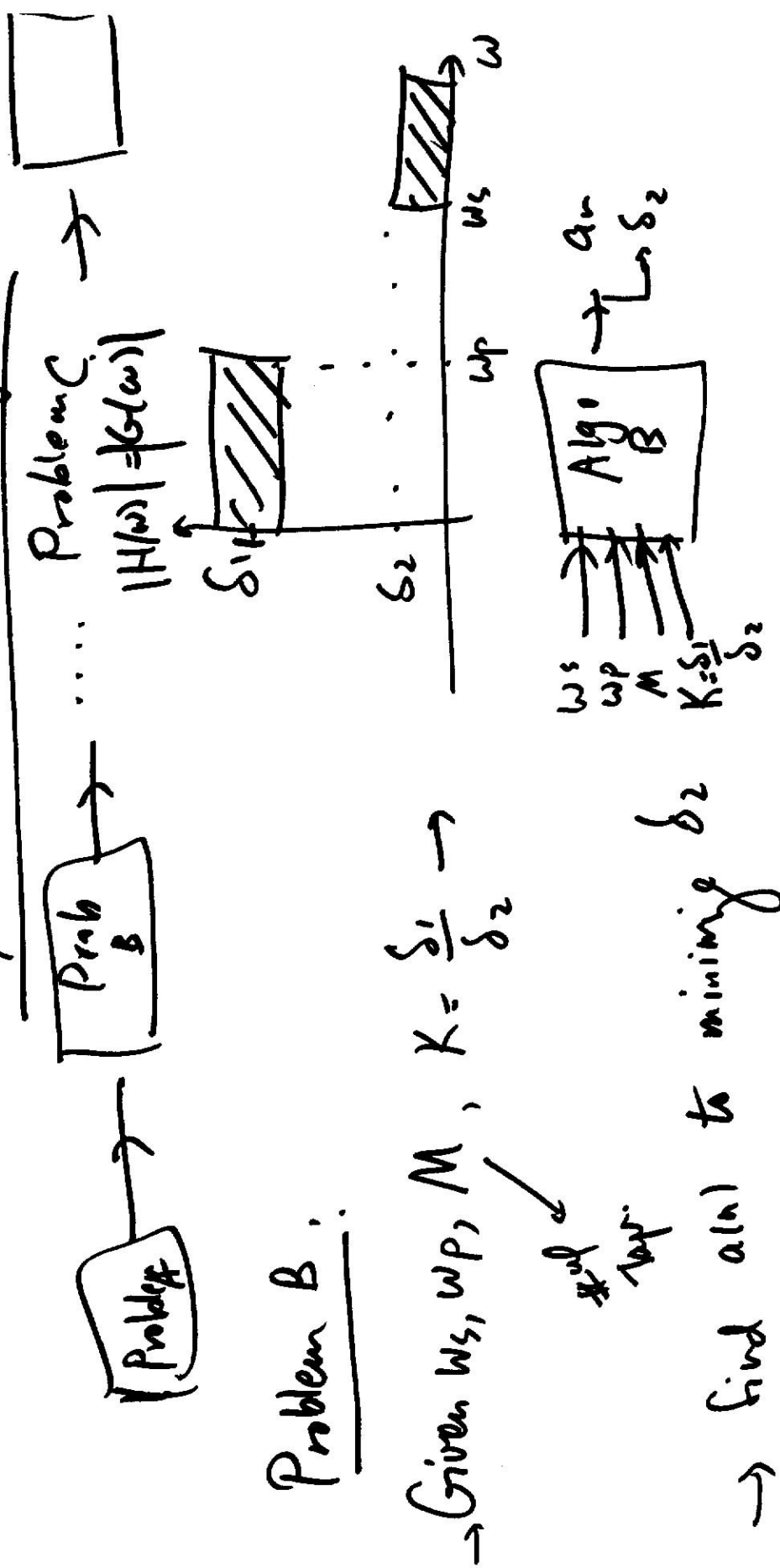


oct 16, 03

# Optimum FIR Filter Design



Problem B:

