

Dual of Logistic regression.

$$\vec{x}_i \in \mathbb{R}^m$$

Recall logistic regression:

Data points (\vec{x}_i) , labels $y_i = +1$ or -1 .

Want: $\vec{w}^T \vec{x} + \beta = \log \frac{p(\vec{x})}{1-p(\vec{x})}$, where $p(\vec{x})$ is the probability that the data point belongs to class 1.

We showed that finding the best \vec{w}^T, β is the same as:
(per the maximum-likelihood estimator)

$$\text{Max } \prod_{i=1}^n P(y_i = y_i)$$

$$\text{Maximize } \prod_{i=1}^n \left(\frac{\exp(y_i (\vec{w}^T \vec{x}_i + \beta))}{1 + \exp(y_i (\vec{w}^T \vec{x}_i + \beta))} \right) \quad \text{Non-convex problem.}$$

We can take logs to turn this into a convex problem.

$$\text{Maximize } \log \left(\prod_{i=1}^n \left(\frac{\exp(y_i (\vec{w}^T \vec{x}_i + \beta))}{1 + \exp(y_i (\vec{w}^T \vec{x}_i + \beta))} \right) \right)$$

$$= \text{maximize } \sum_{i=1}^n \log \left(\frac{1}{1 + \exp(-y_i (\vec{w}^T \vec{x}_i + \beta))} \right)$$

$$= \text{maximize } \sum_{i=1}^n -\log \left(1 + \exp(y_i (\vec{w}^T \vec{x}_i + \beta)) \right)$$

Instead of $\max_x f(x)$, we consider $-\min_x f(x)$. (2)

The solution x^* to the latter also solves the former.

So consider:

$$p^* = \min_{\vec{w}, \beta} \sum_{i=1}^n \log(1 + \exp(y_i (\vec{w}^T \vec{x}_i + \beta)))$$

$\rightarrow \log\left(\frac{1}{P(Y_i = y_i)}\right)$

This is a convex problem.

It is unconstrained? What is its dual?

Let $f(t) = \log(1 + e^{-t})$.

Consider $\beta = 0$ for simplicity.

$$p^* = \min_{\vec{w}} \sum_{i=1}^n \log(1 + \exp(-y_i (\vec{w}^T \cdot \vec{x}_i)))$$

$$= \min_{\vec{w}} \sum_{i=1}^n \log(1 + \exp(-y_i \vec{x}_i^T \cdot \vec{w}))$$

Define $A = [y_1 \vec{x}_1, y_2 \vec{x}_2, \dots, y_n \vec{x}_n]$. and $\vec{v} = A^T \vec{w}$; $\vec{v} = \begin{bmatrix} v_1 \\ v_2 \\ \vdots \\ v_m \end{bmatrix}$

$$= \min_{\vec{v}, \vec{w}} \sum_{i=1}^n \log(1 + \exp(-v_i))$$

s.t. $\vec{v} = A^T \vec{w}$

$$1 + \exp(-v_i) = \frac{1}{P(Y_i = y_i)}$$

Now we have a "constraint".

$$p^* = \min_{\vec{v}, \vec{w}} \sum_{i=1}^n f(v_i). \quad \text{Convex.}$$

$$\vec{v} = A^T \vec{w}$$

(3)

• Only constraints are linear equality + convex problem.

• If there is a feasible point, i.e. some point

s.t. $\vec{v} = A^T \vec{w}$, then Slater's condition holds, i.e. $p^* = d^*$.

What is d^* ?

$$L(\vec{v}, \vec{w}, \vec{z}) = \sum_{i=1}^n f(v_i) + \vec{z}^T (\vec{v} - A^T \vec{w})$$

$$g(\vec{z}) = \min_{\vec{v}, \vec{w}} L(\vec{v}, \vec{w}, \vec{z}) = \min_{\vec{v}} \min_{\vec{w}} L(\vec{v}, \vec{w}, \vec{z})$$

if $A\vec{z} \neq 0$.

$$= \min_{\vec{v}} \begin{cases} -\infty \\ \sum_{i=1}^n f(v_i) + \vec{z}^T \vec{v} \end{cases} \quad \text{if } A\vec{z} = 0.$$

Consider $h(t) = f(t) + 2_i t = \log(1 + e^{-t}) + 2_i t$.

$$\frac{dh}{dt} = \frac{1}{1+e^{-t}} (e^{-t})(-1) + 2_i = \frac{-1}{1+e^t} + 2_i$$

Setting = 0. $\frac{1}{1+e^t} = 2_i \Rightarrow 1+e^t = \frac{1}{2_i}$

$$\Rightarrow e^t = \frac{1-2_i}{2_i} \Rightarrow t = \log\left(\frac{1-2_i}{2_i}\right).$$

Note $\frac{1-2_i}{2_i} \geq 0$ only when $2_i \in [0, 1]$.

if $2_i < 0$, $\lim_{t \rightarrow \infty} h(t) = -\infty$.

if $2_i > 1$, since $\frac{d}{dt} \log(1+e^{-t}) = \frac{-1}{1+e^t}$ has slope < -1 ,

$2_i t$ dominates as $t \rightarrow -\infty \Rightarrow \lim_{t \rightarrow -\infty} h(t) = -\infty$.

