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UC Berkeley CS61C : Machine Structures

Lecture 27 – Single-Cycle CPU Control

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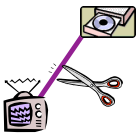
Wireless High Definition? ⇒ Several companies will be working on a "WirelessHD" standard, which will allow HD output devices like laptops, DVD players, video cameras and video games to send HD content @ 5 GB/s!

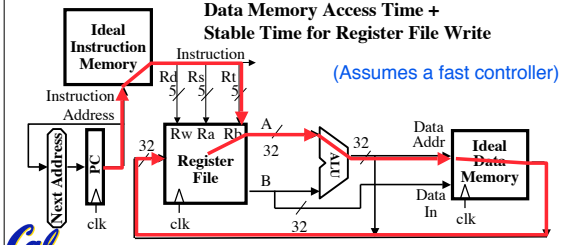


news.com.com/Wireless+HD+specification+due+in+2007/2100-1039_3-6130942.html

An Abstract View of the Critical Path

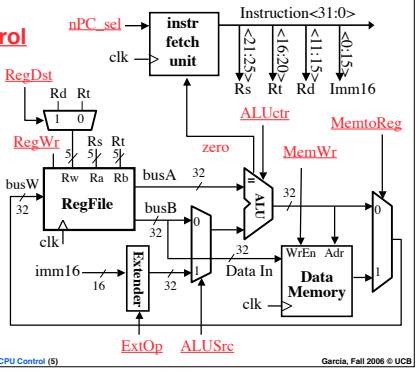
Critical Path (Load Instruction) =
Delay clock through PC (FFs) +
Instruction Memory's Access Time +
Register File's Access Time, +
ALU to Perform a 32-bit Add +
Data Memory Access Time +
Stable Time for Register File Write

(Assumes a fast controller)



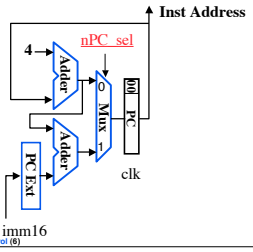
Summary: A Single Cycle Datapath

- **We have everything except control signals**



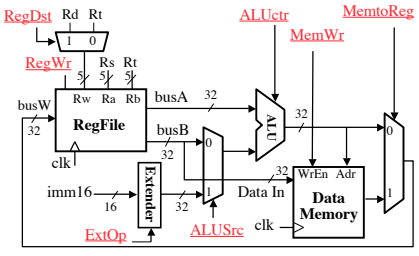
Recap: Meaning of the Control Signals

- **nPC_sel:** "+4" 0 ⇒ PC <- PC + 4
  "br" 1 ⇒ PC <- PC + 4 + {SignExt(Im16) , 00 }
  "n"=next
- **Later in lecture: higher-level connection between mux and branch condition**

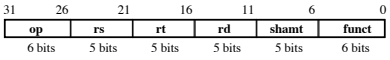


Recap: Meaning of the Control Signals

- **ExtOp:** "zero", "sign"
- **ALUsrc:** 0 ⇒ regB; 1 ⇒ immed
- **ALUctr:** "ADD", "SUB", "OR"
- **MemWr:** 1 ⇒ write memory
- **MemtoReg:** 0 ⇒ ALU; 1 ⇒ Mem
- **RegDst:** 0 ⇒ "rt"; 1 ⇒ "rd"
- **RegWr:** 1 ⇒ write register



RTL: The `Add` Instruction

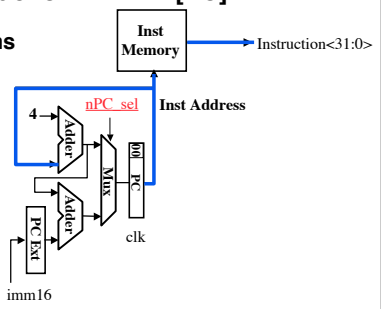


`add rd, rs, rt`

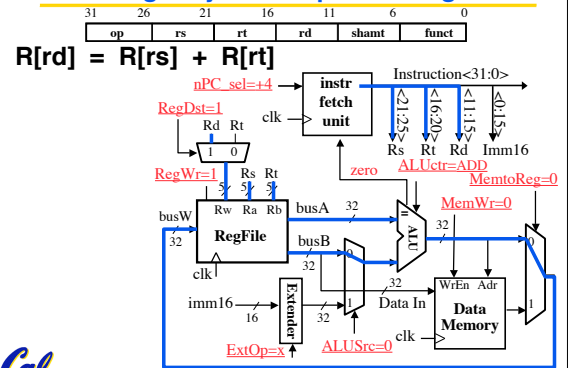
- **MEM[PC]** Fetch the instruction from memory
- **R[rd] = R[rs] + R[rt]** The actual operation
- **PC = PC + 4** Calculate the next instruction's address