→ Given  $w_s, w_p, M, K = \frac{\delta_1}{\delta_2}$  →

# of taps

→ find  $a(n)$  to minimize  $\delta_2$

Problem C:  $E(\omega) = W(\omega) [G(\omega) - D(\omega)]$

where  $W(\omega) =$  positive weighting function  $= \begin{cases} \frac{1}{k} & I_1 \\ 1 & I_2 \end{cases}$

$$G(\omega) = \sum_{n=0}^M a(n) \cos(n\omega)$$

$$D(\omega) = \text{Desired Freq Response} = \begin{cases} 1 & I_1 \\ 0 & I_2 \end{cases}$$

→ Find  $a(n)$  to minimize

$$\text{Max}_{\omega \in F} |E(\omega)|$$

where  $F = I_1 \cup I_2$ .

$F$  is a subset of a closed interval  $0 \leq \omega \leq \pi$



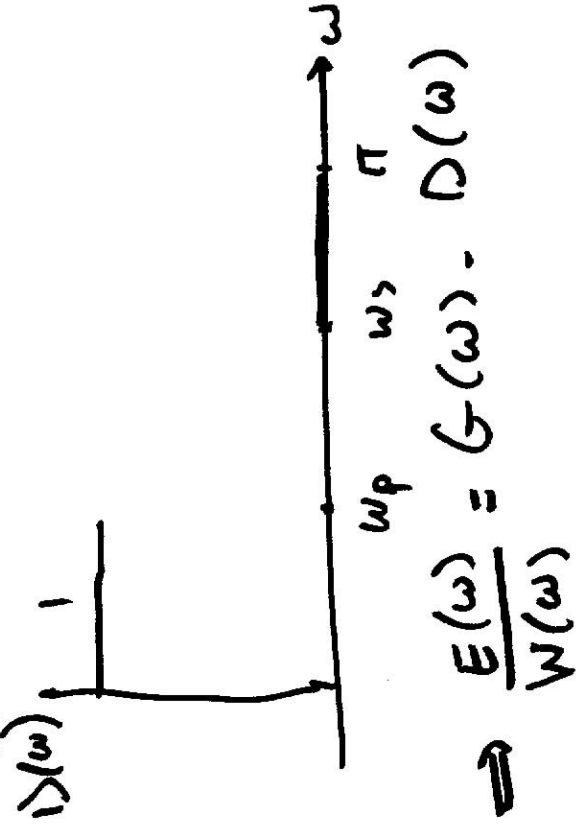
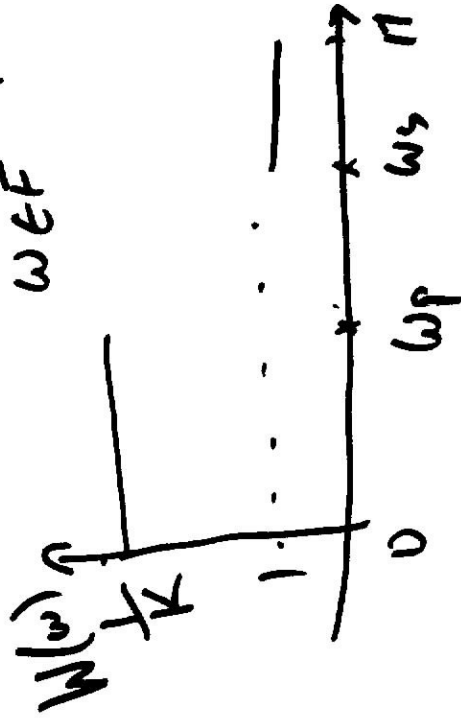
$I_1 = \text{pass band}$

$I_2 = \text{stopband}$

Show: Problem C is equivalent to Problem B.

