

You should also look at the problem set problems as well!! And read the notes.

1. Prove by induction that  $\sum_{i=0}^n i = (n+1)n/2$ .
2. Use induction to prove that  $n^5 - n$  is divisible by 5, whenever  $n$  is a nonnegative integer.
3. Rosen. Problem 40.
4. Rosen page 225 question 6.
  - (a) State the well-ordering property for the set of positive integers.
  - (b) Use this property to show that every positive integer can be written as the product of primes
5. Rosen page 199, problem 22 Prove by induction that 6 divides  $n^3 - n$  whenever  $n$  is a nonnegative integer.
6. **Euclid's argument**  
Consider the following result, first proved many centuries ago.

**Theorem 1 (Euclid)** *There exist infinitely many primes.*

**Proof:** Assume to the contrary that there exist finitely many primes. Let these primes (in increasing order) be  $p_1 = 2, p_2 = 3, p_3 = 5, \dots, p_k$ . Let  $q_k = p_1 p_2 p_3 \cdots p_k + 1$ . Note that  $q_k$  is a new number not in the list of primes  $p_1, \dots, p_k$ . At the same time, it is not divisible by  $p_i$  for any  $i$ , since  $q_k \equiv p_1 p_2 p_3 \cdots p_k + 1 \equiv 1 \pmod{p_i}$ , which would mean that  $q_k$  is a new prime different from  $p_1, \dots, p_k$ , which is a contradiction. This completes the proof.  $\square$

Let  $p_1, \dots, p_k$  represent the first  $k$  primes. Are we guaranteed that  $p_1 p_2 p_3 \cdots p_k + 1$  is always prime for all  $k \geq 1$ ? Is the above proof valid? Explain.

7. **Stable marriage**
  - (a) Consider the following situation. After we find the pairing using the boy-optimal algorithm we described in class, one of the boys  $b$ , who gets paired up with some girl  $g$ , changes his mind. I.e., he changes his preference by moving  $g$  higher up in his list.  
Is the current pairing stable? Is it still boy-optimal? Prove or disprove these statements.
  - (b) Prove or disprove that any boy optimal stable pairing is girl pessimal.
  - (c) Prove that the traditional marriage algorithm is boy optimal.
8. **Number theory**
  - (a) True or false? You should give some explanations.
    - i. 2 has an inverse modulo 6.
    - ii.  $\gcd(32, 1387) = 1$ .

- (b) Find an inverse of 3 modulo 17.
- (c) Find an inverse of 144 mod 171.
- (d) Solve this equation:  $9x = 2 \pmod{40}$ .
- (e) (Rosen page 165, problem 11.) Find  $\gcd(2n + 1, 3n + 2)$ , where  $n$  is a positive integer.
- (f) Which integers leave a remainder of 1 when divided by 2, leave a remainder of 2 when divided by 3, and leave a remainder of 3 when divided by 5? (Chinese remainder theorem. HW 4.)
- (g) A *least common multiple* of positive integers  $a$  and  $b$ , denoted by  $\text{lcm}(a, b)$ , is the smallest positive integer that is divisible by  $a$  and  $b$ . Prove that  $ab = \gcd(a, b)\text{lcm}(a, b)$ .
- (h) (Rosen page 149, problem 15.) Show that if  $p$  is prime, the only solution of  $x^2 - 1 = 0 \pmod{p}$  are integers  $x$  such that  $x = 1 \pmod{p}$  or  $x = -1 \pmod{p}$ .
- (i) What is  $a^{(p-1)(q-1)} \pmod{pq}$  for  $a$  relatively prime to  $p$  and  $q$  and  $p$  and  $q$  are primes?
- (j) Show that  $a^x \pmod{p} = a^{x \pmod{p-1}} \pmod{p}$ . You may use Fermat's Little Theorem.

## 9. RSA

- (a) Describe the RSA encryption and decryption scheme. (Assume  $p$  and  $q$  are prime. Let  $n = pq$ . Use  $x$  as a message,  $(e, n)$  as a public key, and  $(d, n)$  as a private key.)
  - (b) Analyze the encryption and decryption times for RSA.
  - (c) Compare decoding and encoding time. Usually  $e = 3$ , why?
10. (a) Prove that any degree 1 polynomial over a field  $F$  has at most 1 zero. What is the maximum number of points where two distinct degree 1 polynomials can intersect?
- (b) If at least  $d$  points on a degree  $\leq d - 1$  polynomial have value 5, what is the polynomial?
- (c) Given pairs  $(x_1, y_1), (x_2, y_2)$ , describe a degree 1 polynomial that hits these points.

## 11. Applications of polynomials.

- (a) Describe a method to share a secret with  $b$  bits with  $n$  people such that each person is given at most  $b/k + 1$  bits and any  $k$  of the people can reconstruct all of the bits, and any  $k - 1$  people won't know anything about the first  $b/k$  bits of the secret. (Use a Galois Field,  $GF(p)$  for your field, and state how big  $p$  should be.)
- (b) Say, you receive  $(1, 3), (2, 4), (3, 7)$  and  $(4, 9)$  for a degree 1 polynomial code. Which point is corrupted? (Hint: drop each and see if the other three are on a line.) If at most 2 points are corrupted in a degree  $n - 1$  polynomial code, extend this simple algorithm for finding the original message. Extend this algorithm to correct  $k$  points. Assuming that finding a degree  $n$  polynomial that fits  $n$  points takes  $x$  time, how long would your algorithm take.
- (c) Review the notes on error correction, why is the Welsh-Berlekemp method better than your simple method above when the number of errors  $k$  becomes larger.
  - i. What is the number of possible legal number of messages for degree  $n - 1$  polynomial code over a field of size  $f$  where the message consists of the value of the polynomial on  $n + 2k$  points?
  - ii. Consider a noisy channel that can choose  $k$  places to change, and change them to any character. How many actions can the channel take?
  - iii. If you recover the original message, can you figure out exactly what the channel did as well?
  - iv. How many bits must the received message contain to decode?

12.
  - (a) Show that a Gray code exists for any  $b$ -bit strings.
  - (b) Show that any graph where every node has degree larger than  $n/2$  has a Hamiltonian cycle.
  - (c) For what values of  $n$  does the dimension  $n$  hypercube (with  $2^n$  nodes) have an Eulerian path.