



EE16B

# Designing Information Devices and Systems II

Lecture 11A

Moore Penrose Inverse  
Sampling and Interpolation  
(Polynomial Interpolation)

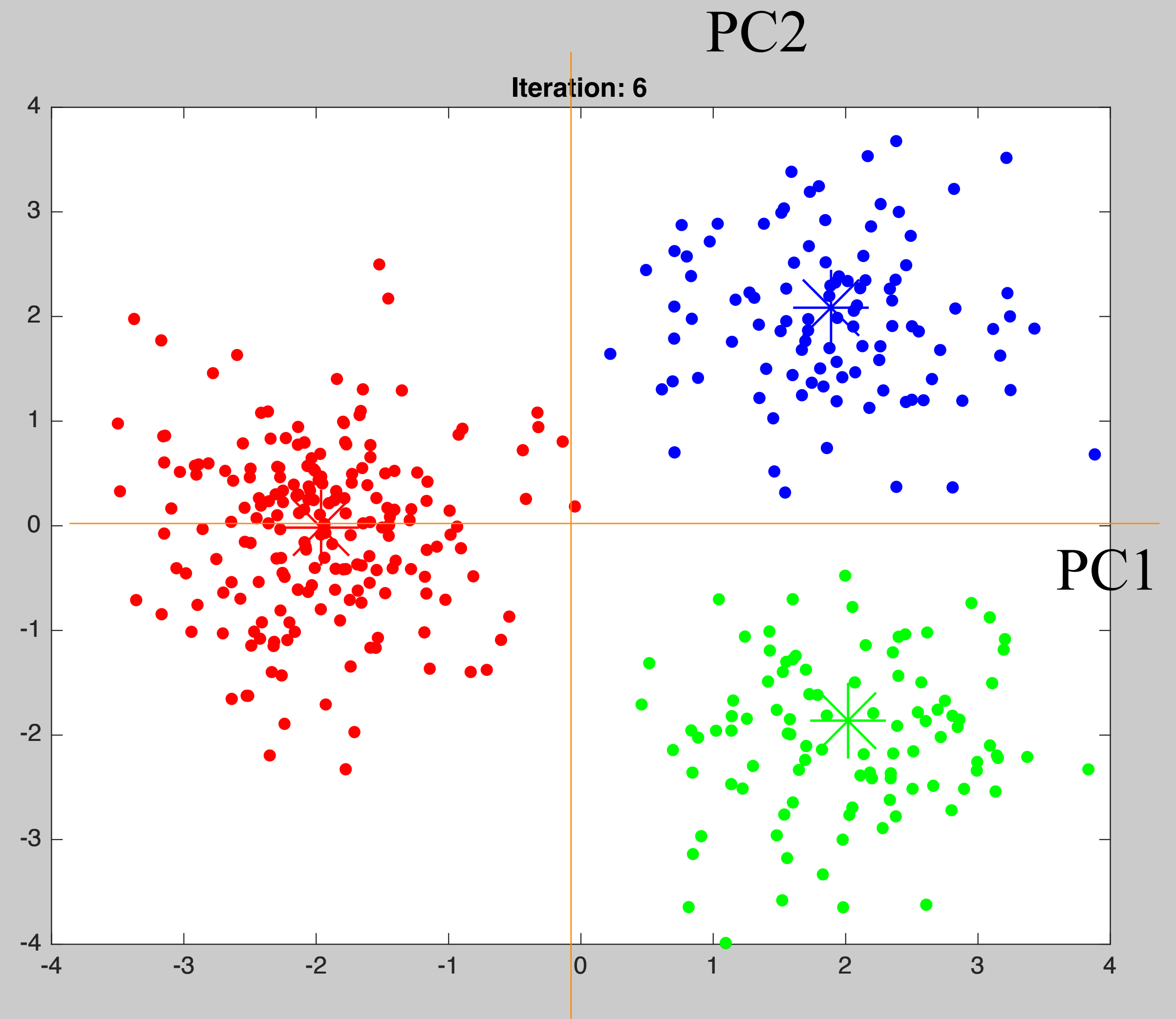
# Intro

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- Last time:
  - Examples of PCA
  - Labeled vs non-labeled clustering
- Today
  - Finish the PCA/SVD Module  
Moore Penrose Pseudo-Inverse
  - New module: Sampling and interpolation
  - Polynomial interpolation