start by showing  $\theta = \delta_2$

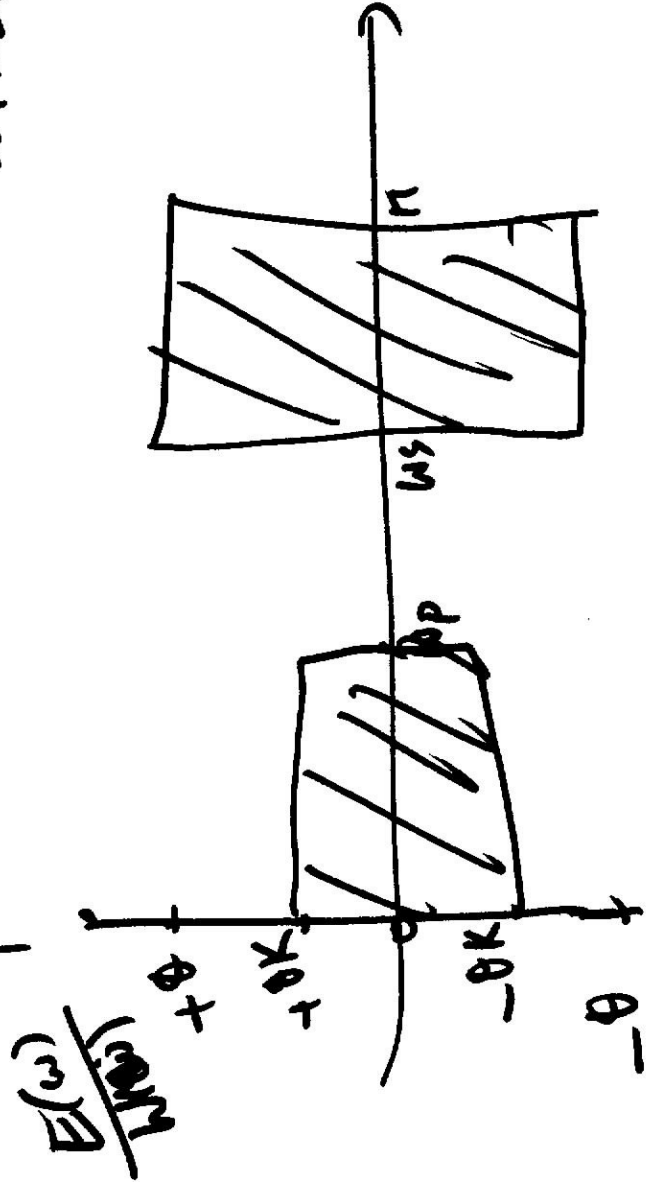
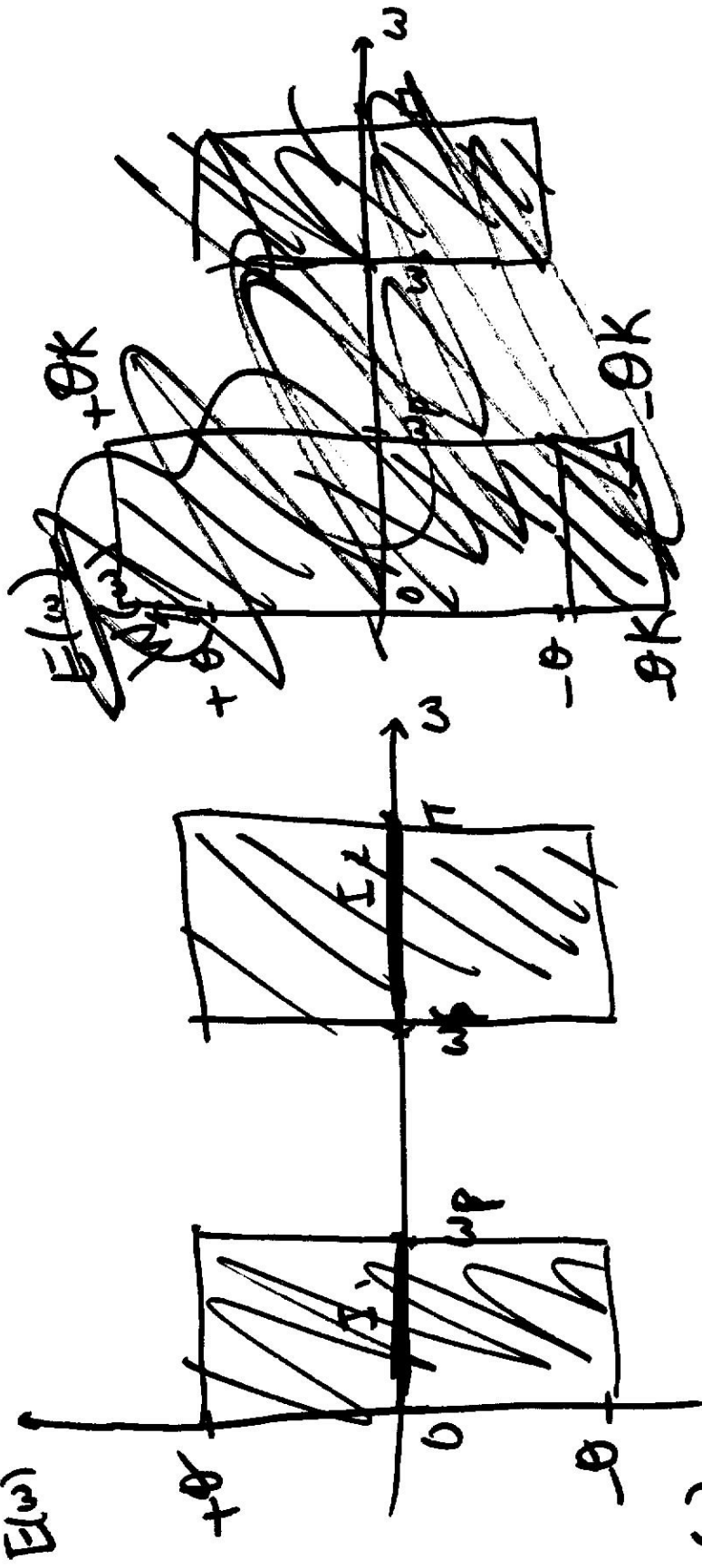
Show  $\max_{\omega \in F} |E(\omega)| \stackrel{??}{=} \delta_2$



