Disclaimer: This discussion worksheet is fairly long and is not designed to be finished in a single section. Some of these questions of the level that you might see on an exam and are meant to provide extra practice with asymptotic analysis.

1 More Running Time

Give the worst case and best case running time in $\Theta(\cdot)$ notation in terms of M and N.

(a) Assume that slam() $\in \Theta(1)$ and returns a boolean.

```
public void comeon(int M, int N) {
       int j = 0;
       for (int i = 0; i < N; i += 1) {</pre>
3
           for (; j < M; j += 1) {
                if (slam(i, j))
5
                    break;
7
           }
       }
       for (int k = 0; k < 1000 * N; k += 1) {
10
           System.out.println("space jam");
11
13 }
```

(b) *Exam Practice:* Give the worst case and best case running time in $\Theta(\cdot)$ notation in terms of N for find.

```
public static boolean find(int tgt, int[] arr) {
       int N = arr.length;
       return find(tgt, arr, 0, N);
4 }
 private static boolean find(int tgt, int[] arr, int lo, int hi) {
       if (lo == hi || lo + 1 == hi) {
           return arr[lo] == tgt;
8
       int mid = (lo + hi) / 2;
       for (int i = 0; i < mid; i += 1) {</pre>
10
           System.out.println(arr[i]);
11
12
       return arr[mid] == tgt || find(tgt, arr, lo, mid)
13
                               || find(tgt, arr, mid, hi);
14
15 }
```

2 Recursive Running Time

For the following recursive functions, give the worst case and best case running time in $\Theta(\cdot)$ notation.

(a) Give the running time in terms of N.

(b) Give the running time for andwelcome (arr, 0, N) where N is the length of the input array arr.

```
public static void andwelcome(int[] arr, int low, int high) {
       System.out.print("[ ");
       for (int i = low; i < high; i += 1) {</pre>
            System.out.print("loyal ");
6
       System.out.println("]");
7
       if (high - low > 0) {
           double coin = Math.random();
8
           if (coin > 0.5) {
9
               andwelcome(arr, low, low + (high - low) / 2);
10
11
           } else {
               andwelcome(arr, low, low + (high - low) / 2);
12
               andwelcome(arr, low + (high - low) / 2, high);
           }
14
15
16 }
```

(c) Give the running time in terms of N.

```
public int tothe(int N) {
    if (N <= 1) {
      return N;
    }
    return tothe(N - 1) + tothe(N - 1) + tothe(N - 1);
}</pre>
```

(d) Exam Practice: Give the running time in terms of N

```
public static void spacejam(int N) {
    if (N == 1) {
        return;
}

for (int i = 0; i < N; i += 1) {
        spacejam(N-1);
}

}</pre>
```

3 Hey you watchu gon do?

For each example below, there are two algorithms solving the same problem. Given the asymptotic runtimes for each, is one of the algorithms **guaranteed** to be faster? If so, which? And if neither is always faster, explain why. Assume the algorithms have very large input (so N is very large).

- (a) Algorithm 1: $\Theta(N)$, Algorithm 2: $\Theta(N^2)$
- (b) Algorithm 1: $\Omega(N)$, Algorithm 2: $\Omega(N^2)$
- (c) Algorithm 1: O(N), Algorithm 2: $O(N^2)$
- (d) Algorithm 1: $\Theta(N^2)$, Algorithm 2: $O(\log N)$
- (e) Algorithm 1: $O(N \log N)$, Algorithm 2: $\Omega(N \log N)$

Why did we need to assume that N was large?

4 Big Ballin' Bounds

- 1. Prove the following bounds by finding some constant M > 0 and input N > 0 for $M \in \mathbb{R}, N \in \mathbb{N}$ such that f and g satisfy the relationship.
 - (a) $f \in O(g)$ for f = 2n, $g = n^2$
 - (b) $f \in \Omega(g)$ for f = 0.1n, g = 40
 - (c) $f \in \Theta(g)$ for $f = \log(n)$, $g = \log(n^a)$, for a > 0.
- 2. Answer the following claims with true or false. If false, provide a counterexample.
 - (a) If $f(n) \in O(g(n))$, then $500f(n) \in O(g(n))$.
 - (b) If $f(n) \in \Theta(g(n))$, then $2^{f(n)} \in \Theta(2^{g(n)})$.