So for $z_i \in [0, 1]$.

$$e^t = \frac{1-z_i}{z_i} \quad ; \quad t = \log\left(\frac{1-z_i}{z_i}\right).$$

$$h(z_i) = \log\left(1 + e^{-\log\left(\frac{1-z_i}{z_i}\right)}\right) + z_i \left(\log\left(\frac{1-z_i}{z_i}\right)\right).$$

$$= \log\left(1 + \frac{z_i}{1-z_i}\right) + z_i \left(\log\left(\frac{1-z_i}{z_i}\right)\right).$$

$$= \log\left(\frac{1}{1-z_i}\right) + z_i \log\left(\frac{1}{z_i}\right) + z_i \cdot \log(1-z_i).$$

$$= -z_i \log z_i - (1-z_i) \log(1-z_i)$$

= entropy of z_i

$$\therefore g(\vec{z}) = \begin{cases} \sum_{i=1}^n -z_i \log z_i - (1-z_i) \log(1-z_i) \\ -\infty \end{cases}$$

if $A\vec{z} = 0$
and $z_i \in [0, 1] \forall i$

otherwise.

Dual: $\max_{\vec{z}} \sum_{i=1}^n -z_i \log z_i - (1-z_i) \log(1-z_i)$

s.t. $z_i \in [0, 1] \forall i$
 $A\vec{z} = 0.$

What is z_i ? $\frac{1}{P(Y_i = y_i)} = \frac{1}{1-z_i}$ So $z_i = P(Y_i \neq y_i)$

$$A\vec{z} = 0 \Rightarrow [\vec{x}_1 \dots \vec{x}_n] \begin{bmatrix} y_1 (1 - P(Y_1 = y_1)) \\ \vdots \\ y_n (1 - P(Y_n = y_n)) \end{bmatrix} = 0$$

Note $y_i (1 - P(Y_i = y_i)) = 1 - P(\vec{x}_i)$ if $y_i = 1$
 $= P(\vec{x}_i)$ if $y_i = -1$.

Total Least Squares.

In a normal least squares problem (sometimes called OLS for ordinary least squares) we try to find \vec{x} such that

$$A\vec{x} \approx \vec{b} \quad \text{and we minimize}$$

$$\|\vec{e}\|_2^2 = \|A\vec{x} - \vec{b}\|_2^2 \quad \text{In this formulation, we are}$$

assuming that the errors in our data are only in \vec{b} . But

what if there are also errors in A ?

i.e. we have $[A + \tilde{A}]\vec{x} = \vec{b} + \tilde{\vec{b}}$

\uparrow error in A
 \uparrow error in \vec{b}

$$\begin{aligned}
 m > n \\
 A \in \mathbb{R}^{m \times n} \\
 \vec{b} \in \mathbb{R}^m
 \end{aligned}$$

\tilde{A} and $\tilde{\vec{b}}$ are perturbations we do not know. We wish to find \vec{x} s.t. $[A + \tilde{A}]\vec{x} = \vec{b} + \tilde{\vec{b}}$ but all we have is A and \vec{b} .

How to measure the perturbations? And minimize them? i.e. find A, \vec{b} closest to $A + \tilde{A}, \vec{b} + \tilde{\vec{b}}$.

minimize $\|\tilde{A} \quad \tilde{\vec{b}}\|_F^2$

s.t. $[A + \tilde{A}]\vec{x} = \vec{b} + \tilde{\vec{b}}$

$$[A + \tilde{A}]\vec{x} = \vec{b} + \tilde{\vec{b}}$$

$$[A + \tilde{A} \mid \vec{b} + \tilde{\vec{b}}] \begin{bmatrix} \vec{x} \\ -1 \end{bmatrix} = 0_{m \times 1}$$

So we want $\begin{bmatrix} \vec{x} \\ -1 \end{bmatrix} \in N \left[\underbrace{[A + \tilde{A} \mid \vec{b} + \tilde{\vec{b}}]}_{\substack{m \text{ rows} \\ n+1 \text{ cols}}} \right]$

So this matrix must be rank deficient.

Define: $[A + \tilde{A} \mid \vec{b} + \tilde{\vec{b}}] = \tilde{\vec{z}} \in \mathbb{R}^{m \times (n+1)}$

$[A \mid \vec{b}] = \vec{z} \in \mathbb{R}^{m \times (n+1)}$

minimize: $\|\vec{z} - \tilde{\vec{z}}\|_F^2$

s.t. $\text{Rank}(\tilde{\vec{z}}) = n$

→ solution by least-sq.