$$E(\omega) = W(\omega) [G(\omega) - D(\omega)] \Rightarrow \frac{E(\omega)}{W(\omega)} = G(\omega) - D(\omega)$$

$$\Rightarrow G(\omega) = D(\omega) + \frac{E(\omega)}{W(\omega)}$$

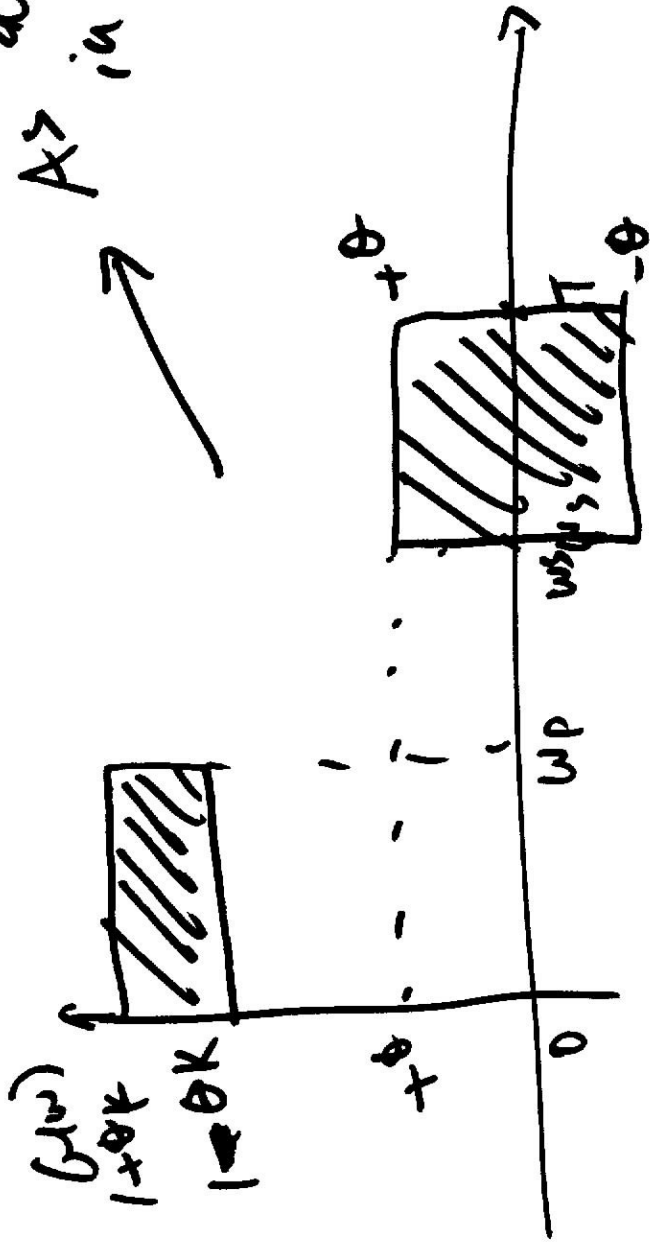
By def:  $\max_{\omega \in F} |E(\omega)| = \theta$