# Labeled VS non labeled Classification

Word1  
Word2  
Word3  
Word4  
Word5  
Word6  
Word7  
Word8



# Labeled VS non labeled Classification

Word1

Word2

Word3

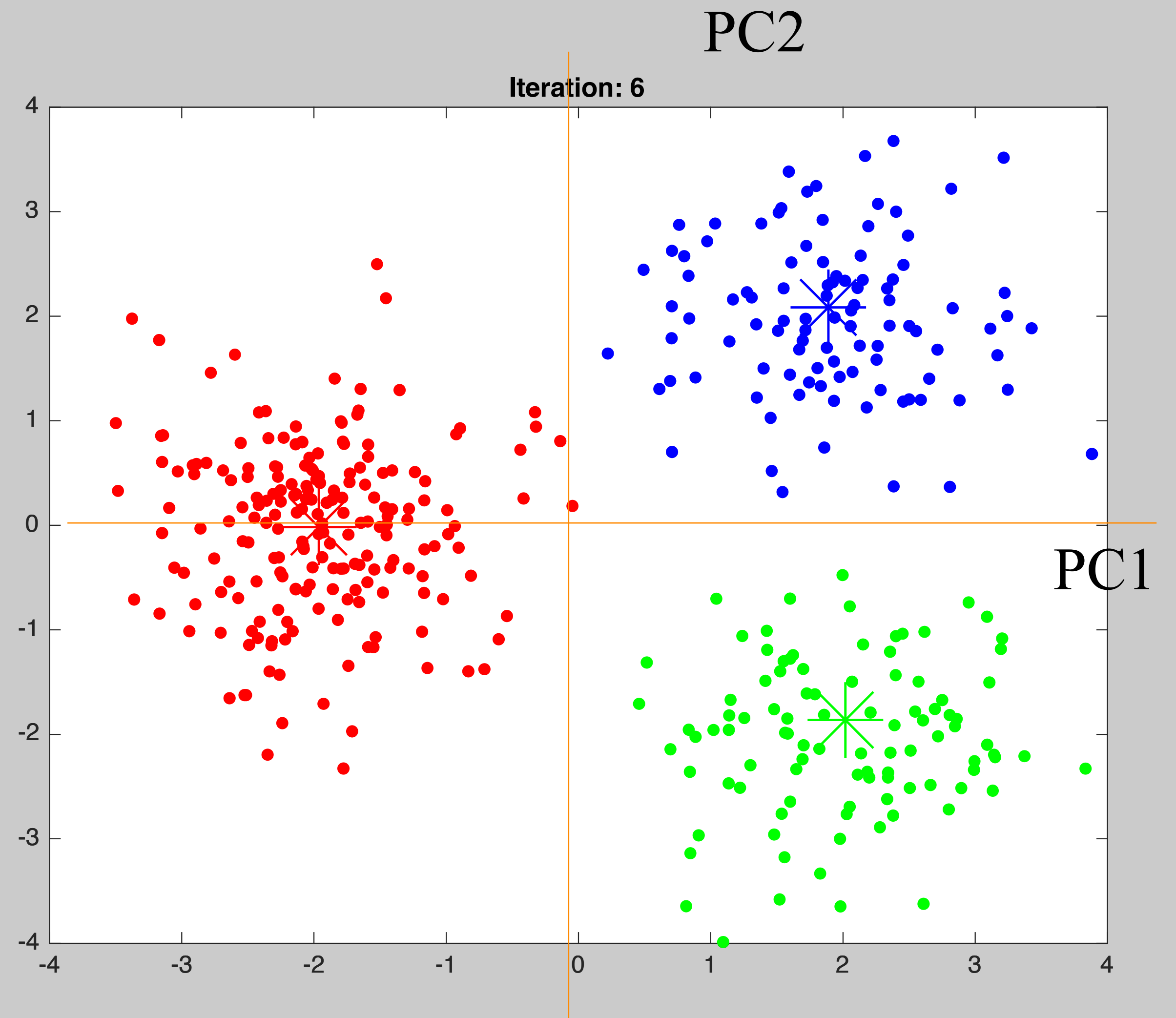
Word4

Word5

Word6

Word7

Word8



# Labeled VS non labeled Classification

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“Banana”

“Banana”

“Banana”

“Mango”

“Mango”

“Mango”

“Chop”

“Chop”

“Chop”

PC2

PC1

# Matrix Inverse, and Pseudo-Inverse

- Square, full rank (N) A:

$$Ax = y$$

$$x = A^{-1}y$$

$$A^{-1}A = AA^{-1} = I_{N \times N}$$

- Tall, full rank (N) A (Least squares):

$$Ax = y$$

$$x = (A^T A)^{-1} A^T y$$

$$(A^T A)^{-1} A^T A = I_{N \times N}$$

Also, left inverse



# Pseudo Inverse and the SVD

$$Ax = y$$

$$x = (A^T A)^{-1} A^T y$$

• SVD:

$$A = U_1 S V_1^T$$

$$(A^T A) = V_1 S U_1^T U_1 S V_1^T = V_1 S^2 V_1^T$$

$$(A^T A)^{-1} = V_1 S^{-2} V_1^T$$

$$(A^T A)^{-1} A^T = V_1 S^{-2} V_1^T V_1^T S U_1^T = V_1 S^{-2} S U_1^T =$$

$$= V_1 S^{-1} U_1^T = A^\dagger$$

Moore-Penrose Pseudo-inverse

# Under-determined Linear Systems

- Fat, full rank (N) A:

$$Ax = y \quad \text{Infinite solutions!}$$



- Claim:

$$AA^\dagger = I_{N \times N} \quad \text{Also, right inverse}$$

- SVD:

$$A = U_1 S V_1^T$$

$$AA^\dagger = U_1 S V_1^T V_1 S^{-1} U_1^T = U_1 S S^{-1} U_1^T = U_1 U_1^T = I_{N \times N}$$



# Under-determined Linear Systems

- Fat, full rank ( $N$ )  $A$ :

$$Ax = y \quad \text{Infinite solutions!}$$

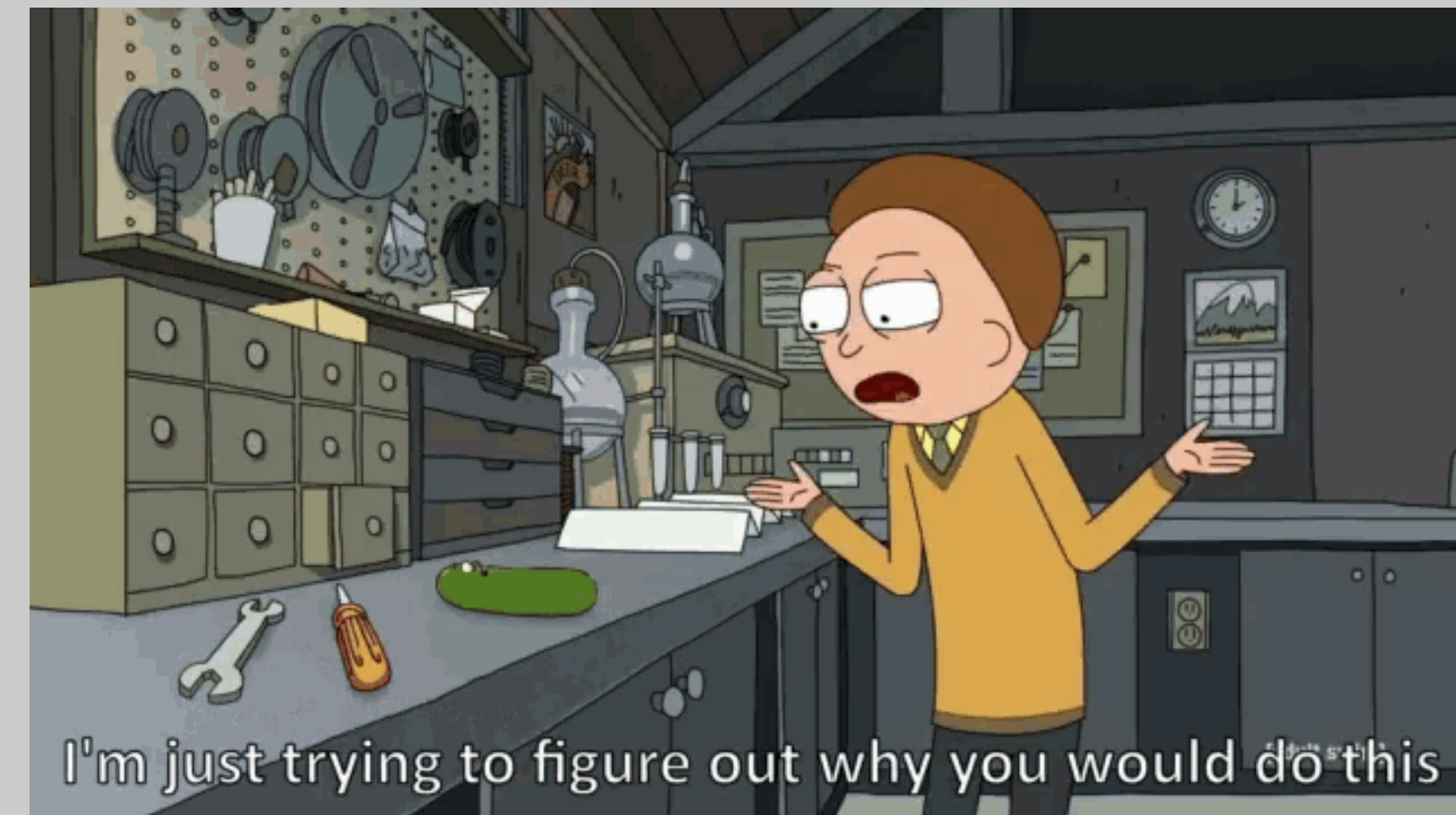
- Fact:

