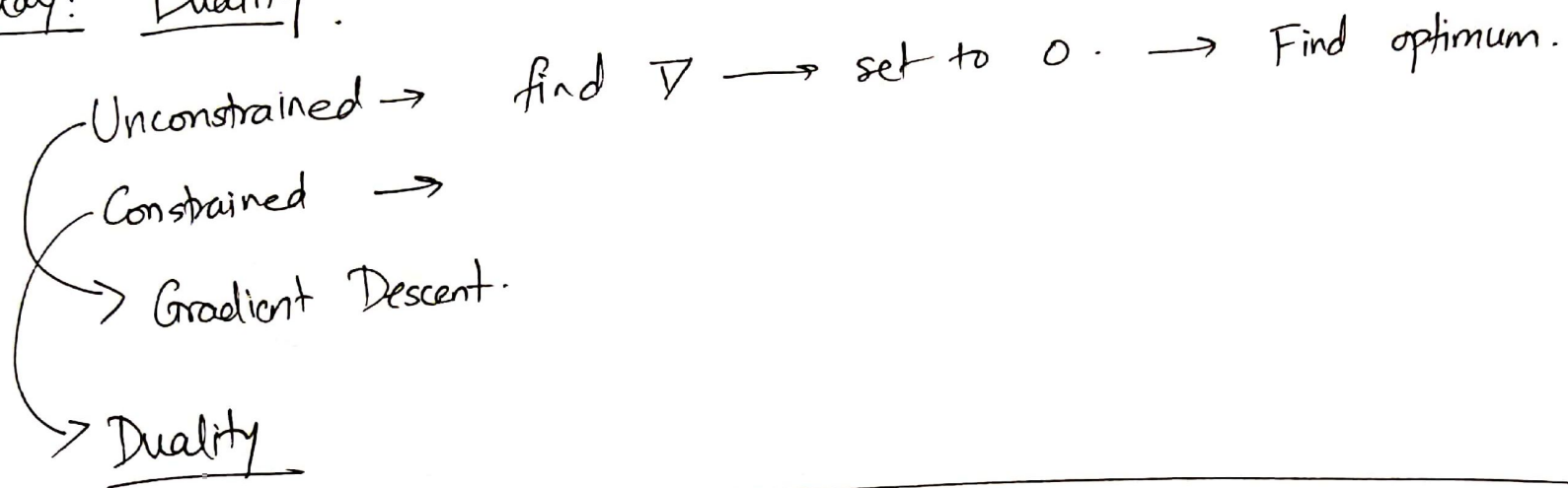


Today: Duality.



$$p^* = \begin{array}{l} \text{minimize } f_0(\vec{x}) \\ \text{s.t. } f_i(\vec{x}) \leq 0 \quad \forall 1 \leq i \leq m \\ h_i(\vec{x}) = 0 \quad \forall i \leq i < p \end{array} \quad] \quad \text{THE PRIMAL}$$

For this: we define the Lagrangian

$$L(\vec{x}, \vec{\lambda}, \vec{v}) = f_0(\vec{x}) + \sum_{i=1}^m \lambda_i f_i(\vec{x}) + \sum_{i=1}^p v_i h_i(\vec{x}).$$

When $\lambda_i \geq 0$

$\vec{\lambda}, \vec{v}$ are called Lagrange multipliers.
dual variables.

o

$$\min_{\vec{x}} L(\vec{x}, \vec{\lambda}, \vec{v}) := \underbrace{g(\vec{\lambda}, \vec{v})}_{\text{function } \vec{\lambda}, \vec{v}}$$

$$x = x_1$$

$$x = x_2$$

$$\vdots$$

$$L_1(\vec{\lambda}, \vec{v}) = L(\vec{x}_1, \vec{\lambda}, \vec{v})$$

$$L_2(\vec{\lambda}, \vec{v}) = L(\vec{x}_2, \vec{\lambda}, \vec{v})$$

$$\vdots$$

$$\inf_{\vec{x}} \{L_1, L_2, \dots, L\}$$

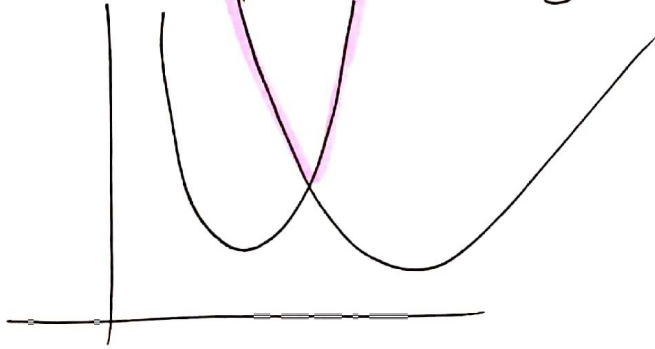
(2)