Recall  
 $G(\omega) = D(\omega) + \frac{E(\omega)}{W(\omega)}$

done  
As in Problem C.

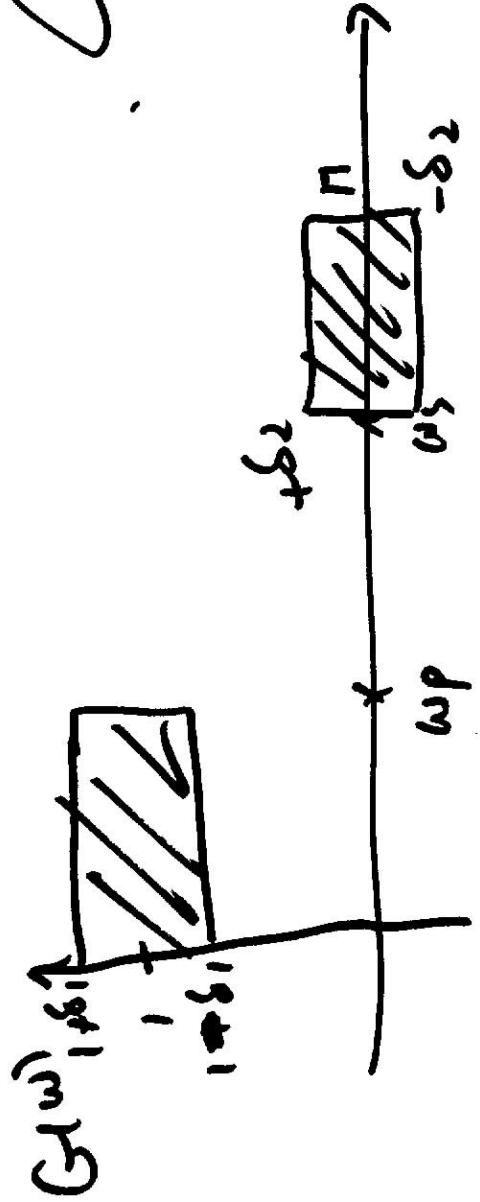
$$K = \frac{\delta_1}{\delta_2}$$



How do the  $G(w)$  look like as in Problem A or B?

Conclusion

$$\theta = \delta_2$$



Conclusion:

Problem C: find  $a(n)$  To minimize  $S_2$ .

Problem D:  $n$  To minimize  $S_2$ .

$\Rightarrow$  Same problem.

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Problem A  $\rightarrow$  Problem B  $\rightarrow$  Problem C.

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Solve Problem C:

Alternation Thm:

$F =$  union of closed interval.

Let  $P(x)$  be an  $r$ th order polynomial.

$$P(x) = \sum_{k=0}^r a_k x^k$$

$D =$  desired  $f_1$  continuous in  $F$ .

$W =$  positive  $f_2$  [  $D(x) - P(x)$  ]

$$E(x) = W(x) \left| \max_{x \in F} |E(x)| \right|$$

$$\text{Define } \|E\| = \max_{x \in F} |E(x)|$$

Necessary & Sufficient Condition for  $P(x)$

To be a unique  $r$ th order polynomial that minimizes  $\|E\|$  is that  $E(x)$  exhibits at least  $r+2$  alternations

i.e. There is at least  $r+2$  values of  $x$

Assue.

