## Integer Types and Literals

| Signed? | Literals |
| :---: | :---: |
| $\begin{aligned} & \text { Yes } \\ & \text { Yes } \end{aligned}$ | Cast from int: (byte) 3 <br> None. Cast from int: (short) 4096 |
| No | 'a' // (char) 97 ' $\backslash n$ ' // newline ((char) 10) ' lt ' // tab ((char) 8) ' $\$ ', // backslash 'A', '\101', ' $\backslash u 0041$ ' // == (char) |
| Yes | ```123 0100 // Octal for 64 0x3f, Oxfffffffff // Hexadecimal 63,``` |
| Yes | $\begin{aligned} & \text { 123L, 01000L, Ox3fL } \\ & \text { 1234567891011L } \end{aligned}$ |

nerals are just negated (positive) literals.
ns that there are $2^{N}$ integers in the domain of the type:
range of values is $-2^{N-1}$.. $2^{N-1}-1$.
ed, only non-negative numbers, and range is $0 . .2^{N}-1$.
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## CS61B Lecture \#14: Integers

## Modular Arithmetic: Examples

8) yields 0 , since $512-0=2 \times 2^{8}$.
9) and (byte) $(127+1)$ yield -128 , since $128-(-128)=$
*99) yields 15 , since $9999-15=39 \times \cdot 2^{8}$.
*13) yields 122 , since $-390-122=-2 \times 2^{8}$.
yields $2^{16}-1$, since $-1-\left(2^{16}-1\right)=-1 \times 2^{16}$.

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## Modular Arithmetic

$w$ do we handle overflow, such as occurs in $10000 * 10000 * 10000$ ? ges throw an exception (Ada), some give undefined re)
; the result of any arithmetic operation or conversion pes to "wrap around"-modular arithmetic.
"next number" after the largest in an integer type is (like "clock arithmetic").

128 == (byte) ( $127+1$ ) == (byte) -128
sult of some arithmetic subexpression is supposed to $T$, an $n$-bit integer type,
ompute the real (mathematical) value, $x$,
a number, $x^{\prime}$, that is in the range of $T$, and that is to $x$ modulo $2^{n}$.
ans that $x-x^{\prime}$ is a multiple of $2^{n}$.)

## Negative numbers

resentation for -1?

$$
\left.\begin{array}{r|r}
1 & 00000001_{2} \\
+\quad-1 & 11111111_{2} \\
= & 0
\end{array} \right\rvert\, \begin{array}{r|}
\hline 00000000_{2}
\end{array}
$$

$n$ a byte, so bit 8 falls off, leaving 0.
ed bit is in the $2^{8}$ place, so throwing it away gives an modulo $2^{8}$. All bits to the left of it are also divisible
types (char), arithmetic is the same, but we choose to ly non-negative numbers modulo $2^{16}$.

$$
\begin{array}{r|r}
1 & 0000000000000001_{2} \\
+2^{16}-1 & 1111111111111111_{2} \\
=2^{16}+0 & 1 \mid 0000000000000000_{2}
\end{array}
$$

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## Modular Arithmetic and Bits

ound?
tion is the natural one for a machine that uses binary
consider bytes (8 bits):

| Decimal | Binary |
| ---: | ---: |
| 101 | 1100101 |
| $\times 99$ | 1100011 |
| 9999 | $100111 \mid 00001111$ |
| -9984 | $100111 \mid 00000000$ |
| 15 | 00001111 |

it $n$, counting from 0 at the right, corresponds to $2^{n}$. he left of the vertical bars therefore represent multi256.
them away is the same as arithmetic modulo 256
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## Promotion