Observations

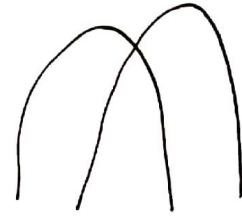
- ① g depends on $\vec{\lambda}, \vec{v}$.
- ② $L(\vec{x}, \vec{\lambda}, \vec{v})$ is an affine function of $\vec{\lambda}, \vec{v}$.

Reminder: Lemma:

Pointwise maximum of convex functions \Rightarrow Convex.

$$f_3(\vec{x}) = \max \{f_1(\vec{x}), f_2(\vec{x})\}$$


Pointwise minimum of concave functions \Rightarrow concave.



- ③ $g(\vec{\lambda}, \vec{v})$ is concave!! $\rightarrow \vec{\lambda}, \vec{v}$
 \hookrightarrow Does not depend on $f_0(x)$ at all.

④ Property of g :

$$g(\vec{\lambda}, \vec{v}) \leq p^*$$

for all $\vec{\lambda} \geq 0$, \vec{v} .
element-wise.

③

Proof:

Consider: \tilde{x} a feasible point for the primal

$$f_i(\tilde{x}) \leq 0 \quad (1) \quad 1 \leq i \leq m$$

$$h_i(\tilde{x}) = 0 \quad (2)$$

$$\lambda_i f_i(\tilde{x}) \leq 0 \quad (1^*)$$

$$v_i h_i(\tilde{x}) = 0 \quad (2^*)$$

$$\Rightarrow \sum \lambda_i f_i(\tilde{x}) \leq 0$$

$$\Rightarrow \sum v_i h_i(\tilde{x}) = 0$$

Consider:

$$L(\tilde{x}, \vec{\lambda}, \vec{v}) = f_0(\tilde{x}) + \underbrace{\sum \lambda_i f_i(\tilde{x})}_{\leq 0} + \underbrace{\sum v_i h_i(\tilde{x})}_{=0}$$

$$\leq f_0(\tilde{x})$$

$$g(\vec{\lambda}, \vec{v}) = \inf_{\tilde{x}} L(\tilde{x}, \vec{\lambda}, \vec{v})$$

$g(\vec{\lambda}, \vec{v})$ lowerbounds all values of f_0 for feasible \tilde{x}

$$g(\vec{\lambda}, \vec{v}) \leq p^*$$

An interpretation:

$$M(\vec{x}) = f_0(\vec{x}) + \sum \mathbb{1}_{\text{neg}} \{f_i(x)\} + \sum \mathbb{1}_0(h_i(x)).$$

$$\mathbb{1}_{\text{neg}} = \begin{cases} 0 & \text{if } f_i(x) \leq 0 \\ \infty & \text{if } f_i(x) > 0 \end{cases}$$

$$\mathbb{1}_0 = \begin{cases} 0 & \text{if } h_i(x) = 0 \\ \infty & \text{otherwise.} \end{cases}$$

$$\min_{\vec{x}} M(\vec{x})$$

→ Hard thresholding penalty for violating a constraint : $M(\vec{x})$
 Instead : the Lagrangian gives you only a linear penalty in $\vec{\lambda}$ or $\vec{\nu}$

example:

Minimum norm problem.

A is full row rank.

$$P^* = \min_{\vec{x}} \vec{x}^T \vec{x} \\ \text{st. } A\vec{x} = \vec{b} \Leftrightarrow A\vec{x} - \vec{b} = 0$$



$$L(\vec{x}, \vec{v}) = \vec{x}^T \vec{x} + \vec{v}^T (A\vec{x} - \vec{b}) \quad : \text{Convex function } \vec{x}$$