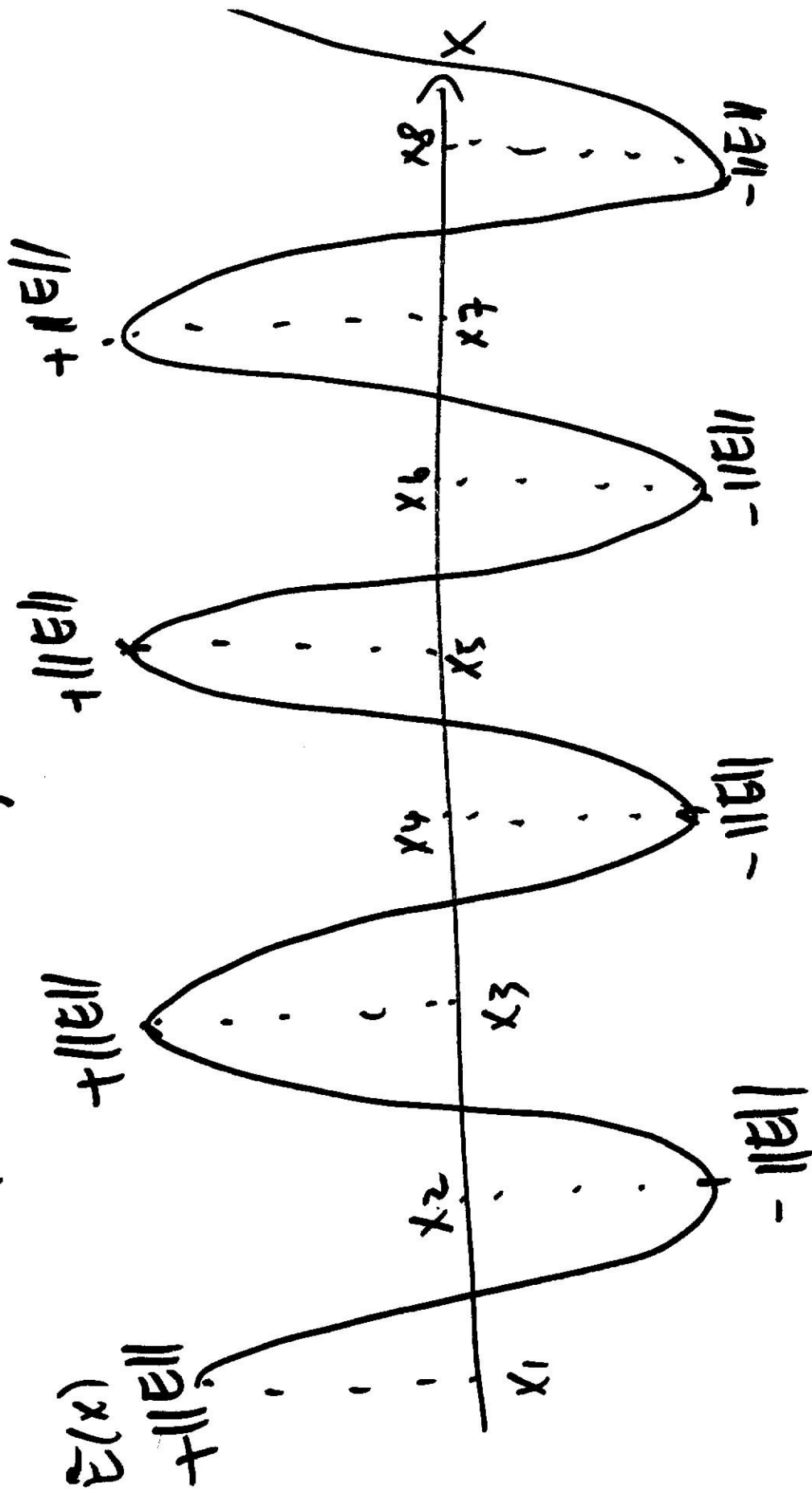
If

Then.