Instruction Fetch Unit at the Beginning of Add

- Fetch the instruction from Instruction memory: Instruction = MEM[PC]
  - same for all instructions

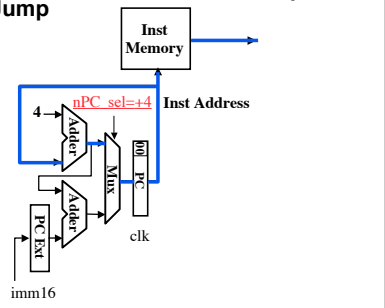


The Single Cycle Datapath during Add

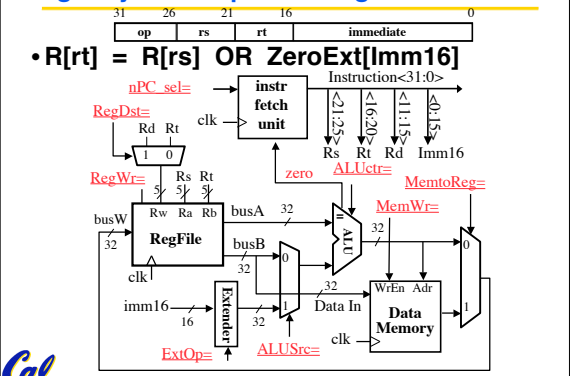


Instruction Fetch Unit at the End of Add

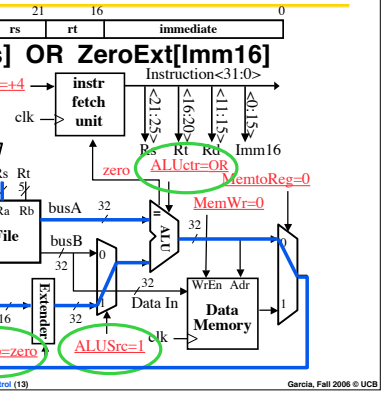
- PC = PC + 4
  - This is the same for all instructions except: Branch and Jump



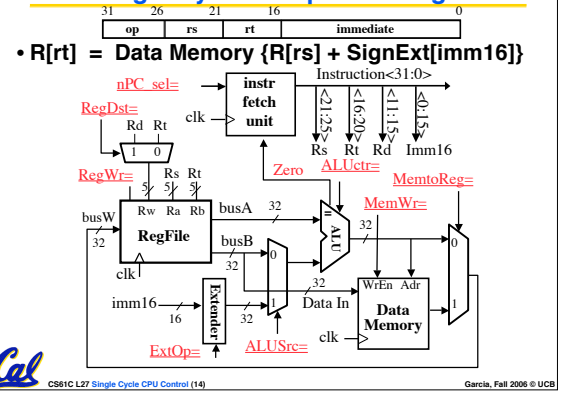
Single Cycle Datapath during Or Immediate?



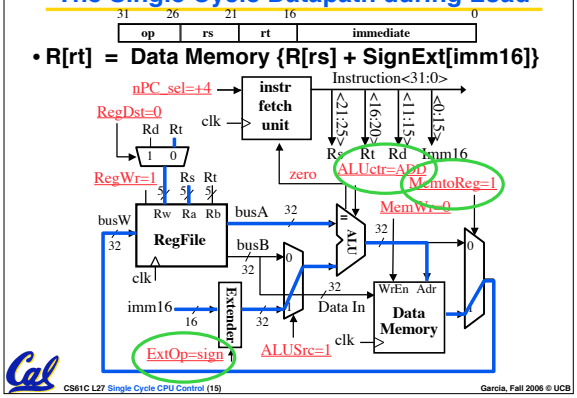
Single Cycle Datapath during Or Immediate?



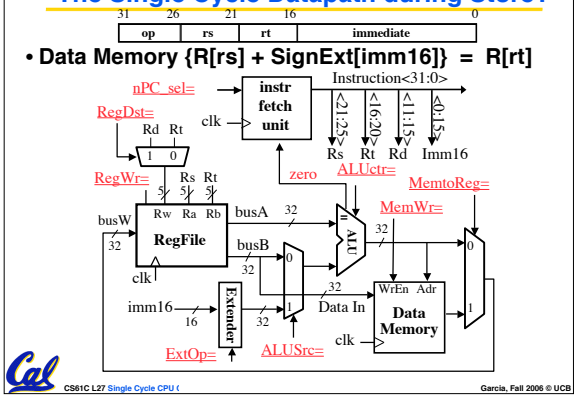
The Single Cycle Datapath during Load?