$$g(\vec{v}) = \min_{\vec{x}} L(\vec{x}, \vec{v})$$

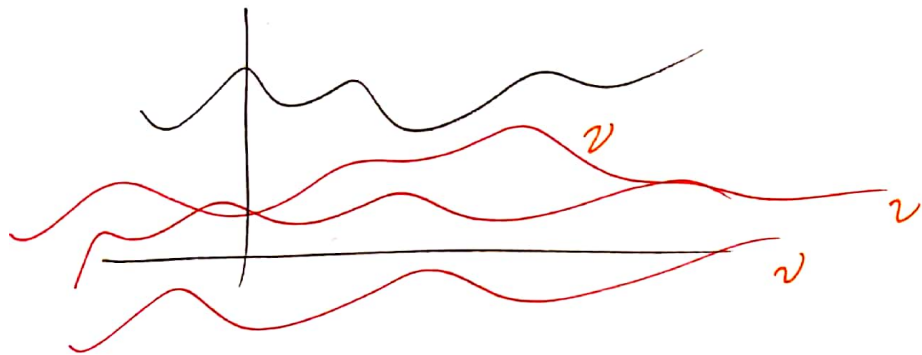
$$\text{Minimize } : \nabla_{\vec{x}} L = 0.$$

$$\nabla_{\vec{x}} L(\vec{x}, \vec{v}) = 2\vec{x} + A^T \vec{v}$$

$$\text{setting to } 0 \Rightarrow \vec{x} = -\frac{1}{2} A^T \vec{v}$$

$$g(\vec{v}) = L\left(-\frac{1}{2} A^T \vec{v}, \vec{v}\right) = \frac{1}{4} \vec{v}^T A A^T \vec{v} + \vec{v}^T (A(-\frac{1}{2} A^T \vec{v}) - \vec{b}) \\ = -\frac{1}{4} \vec{v}^T A A^T \vec{v} - \vec{v}^T \vec{b}$$

$$P^* \geq g(\vec{v})$$



Lagrange Dual Problem!

$$d^* = \max_{\substack{\vec{\lambda} \geq 0 \\ \vec{v}}} g(\vec{\lambda}, \vec{v})$$

$$g(\vec{\lambda}, \vec{v})$$

CONVEX PROGRAM.

$$\max g(\vec{v}) = d^*$$

$$\nabla_{\vec{v}} g(\vec{v}) = -\frac{1}{4} (2AA^T) \vec{v} - \vec{b}$$

$$\text{Set} = 0.$$

$$\vec{v}_* = -2 (AA^T)^{-1} \cdot \vec{b}$$

$$\vec{x}_1 = -\frac{1}{2} A^T \vec{v}_* = A^T (AA^T)^{-1} \cdot \vec{b}.$$

We know:

$$d^* \leq p^*$$

Weak duality.

$$d^* = p^*$$

Strong duality.

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Example:

$$p^* = \begin{array}{l} \text{min.} \\ \text{st. } x_i^2 = 1 \end{array}$$

$$\vec{x}^T W \vec{x}$$

$$i=1, \dots, n.$$

$$W \in S^n$$

$$x_i = +1 \text{ or } x_i = -1.$$

Partitioning problem.

w_{ij}

$$\vec{x}^T \begin{bmatrix} w_1 & & & 0 \\ & w_2 & & \\ & & \dots & \\ 0 & & & w_n \end{bmatrix} \vec{x}$$

$$L(x, \vec{v}) = x^T W x + \sum_{i=1}^n v_i (x_i^2 - 1).$$

$$= \underbrace{x^T (W + \text{diag}(\vec{v})) x}_{\text{quadratic}} - \sum_{i=1}^n v_i$$

$$x^T Q \cdot x \geq 0 \quad \text{if } Q \text{ PSD.}$$

if $W + \text{diag}(\vec{v})$ is not PSD.

$$g(\vec{v}) = \inf_{\vec{x}} L(\vec{x}, \vec{v}) = \begin{cases} -\infty \\ -\sum_{i=1}^n v_i \end{cases}$$

if $W + \text{diag}(\vec{v})$ is PSD ≥ 0

$$\max g(\vec{v})$$

$$\text{st. } W + \text{diag}(\vec{v}) \text{ is PSP}$$

Let us pick some value for \vec{v} .

$$\vec{v} = - \begin{bmatrix} \lambda_{\min}(W) \\ \lambda_{\min}(W) \\ \vdots \\ \lambda_{\min}(W) \end{bmatrix}$$

$$W - \lambda_{\min} I \succeq 0.$$

$$\text{Lower bound} = n \cdot \lambda_{\min}(W) = \sum_{i=1}^n \nu_i$$

$$p^* \geq n \cdot \lambda_{\min}(W).$$

$$\nu_i = \lambda_{\min}(W).$$

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