$$AA^\dagger = I_{N \times N}$$

- A Solution:

$$\hat{x} = A^\dagger y$$

$$A\hat{x} = AA^\dagger y = Iy = y$$



- A minimum-norm solution!

$$\|\hat{x}\| < \|\vec{x}\| \quad \forall A\vec{x} = y$$

# Minimum - Norm

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- A minimum-norm solution!  $\hat{x} = A^\dagger y$
- Proof outline:  $\|\hat{x}\| < \|\vec{x}\| \quad \forall A\vec{x} = y$

$$\vec{x} = \hat{x} + \tilde{x}$$

– Show that  $A\tilde{x} = 0$

– Show that  $\hat{x}^T \tilde{x} = 0$

– Show that  $\|\vec{x}\|^2 = \|\hat{x} + \tilde{x}\|^2 > \|\hat{x}\|^2$

# Back to Open Loop Control

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$$\vec{x}(t) - A^t \vec{x}(0) = \underbrace{\begin{bmatrix} A^{t-1}B & A^{t-2}B & \dots & AB & B \end{bmatrix}}_{R_t} \begin{bmatrix} u(0) \\ u(1) \\ \vdots \\ u(t-1) \end{bmatrix}$$

System is controllable. What if  $t \gg N$ ?  
Whats a good sequence of  $u(t)$ ?

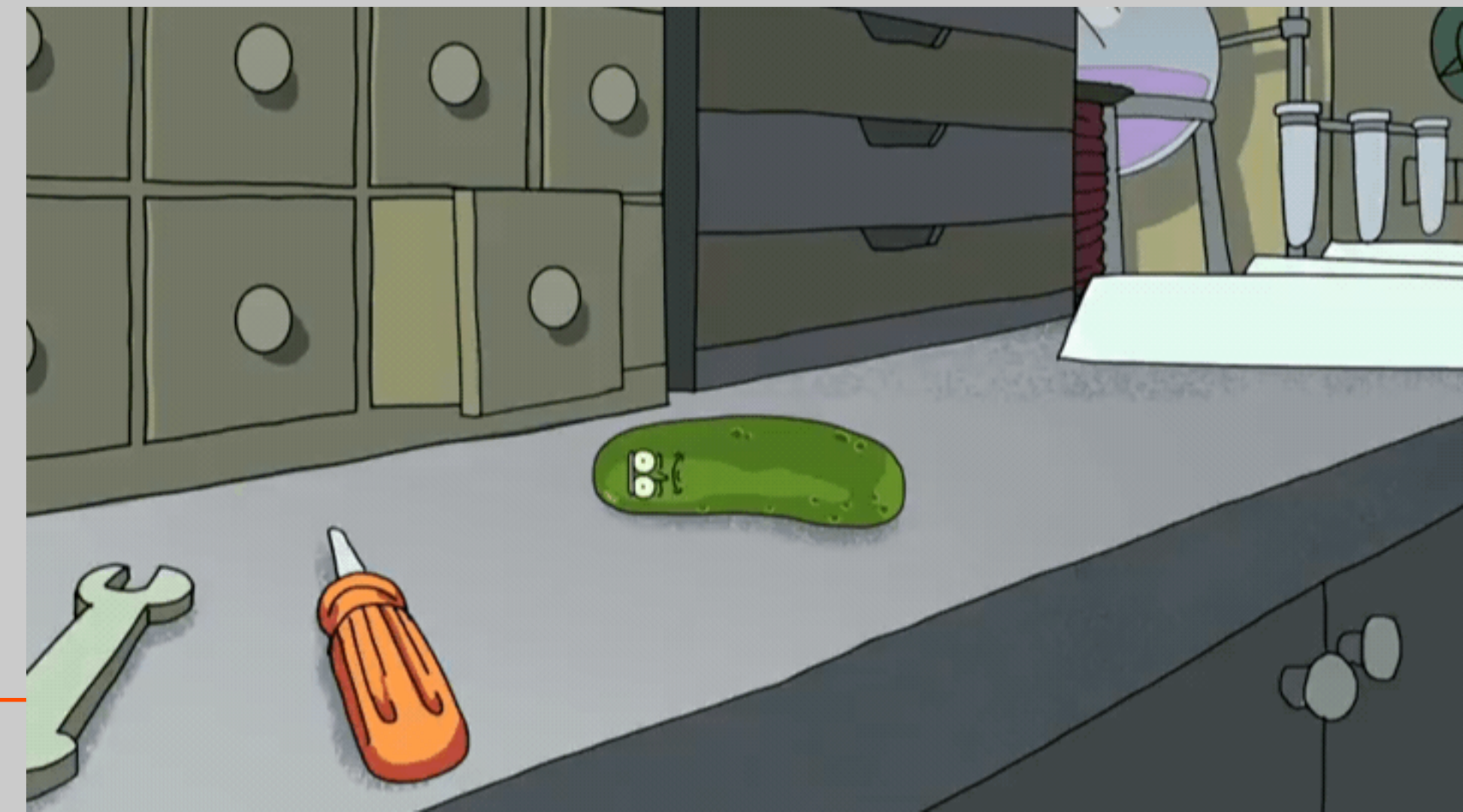
# Summary:

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- Moore penrose Pseudo Inverse