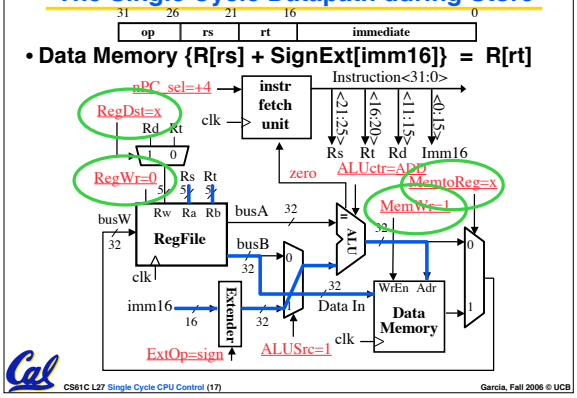
The Single Cycle Datapath during Load



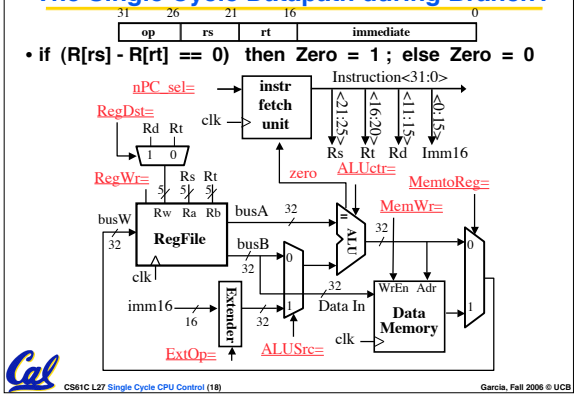
The Single Cycle Datapath during Store?



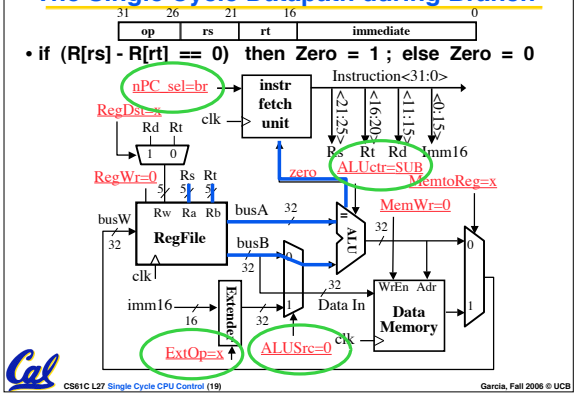
The Single Cycle Datapath during Store



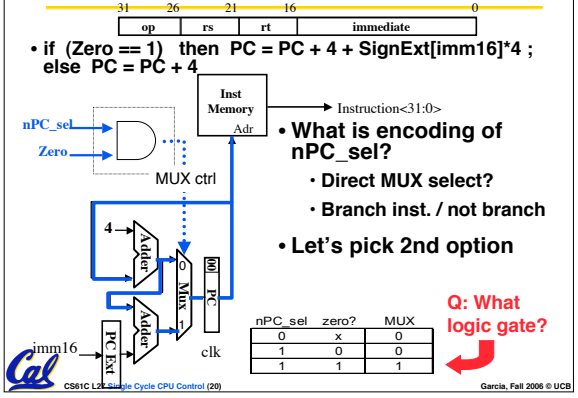
The Single Cycle Datapath during Branch?



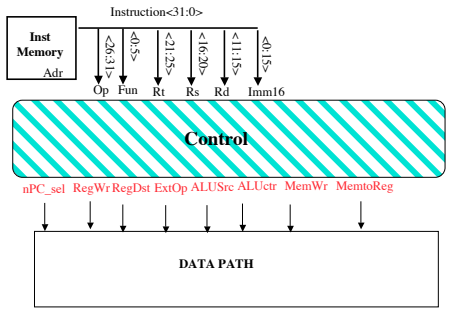
The Single Cycle Datapath during Branch



Instruction Fetch Unit at the End of Branch



Step 4: Given Datapath: RTL -> Control

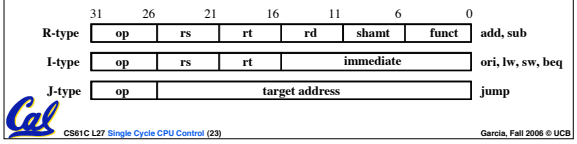


A Summary of the Control Signals (1/2)