perations (+, *, ...) promote operands as needed.
just implicit conversion.
pperations,

```
rand is long, promote both to long.
promote both to int.
\beta== (int) aByte + 3 // Type int
\beta== aLong + (long) 3 // Type long
= (int) 'A' + 2 // Type int
Byte + 1 // ILLEGAL (why?)
ely,
1; // Defined as aByte = (byte) (aByte+1)
mple:
& aChar is an upper-case letter
prCaseChar = (char) ('a' + aChar - 'A'); // why cast?
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```


## Conversion

ava will silently convert from one type to another if this and no information is lost from value.
ast explicitly, as in (byte) x.
e; char aChar; short aShort; int anInt; long aLong;
aByte; anInt = aByte; anInt = aShort;
Char; aLong = anInt;
;, might lose information.
Long; aByte = anInt; aChar = anInt; aShort = anInt;
aChar; aChar = aShort; aChar = aByte;
special dispensation:
3; // 13 is compile-time constant
2+100// 112 is compile-time constant

## Bit twiddling

C++) allow for handling integer types as sequences of iversion to bits" needed: they already are.
Ind their uses:

| Set | Flip | Flip all |
| ---: | ---: | ---: | :--- |
| 00101100 | 00101100 |  |
| 10100111 | -10100111 | $\sim 10100111$ |
| 10101111 | 10001011 | 01011000 |

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| 00101100 | 00101100 |  |  |
| 10100111 | $\sim$ | 10100111 | $\sim 10100111$ |
| 10101111 | 10001011 | 01011000 |  |


|  | Arithmetic Right | Logical Right |
| :---: | :---: | :---: |
| $1 \ll 3$ | 10101101 >> 3 | 10101100 >>> 3 |
|  | 11110101 | 00010101 |
| 1) >>> 29? |  |  |
| << $n$ ? |  |  |
| >> $n$ ? |  |  |
| >>> 3) | \& ( $(1 \ll 5)-1) ?$ |  |

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| Set | Flip | Flip all |
| :---: | :---: | :---: |
| 00101100 | 00101100 |  |
| \| 10100111 | 10100111 | 10100111 |
| 10101111 | 10001011 | 01011000 |


|  | Arithmetic Right | Logical Right |
| :---: | :---: | :---: |
| $1 \ll 3$ | 10101101 >> 3 | 10101100 >>> 3 |
| 0 | 11110101 | 00010101 |
| 1) >>> 29? |  |  |
| << $n$ ? |  |  |
| > $>n ?$>> |  | 」 (i.e., rounded down). |
|  | \& $((1 \ll 5)-1) ?$ |  |

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## Bit twiddling

C++) allow for handling integer types as sequences of version to bits" needed: they already are

## ind their uses:

| Set | Flip | Flip all |
| ---: | :---: | :---: | :---: |
| 00101100 | 00101100 |  |
| 10100111 | -10100111 | $\sim 10100111$ |
| 10101111 | 10001011 | 01011000 |


|  | Arithmetic Right | Logical Right |
| :---: | :---: | :---: |
| $1 \ll 3$ | 10101101 >> | 10101100 >>> 3 |
| 0 | 11110101 | 00010101 |
| 1) >>> 29? $=$ |  | $=7$. |
| << $n$ ? |  | $=x \cdot 2^{n}$. |
| >> $n$ ? |  |  |
| >>> 3) | \& $((1 \ll 5)-1) ?$ |  |

## Bit twiddling

C++) allow for handling integer types as sequences of iversion to bits" needed: they already are.

| ind their uses: |  |  |
| :---: | :---: | :---: |
| Set | Flip | Flip all |
| 00101100 | 00101100 |  |
| \| 10100111 | - 10100111 | 10100111 |
| 10101111 | 10001011 | 01011000 |
| Arithmetic Right |  | Logical Right |
| $1 \ll 3 \quad 10$ | 101101 >> 3 | 310101100 >>> 3 |
|  | 110101 | 00010101 |
| 1) >>> 29? = |  | 7. |
| << $n$ ? |  | $x \cdot 2^{n}$. |
| >> $n$ ? $=$ |  | $\left\lfloor x / 2^{n}\right\rfloor$ (i.e., rounded |
| >>> 3) \& (( | 1<<5)-1)? 5- | -bit integer, bits 3-7 of |