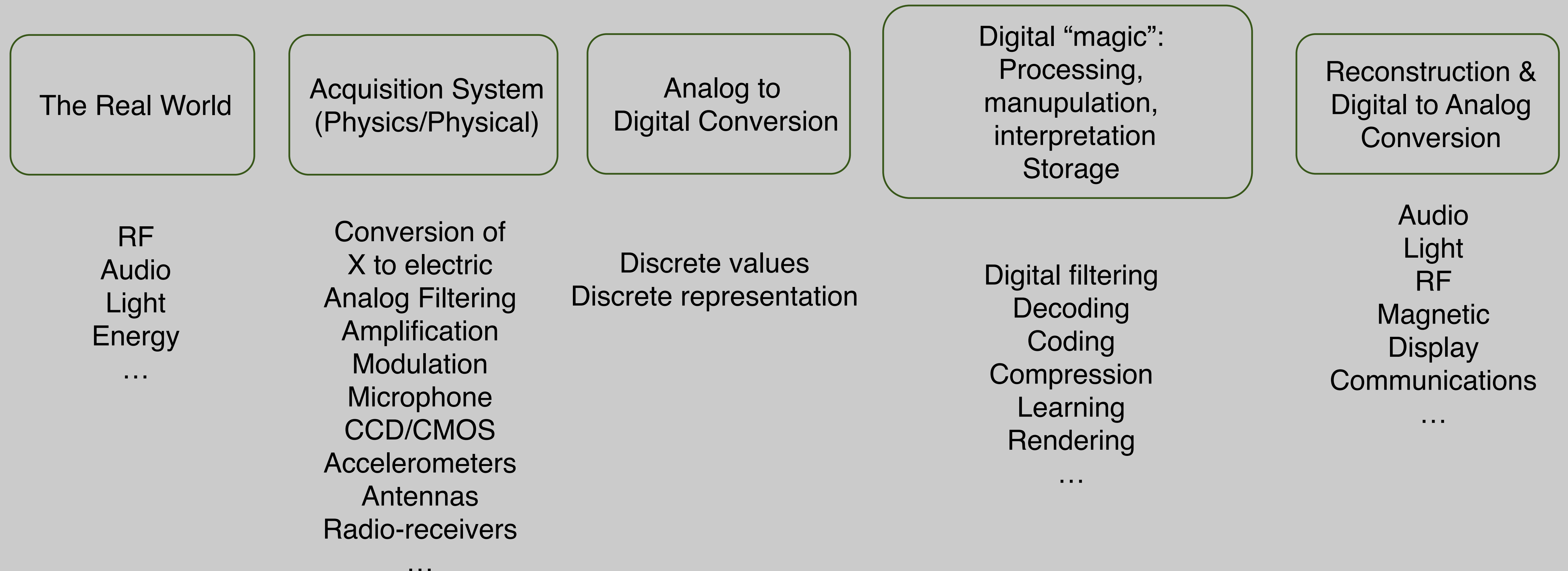
$$A^\dagger = V_1 S^{-1} U_1^T$$

- For tall matrices is the least squares solution!
- For fat matrices is the least-norm solution!
- Computed via the SVD!

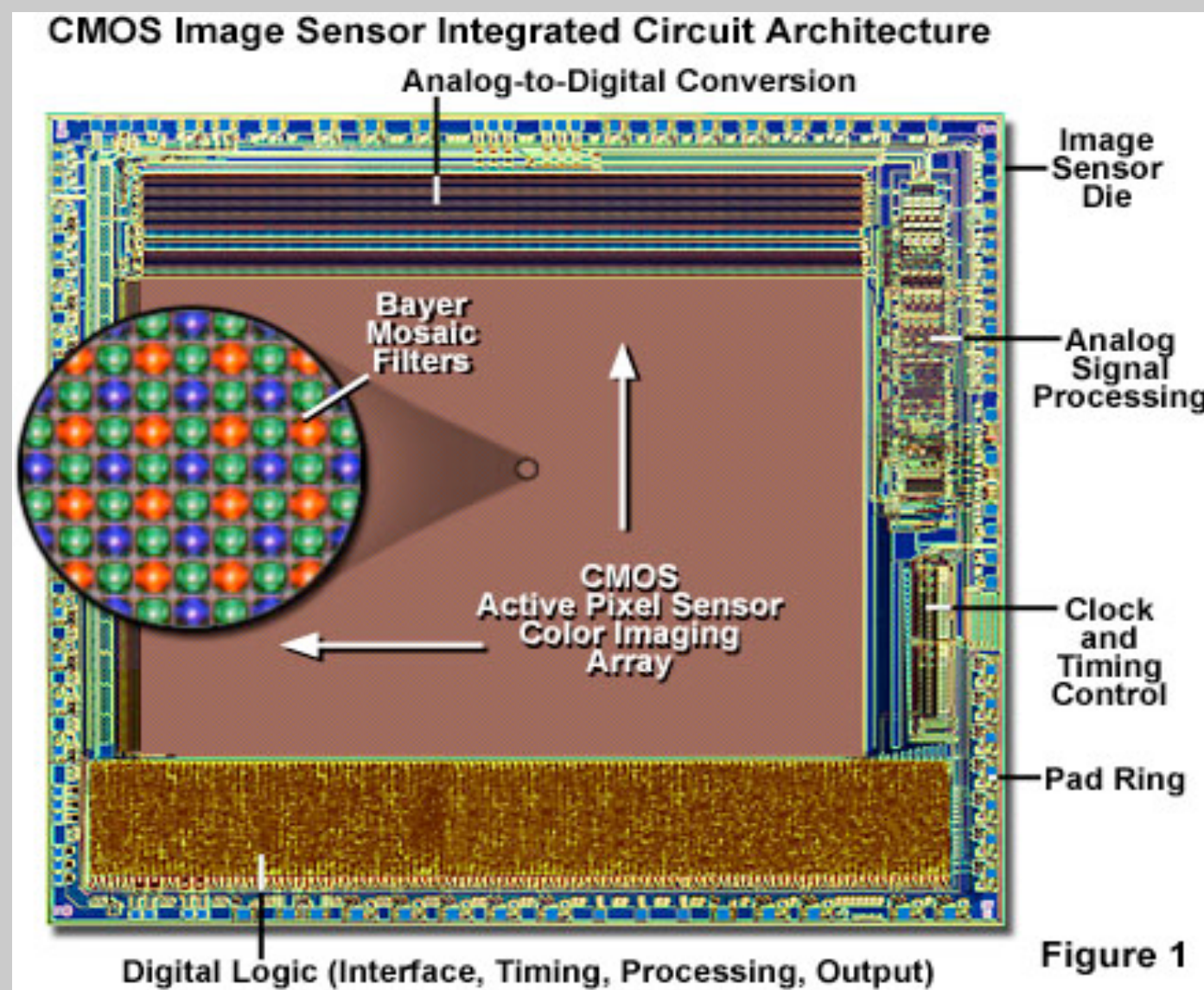


# Sampling and Interpolation

- (Digital) Signal Processing - Only going to touch the surface
  - EE120, EE123, EE145A, EE121, EE225A, EE225B, CS 194-026, CS280



