 2, 3 };
for declarations:
 8, 9 };
```

Results



Example: Accumulate Values

up the elements of array A.

```
(int[] A) {

= 0; i < A.length; i += 1)
```

```
// New (1.5) syntax
for (int x : A)
   N += x;
```

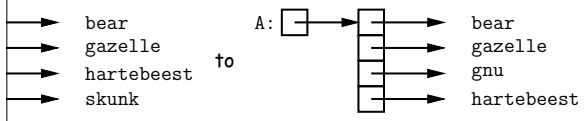
```
d-core: could have written

 i<A.length; N += A[i], i += 1)
just ;

don't: it's obscure.
```


Example: Insert into an Array

t a call like `insert(A, 2, "gnu")` to convert (destruc-

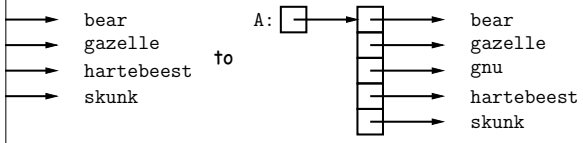



```
location K in ARR, moving items K, K+1, ... to locations
..  The last item in ARR is lost. */
rt (String[] arr, int k, String x) {
rr.length-1; i > k; i -= 1) // Why backwards?
[i-1];
 to this loop:
rraycopy(arr, k, arr, k+1, arr.length-k-1);*/
```

from *to* *# to copy*

Growing an Array

ose that we want to change the description above, so
t2 (A, 2, "gnu") does *not* shove "skunk" off the end,
ws" the array.




```
, r, where r.length = ARR.length+1; r[0..K-1]
ARR[0..K-1], r[k] = x, r[K+1..] same as ARR[K..]. */
insert2(String[] arr, int k, String x) {
t = new String[arr.length + 1];
 0, result, 0, k);
 k, result, k+1, arr.length-k);
```

d a different return type from `insert2`??

(Aside) Java Shortcut

Can write just '`arraycopy`' by including at the top of the

```
ic java.lang.System.arraycopy;
```

define the simple name `arraycopy` to be the equivalent
g.System.arraycopy in the current source file."

ame for `out` so that you can write

```
(...);
```

```
println(...);
```

laration like

```
ic java.lang.Math.*;
```

all the (public) static definitions in `java.lang.Math` and
vailable in this source file by their simple names (the
he last dot)."

unctions like `sin`, `sqrt`, etc.

Example: Merging Program

n two sorted arrays of ints, A and B, produce their
l array containing all from A and B.
der to solve this recursively, it is useful to *generalize*
ction to allow merging *portions* of the arrays.

```
nd B are sorted, returns their merge. */
t[] merge(int[] A, int[] B) {
A, 0, B, 0);

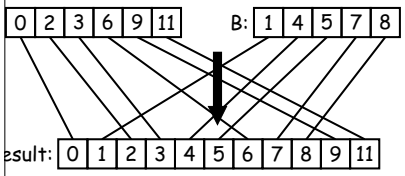

 A[L0..] and B[L1..] assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1) {
gth - L0 + B.length - L1; int[] C = new int[N];
ength) arraycopy(B, L1, C, 0, N);
= B.length) arraycopy(A, L0, C, 0, N);
] <= B[L1]) {
0]; arraycopy(merge(A, L0+1, B, L1), 0, C, 1, N-1);

1]; arraycopy(merge(A, L0, B, L1+1), 0, C, 1, N-1);
```

What is wrong with this implementation?

Example: Merging

n two sorted arrays of ints, A and B, produce their
l array containing all items from A and B.



A Tail-Recursive Solution

```
t[] merge(int[] A, int[] B) {
A, 0, B, 0, new int[A.length+B.length], 0);


] and B[L1..] into C[K..], assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1, int[] C, int k){
ength) /* ? */
= B.length) /* ? */
] <= B[L1]) {
0];

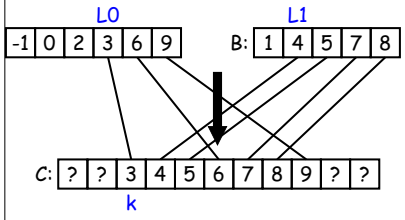

1];
```


A Tail-Recursive Strategy

```
t[] merge(int[] A, int[] B) {
A, 0, B, 0, new int[A.length+B.length], 0);