inst	Register Transfer
add	R[rd] ← R[rs] + R[rt]; PC ← PC + 4
	ALUsrc = RegB, ALUctr = "ADD", RegDst = rd, RegWr, nPC_sel = "+4"
sub	R[rd] ← R[rs] – R[rt]; PC ← PC + 4
	ALUsrc = RegB, ALUctr = "SUB", RegDst = rd, RegWr, nPC_sel = "+4"
ori	R[rt] ← R[rs] + zero_ext(Imm16); PC ← PC + 4
	ALUsrc = Im, Extop = "Z",ALUctr = "OR", RegDst = rt,RegWr, nPC_sel ="+4"
lw	R[rt] ← MEM[R[rs] + sign_ext(Imm16)]; PC ← PC + 4
	ALUsrc = Im, Extop = "sn", ALUctr = "ADD", MemtoReg, RegDst = rt, RegWr, nPC_sel = "+4"
sw	MEM[R[rs] + sign_ext(Imm16)] ← R[rs]; PC ← PC + 4
	ALUsrc = Im, Extop = "sn", ALUctr = "ADD", MemWr, nPC_sel = "+4"
beq	if (R[rs] == R[rt]) then PC ← PC + sign_ext(Imm16)] ‖ 00 else PC ← PC + 4
	nPC_sel = "br", ALUctr = "SUB"

A Summary of the Control Signals (2/2)

See Appendix A								
	func	10 0000	10 0010	We Don't Care :-)				
	op	00 0000	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
		add	sub	ori	lw	sw	beq	jump
RegDst		1	1	0	0	x	x	x
ALUSrc		0	0	1	1	1	0	x
MemtoReg		0	0	0	1	x	x	x
RegWrite		1	1	1	1	0	0	0
MemWrite		0	0	0	0	1	0	0
nPCsel		0	0	0	0	0	1	0
Jump		0	0	0	0	0	0	1
ExtOp		x	x	0	1	1	x	x
ALUctr<2:0>		Add	Subtract	Or	Add	Add	Subtract	xxx



Boolean Expressions for Controller

RegDst = add + sub
ALUSrc = ori + lw + sw
MemtoReg = lw
RegWrite = add + sub + ori + lw
MemWrite = sw
nPCsel = beq
Jump = jump
ExtOp = lw + sw
ALUctr[0] = sub + beq (assume ALUctr is 0 ADD, 01: SUB, 10: OR)
ALUctr[1] = or

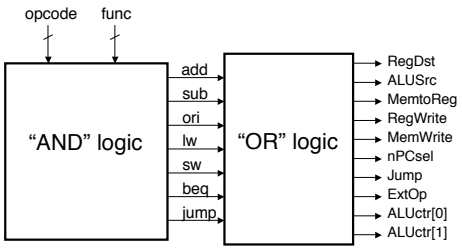
where,

$rtype = \sim op_5 \bullet \sim op_4 \bullet \sim op_3 \bullet \sim op_2 \bullet \sim op_1 \bullet \sim op_0$,
$ori = \sim op_5 \bullet \sim op_4 \bullet op_3 \bullet op_2 \bullet \sim op_1 \bullet op_0$
$lw = op_5 \bullet \sim op_4 \bullet \sim op_3 \bullet \sim op_2 \bullet op_1 \bullet op_0$
$sw = op_5 \bullet \sim op_4 \bullet op_3 \bullet \sim op_2 \bullet op_1 \bullet op_0$
$beq = \sim op_5 \bullet \sim op_4 \bullet \sim op_3 \bullet op_2 \bullet \sim op_1 \bullet \sim op_0$
$jump = \sim op_5 \bullet \sim op_4 \bullet \sim op_3 \bullet \sim op_2 \bullet op_1 \bullet \sim op_0$

How do we implement this in gates?

$add = rtype \bullet func_5 \bullet \sim func_4 \bullet \sim func_3 \bullet \sim func_2 \bullet \sim func_1 \bullet \sim func_0$
$sub = rtype \bullet func_5 \bullet \sim func_4 \bullet \sim func_3 \bullet \sim func_2 \bullet func_1 \bullet \sim func_0$

Controller Implementation



Summary: Single-cycle Processor

° 5 steps to design a processor
- 1. Analyze instruction set => datapath requirements
- 2. Select set of datapath components & establish clock methodology
- 3. Assemble datapath meeting the requirements
- 4. Analyze implementation of each instruction to determine setting of control points that effects the register transfer.
- 5. Assemble the control logic
  - Formulate Logic Equations
  - Design Circuits