lly Efficient Use of Keys: the Trie

much about cost of comparisons. worst case is length of string. hould throw extra factor of key length, L, into costs: parisons really means $\Theta(ML)$ operations. for key X, keep looking at same chars of X M times. tter? Can we get search cost to be O(L)?

multi-way decision tree, with one decision per character

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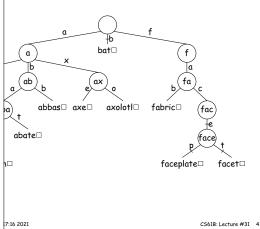
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ed search structures (DS(IJ), Chapter 9

Adding Item to a Trie

ding bat and faceplate.

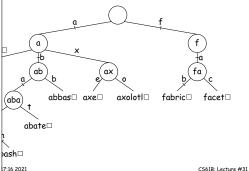
icked.



The Trie: Example

e, abash, abate, abbas, axolotl, axe, fabric, facet} show paths followed for "abash" and "fabric" I node corresponds to a possible prefix.

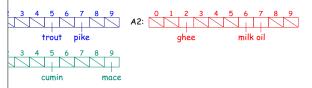
h path to node = that prefix.



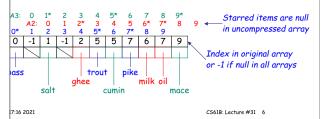
Scrunching Example

(unrelated to Tries on preceding slides)

, each indexed 0..9



them, but keep track of the original index of each item:



A Side-Trip: Scrunching

bvious implementation for internal nodes is array inaracter.

erformance, L length of search key.

independent of N, number of keys. Is there a depen-

ays are sparsely populated by non-null values—waste of

arrays on top of each other!

mpty) entries of one array to hold non-null elements of

arkers to tell which entries belong to which array.

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Practicum

ing idea is cute, but

od if we want to expand our trie.

plicated.

nore useful for representing large, sparse, fixed tables rows and columns.

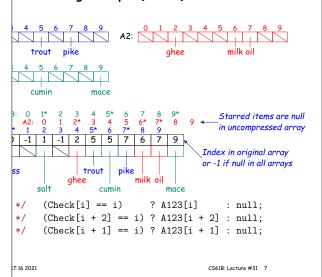
, number of children in trie tends to drop drastically s a few levels down from the root.

ce, might as well use linked lists to represent set of en...

rays for the first few levels, which are likely to have n.

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Scrunching Example (contd.)



robabilistic Balancing: Skip Lists

n be thought of as a kind of n-ary search tree in which put the keys at "random" heights.

hought of as an ordered list in which one can skip large

ple:



tart at top layer on left, search until next step would nen go down one layer and repeat.

, we search for 125 and 127. Gray nodes are looked at; nodes are overshoots.

he nodes were chosen randomly so that there are about nodes that are >k high as there are that are k high.

hes fast with high probability.

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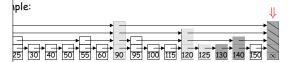
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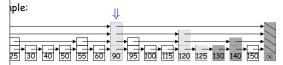
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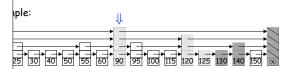
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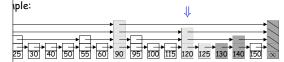
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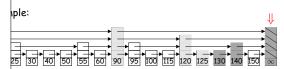
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robabilistic Balancing: Skip Lists

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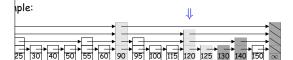
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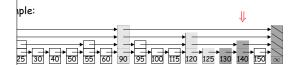
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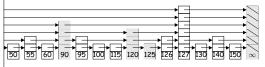
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Example: Adding and deleting

m initial list:



r, we add 126 and 127 (choosing random heights for emove 20 and 40:



s here have been modified.

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robabilistic Balancing: Skip Lists

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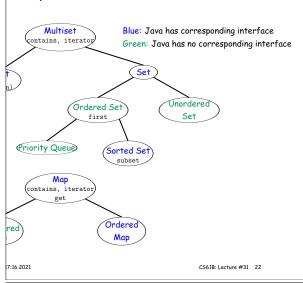
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immary of Collection Abstractions



Summary

arch trees allows us to realize $\Theta(\lg N)$ performance.

-black trees:

 $\left|V\right>$ performance for searches, insertions, deletions. ood for external storage. Large nodes minimize # of ations

performance for searches, insertions, and deletions, s length of key being processed. to manage space efficiently.

idea: scrunched arrays share space.

able $\Theta(\lg N)$ performace for searches, insertions, dele-

plement.

for *interesting ideas*: probabilistic balance, randomstructures.

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Corresponding Classes in Java

ction)

ist, LinkedList, Stack, ArrayBlockingQueue,

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Queue: PriorityQueue Set (SortedSet): TreeSet

d Set: HashSet

Nap: HashMap

(SortedMap): TreeMap

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tructures that Implement Abstractions

linked lists, circular buffers

et

Queue: heaps

Set: binary search trees, red-black trees, B-trees,

arrays or linked lists

d Set: hash table

Nap: hash table

: red-black trees, B-trees, sorted arrays or linked lists

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