

CS61c Summer 2014 Discussion 4 – MIPS Procedures

1 MIPS Control Flow

There are only two instructions necessary for creating and calling functions: `jal` and `jr`. If you follow register conventions when calling functions, you will be able to write much simpler and cleaner MIPS code.

2 Conventions

- How should `$sp` be used? When do we add or subtract from `$sp`?
`$sp` points to a location on the stack to load or store into. Subtract from `$sp` before storing, and add to `$sp` after restoring.
- Which registers need to be saved or restored before using `jr` to return from a function?
 All `$s*` registers that were modified during the function must be restored to their value at the start of the function
- Which registers need to be saved before using `jal`?
`$ra`, and all `$t*`, `$a*`, and `$v*` registers if their values are needed later after the function call.
- How do we pass arguments into functions?
`$a0`, `$a1`, `$a2`, `$a3` are the four argument registers
- What do we do if there are more than four arguments to a function?
 Use the stack to store additional arguments
- How are values returned by functions?
`$v0` and `$v1` are the return value registers.

When calling a function in MIPS, who needs to save the following registers to the stack? Answer “caller” for the procedure making a function call, “callee” for the function being called, or “N.A.” for neither.

\$0	\$v*	\$a*	\$t*	\$s*	\$sp	\$ra
N/A	Caller	Caller	Caller	Callee	N/A	Caller

Now assume a function `foo` calls another function `bar`, which is known to call some other functions. `foo` takes one argument and will modify and use `$t0` and `$s0`. `bar` takes two arguments, returns an integer, and uses `$t0-$t2` and `$s0-$s1`. In the boxes below, draw a possible ordering of the stack just before `bar` calls a function. The top left box is the address of `$sp` when `foo` is first called, and the stack goes downwards, continuing at each next column. Add ‘(f)’ if the register is stored by `foo` and ‘(b)’ if the register is stored by `bar`. The first one is written in for you.

1 <code>\$ra</code> (f)	5 <code>\$t0</code> (f)	9 <code>\$v0</code> (b)	13 <code>\$t1</code> (b)
2 <code>\$s0</code> (f)	6 <code>\$ra</code> (b)	10 <code>\$a0</code> (b)	14 <code>\$t2</code> (b)
3 <code>\$v0</code> (f)	7 <code>\$s0</code> (b)	11 <code>\$a1</code> (b)	15
4 <code>\$a0</code> (f)	8 <code>\$s1</code> (b)	12 <code>\$t0</code> (b)	16

3 A Guide to Writing Functions

```
FunctionFoo: # PROLOGUE
             # begin by reserving space on the stack
             addiu $sp, $sp, -FrameSize

             # now, store needed registers
             sw $ra, 0($sp)
             sw $s0, 4($sp)
             ...
             # BODY
             ...
             # EPILOGUE
             # restore registers
             lw $s0 4($sp)
             lw $ra 0($sp)

             # release stack spaces
             addiu $sp, $sp, FrameSize

             # return to normal execution
             jr $ra
```

4 C to MIPS

1. Assuming \$a0 and \$a1 hold integer pointers, swap the values they point via the stack and return control.

```
void swap(int *a, int *b) {
    int tmp = *a;
    *a = *b;
    *b = tmp;
}
                                     addiu $sp, $sp, -4
                                     lw  $t0, 0($a0)
                                     sw  $t0, 0($sp)
                                     lw  $t0, 0($a1)
                                     sw  $t0, 0($a0)
                                     lw  $t0, 0($sp)
                                     sw  $t0, 0($a1)
                                     addiu $sp, $sp, 4
                                     jr  $ra
```

2. Translate the following algorithm that finds the N^{th} Fibonacci number to MIPS assembly. Assume \$s0 holds N, \$s1 holds fib, \$t0 holds i, and \$t1 hold j.

```
int fib = 1, i = 1, j = 1;
                                     beq  $s0, $0, Ret0
                                     addi $t2, $0, 1
if (N==0)      return 0;
                                     beq  $s0, $t2, Ret1
else if (N==1) return 1;
                                     subi $s0, $s0, 2
N -= 2;
Loop: beq  $s0, $s0, RetF
      add  $s1, $t0, $t1
      addi $t0, $t1, 0
      addi $t1, $s1, 0
      subi $s0, $s0, 1
      j   Loop
Ret0: addi $v0, $0, 0
      j   Done
Ret1: addi $v0, $0, 1
      j   Done
RetF: addi $v0, $0, $s1
Done: jr  $ra
```

What must be done to make this algorithm a callable MIPS function?

Add a prologue and epilogue to reserve space on the stack and store all necessary variables (see #3). Use `$a0` instead of `$s0` to store N, the function's argument.