# Example Digital Imaging Camera



Focus/exposure Control

preprocessing

white-balancing

Post-processing

Color transform

demosaic

Compression

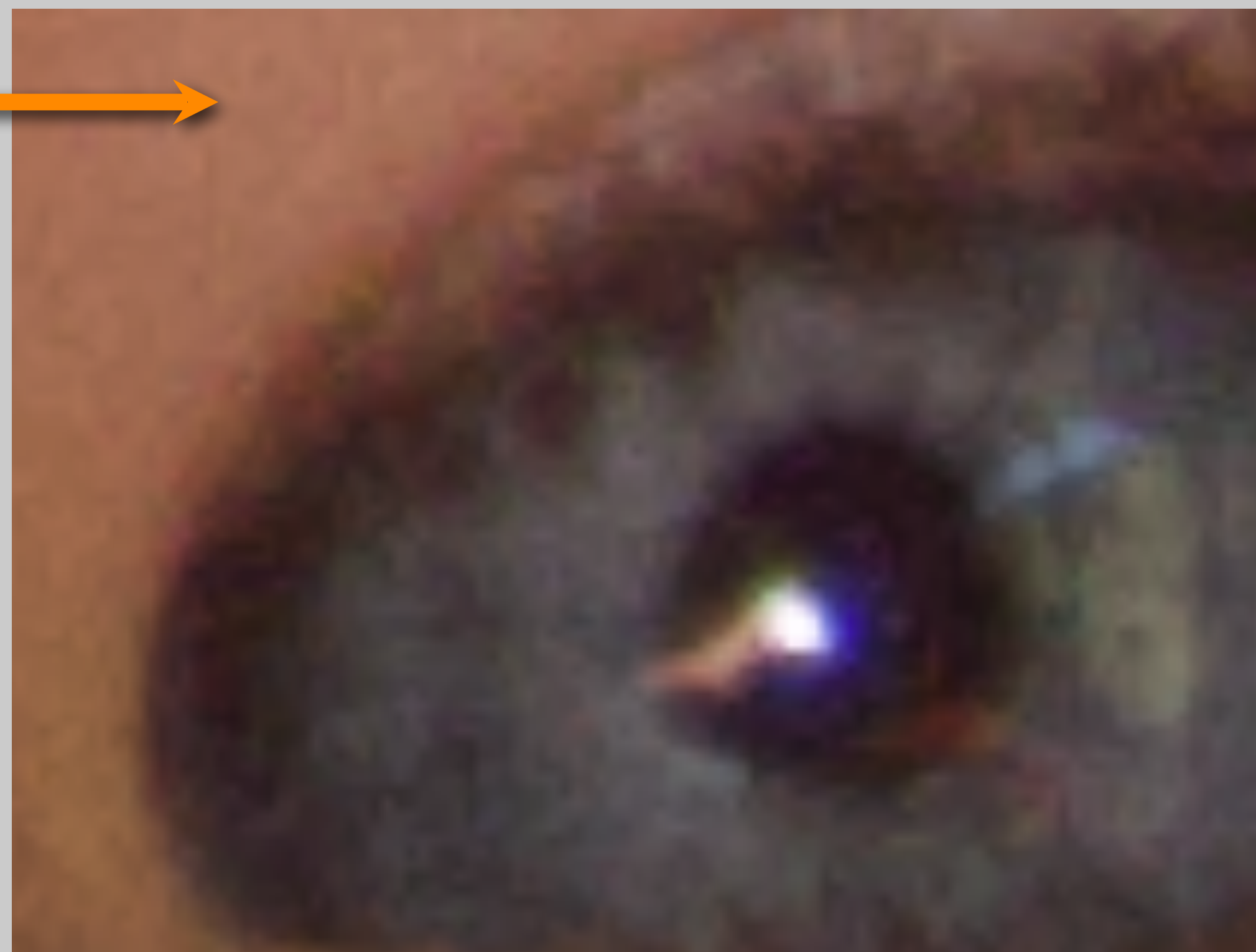
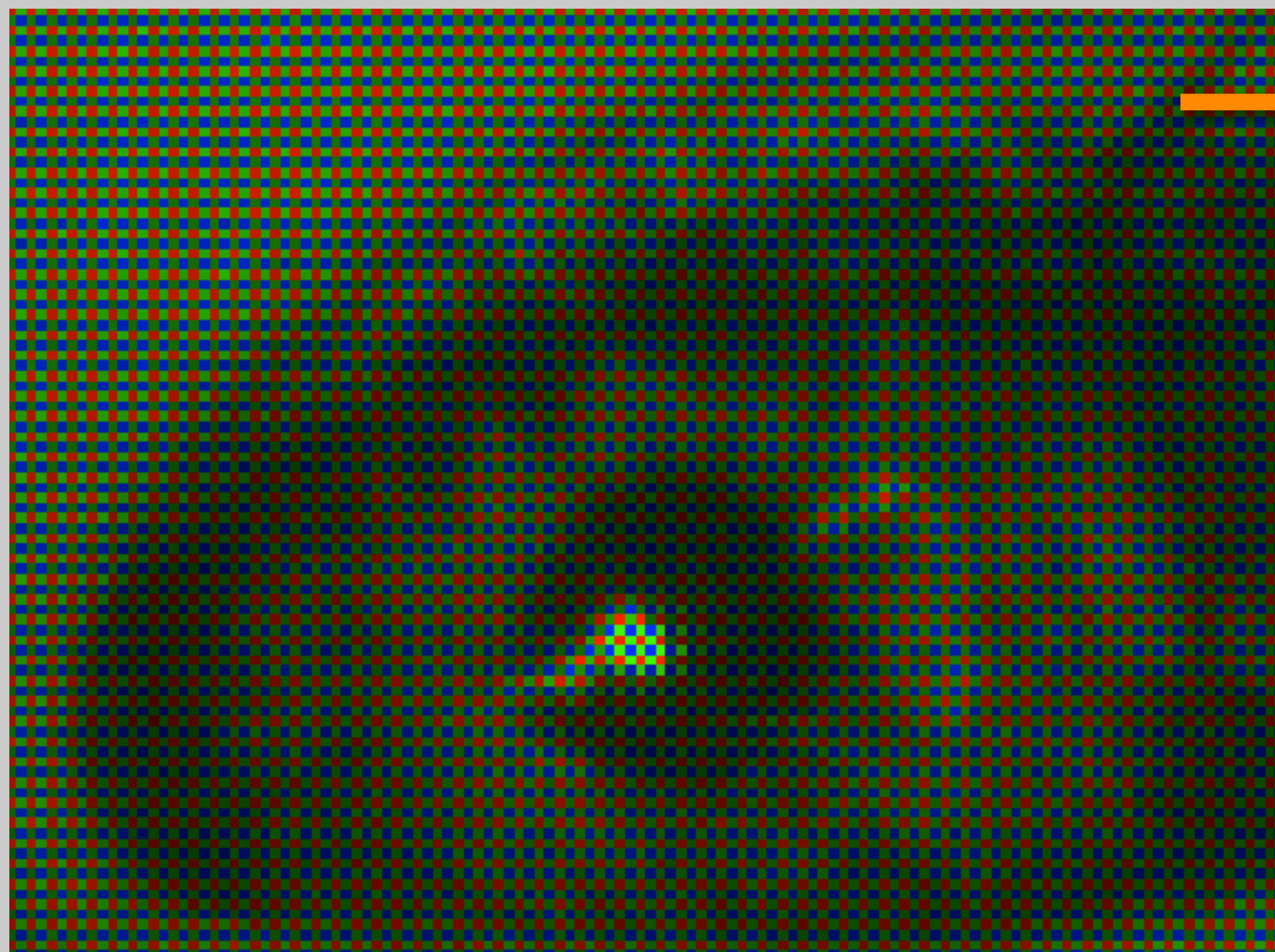
<http://micro.magnet.fsu.edu/primer/digitalimaging/cmosimagesensors.html>