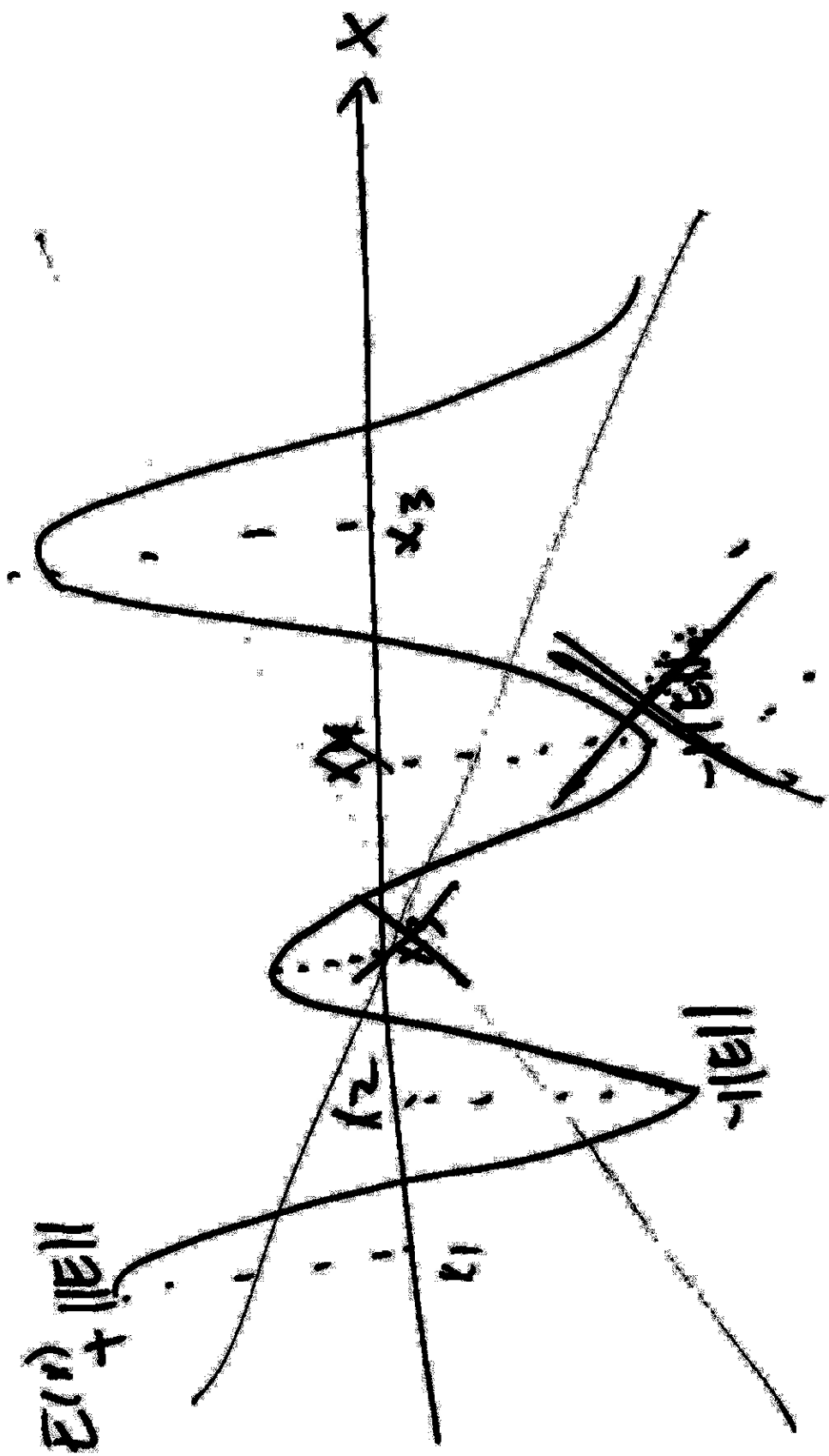
Such that

$$x_1 < x_2 < \dots < x_{r+2}$$

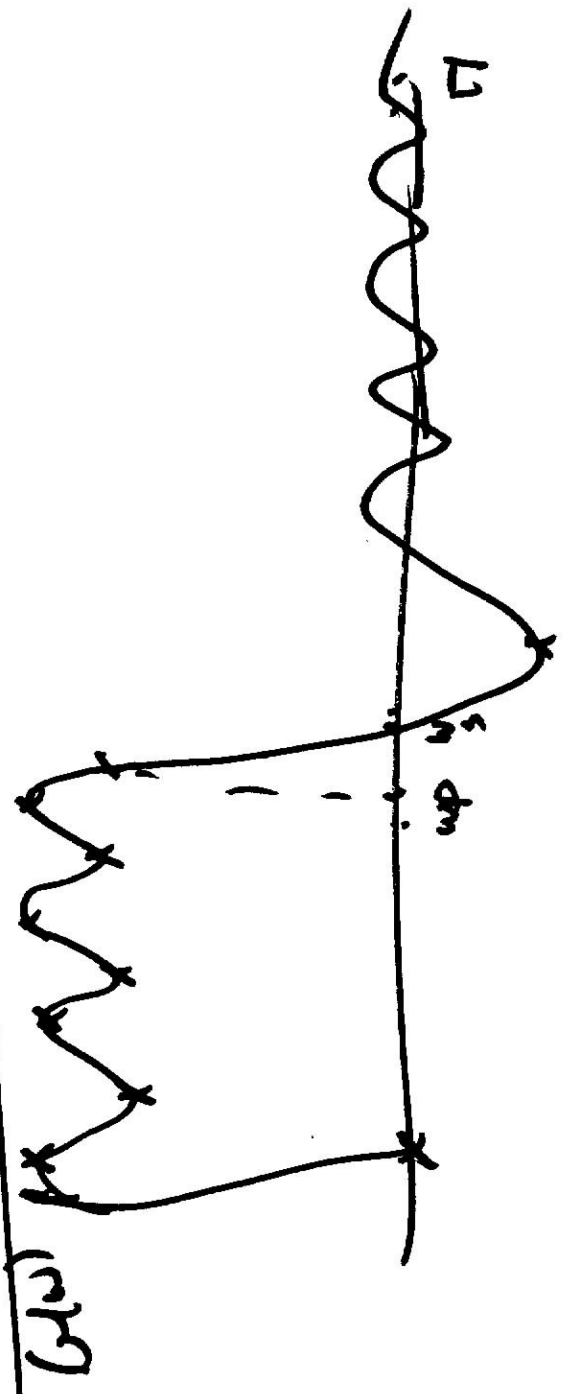
$$E(x_i) = -E(x_{i+1}) = \pm \|E\|$$





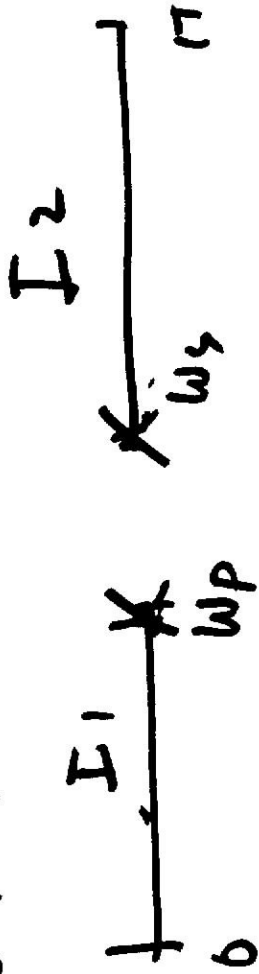


Recall  $G(w)$  (lost test) can have at most  $n-1$  local extrema.



Alt. Then.  
 Case 5 about  $F = I_1, I_2, \dots$ , possible to have alternation at  $w_p$  & at  $w_s$  even though  $w_p, w_s$  are not local extrema for  $G$ .

$$F := I_1 \cup I_2$$



What can we say about ~~local~~  $\max$  #  
 of local extrema of  $G(x)$  in  $M+1+2$

$$F = I_1 \cup I_2 \quad \underline{M+3} \text{ local extrema}$$

$\Rightarrow$  AT most  $\underline{M+3}$  for  $G$  in  $I_1 \cup I_2$

- AT Th at least  $M+2$   $\Rightarrow$  either  $M+2$  or  $M+3$  local extrema in  $F$