] and B[L1..] into C[K..], assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1, int[] C, int k){
```

d merges *part* of `A` with part of `B` into part of `C`. For
er a possible call `merge(A, 3, B, 1, C, 2)`



A Tail-Recursive Solution

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] and B[L1..] into C[K..], assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1, int[] C, int k){
ength) arraycopy(B, L1, C, k, B.length-L1);
= B.length) arraycopy(A, L0, C, k, A.length-L0);
] <= B[L1]) {
0];


1];
```


A Tail-Recursive Solution

```
t[] merge(int[] A, int[] B) {
A, 0, B, 0, new int[A.length+B.length], 0);


] and B[L1..] into C[K..], assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1, int[] C, int k){
ength) /* ? */
= B.length) /* ? */
] <= B[L1]) {
0];


1];
```


Iterative Solution

lon't use either of the previous approaches in languages
Array manipulation is most often iterative:

```
 int[] merge(int[] A, int[] B) {
ew int[A.length + B.length];
```


A Tail-Recursive Solution

```
t[] merge(int[] A, int[] B) {
A, 0, B, 0, new int[A.length+B.length], 0);


] and B[L1..] into C[K..], assuming A and B sorted. */
ge(int[] A, int L0, int[] B, int L1, int[] C, int k){
ength) arraycopy(B, L1, C, k, B.length-L1);
= B.length) arraycopy(A, L0, C, k, A.length-L0);
] <= B[L1]) {
0];
0+1, B, L1, C, k+1);

1];
), B, L1+1, C, k+1);
```


Alternative Solution: Removing k

ant of the loop is that k=L0+L1.

```
t[] merge(int[] A, int[] B) {
 int[A.length + B.length];


1 < C.length) {
 A.length) {
+ L1] = B[L1]; L1 += 1;
 (L1 >= B.length) {
+ L1] = A[L0]; L0 += 1;
 (A[L0] <= B[L1]) {
+ L1] = A[L0]; L0 += 1;

+ L1] = B[L1]; L1 += 1;
```


Iterative Solution II

```
t[] merge(int[] A, int[] B) {
 int[A.length + B.length];


0; k < C.length; k += 1) {
 A.length) {
= B[L1]; L1 += 1;
 (L1 >= B.length) {
= A[L0]; L0 += 1;
 (A[L0] <= B[L1]) {
= A[L0]; L0 += 1;

= B[L1]; L1 += 1;
```

Multidimensional Arrays

- or higher-dimensional layouts, such as

$$A = \begin{array}{|c|c|c|c|} \hline 2 & 3 & 4 & 5 \\ \hline 4 & 9 & 16 & 25 \\ \hline 8 & 27 & 64 & 125 \\ \hline \end{array}$$?

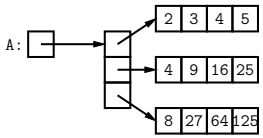
Multidimensional Arrays in Java

primitive in Java, but we can build them as arrays of

```
ew int[3][];
t[] {2, 3, 4, 5};
t[] {4, 9, 16, 25};
t[] {8, 27, 64, 125};
```

```
[] { {2, 3, 4, 5},
     {4, 9, 16, 25},
     { 8, 27, 64, 125} };

 {2, 3, 4, 5},
 {4, 9, 16, 25},
 {8, 27, 64, 125} };
```

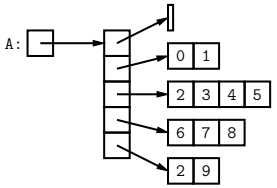



```
ew A[3][4];
0; i < 3; i += 1)
j = 0; j < 4; j += 1)
j] = (int) Math.pow(j + 2, i + 1);
```


Exotic Multidimensional Arrays

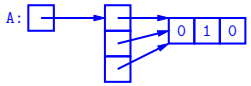
element of an array is independent, there is no single neral:

```
= new int[5][];
 int[] {};
 int[] {0, 1};
 int[] {2, 3, 4, 5};
 int[] {6, 7, 8};
 int[] {9};
```



his print?

```
RO = new int[3][];
ZERO[1] = ZERO[2] =
t[] {0, 0, 0};
 = 1;
.println(ZERO[2][1]);
```



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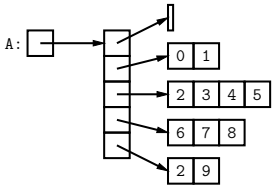
```



his print?

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ZERO[1] = ZERO[2] =
t[] {0, 0, 0};
 = 1;
.println(ZERO[2][1]);
```