# Example: Digital Camera

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DSP

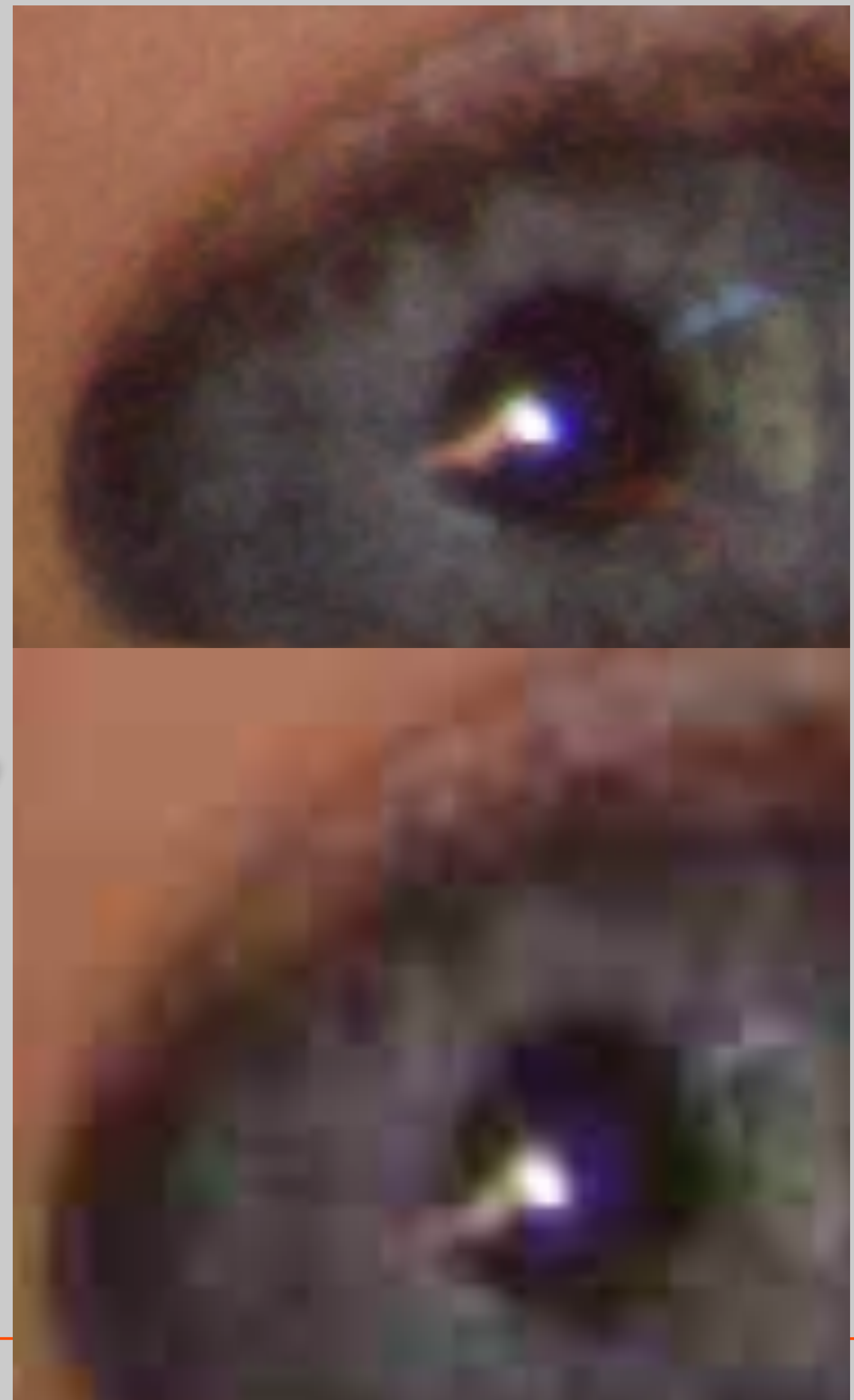


# Example: Digital Camera

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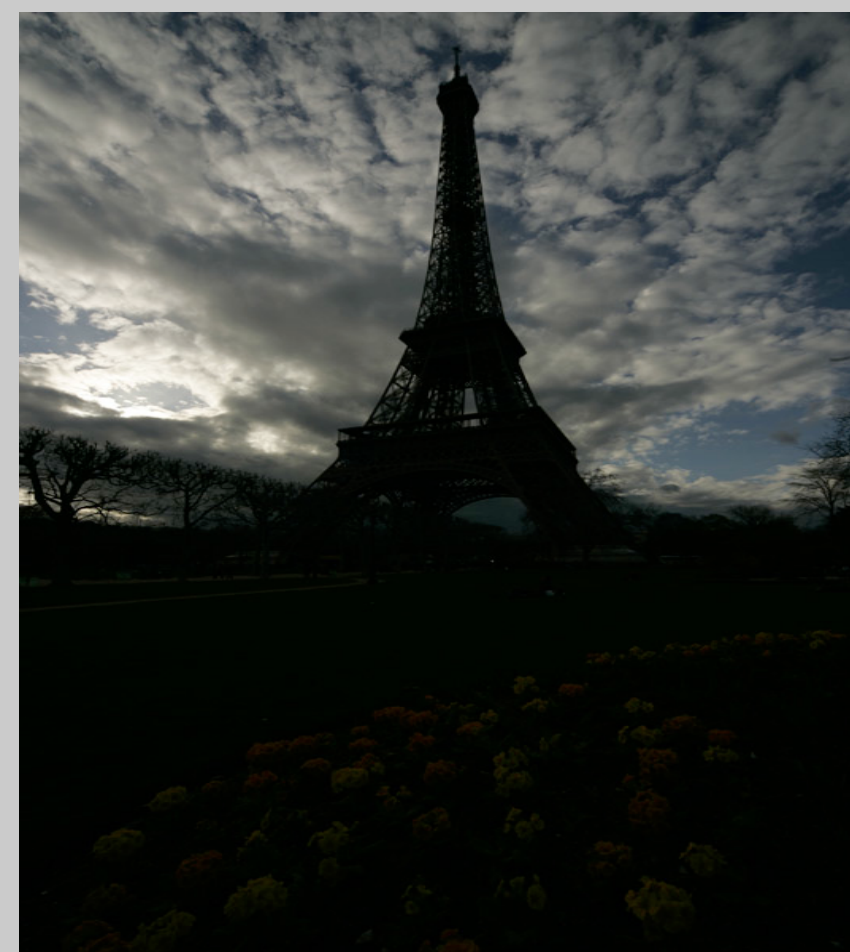
- Compression of 40x without perceptual loss of quality.
- Example of slight overcompression: difference enables x60 compression!

DSP





# Computational Photography



DSP



Implemented in all smart phones (HDR)

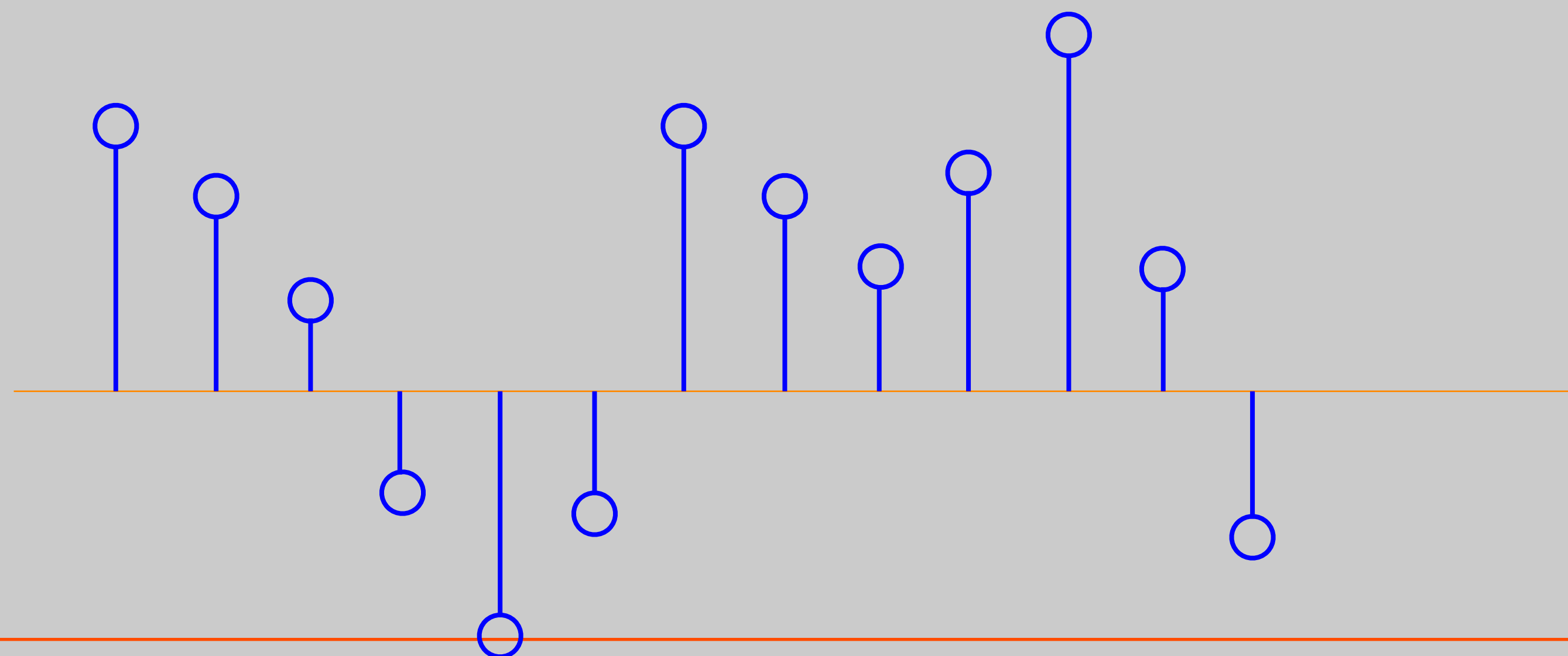
\*[www.hdrsoft.com](http://www.hdrsoft.com)

# Interpolation

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- Given data points  $(x_i, y_i)$   $i=1,2,\dots, n$   
find a continuous function that exactly matches the points.

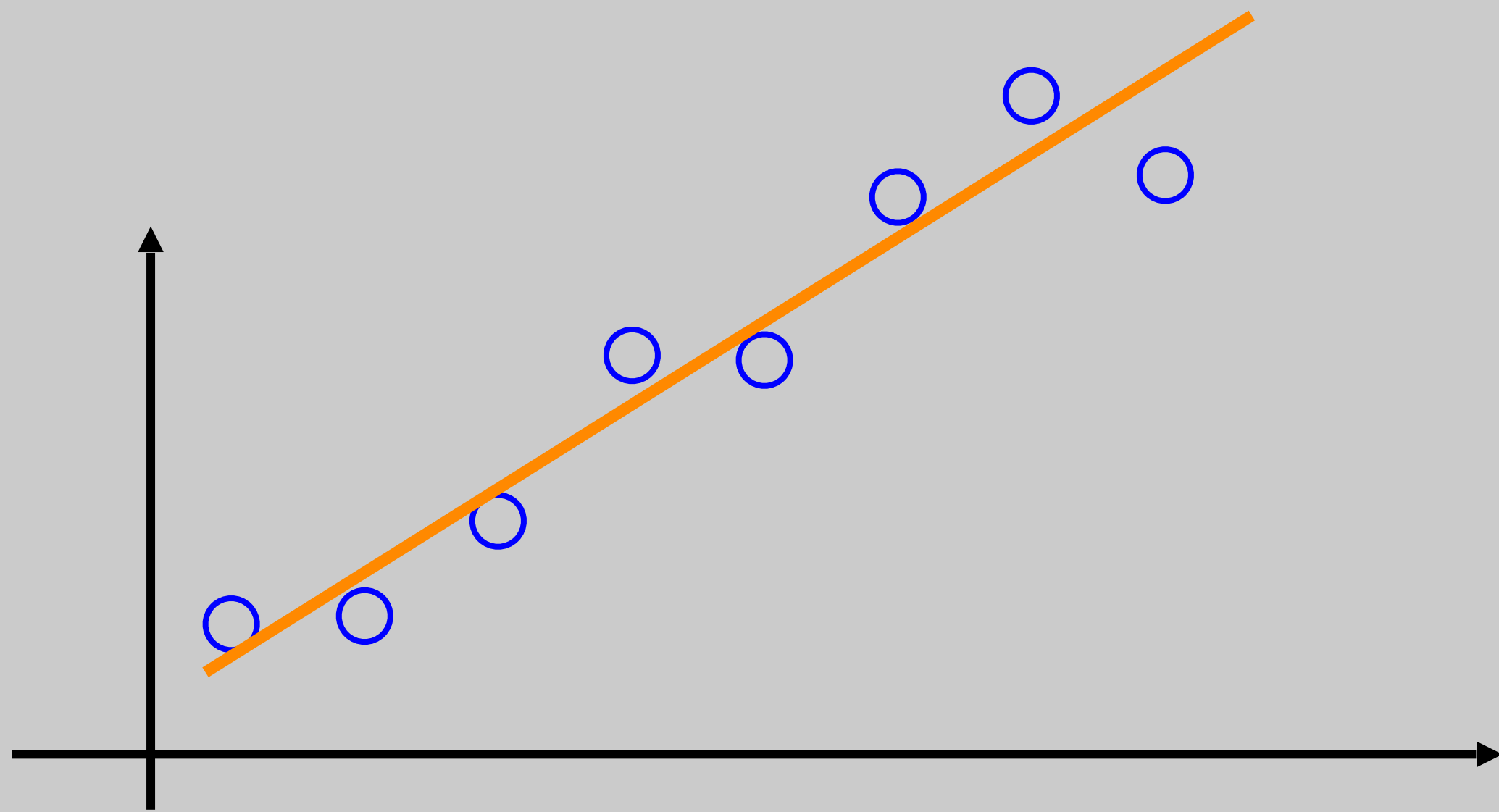
$$y = f(x) \quad \Rightarrow \quad f(x_i) = y_i, \quad i = 1, 2, 3, \dots, n$$



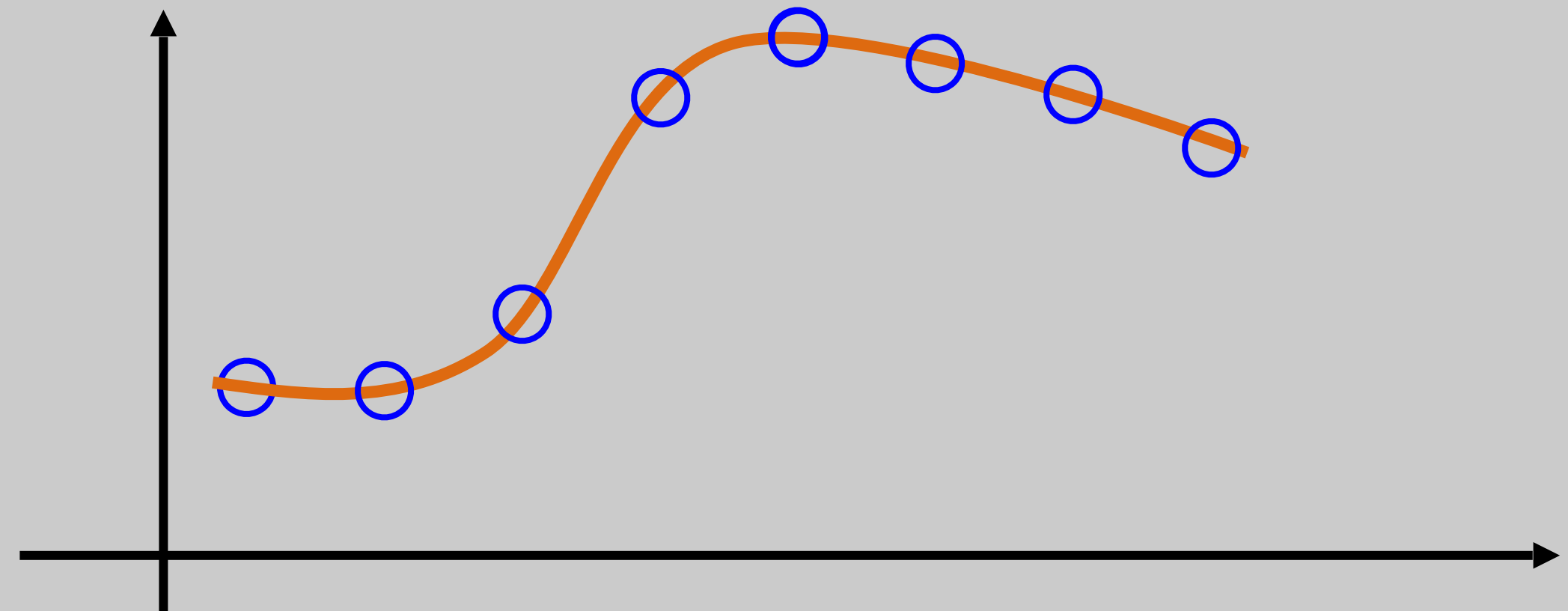
# Regression Vs Interpolation

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regression

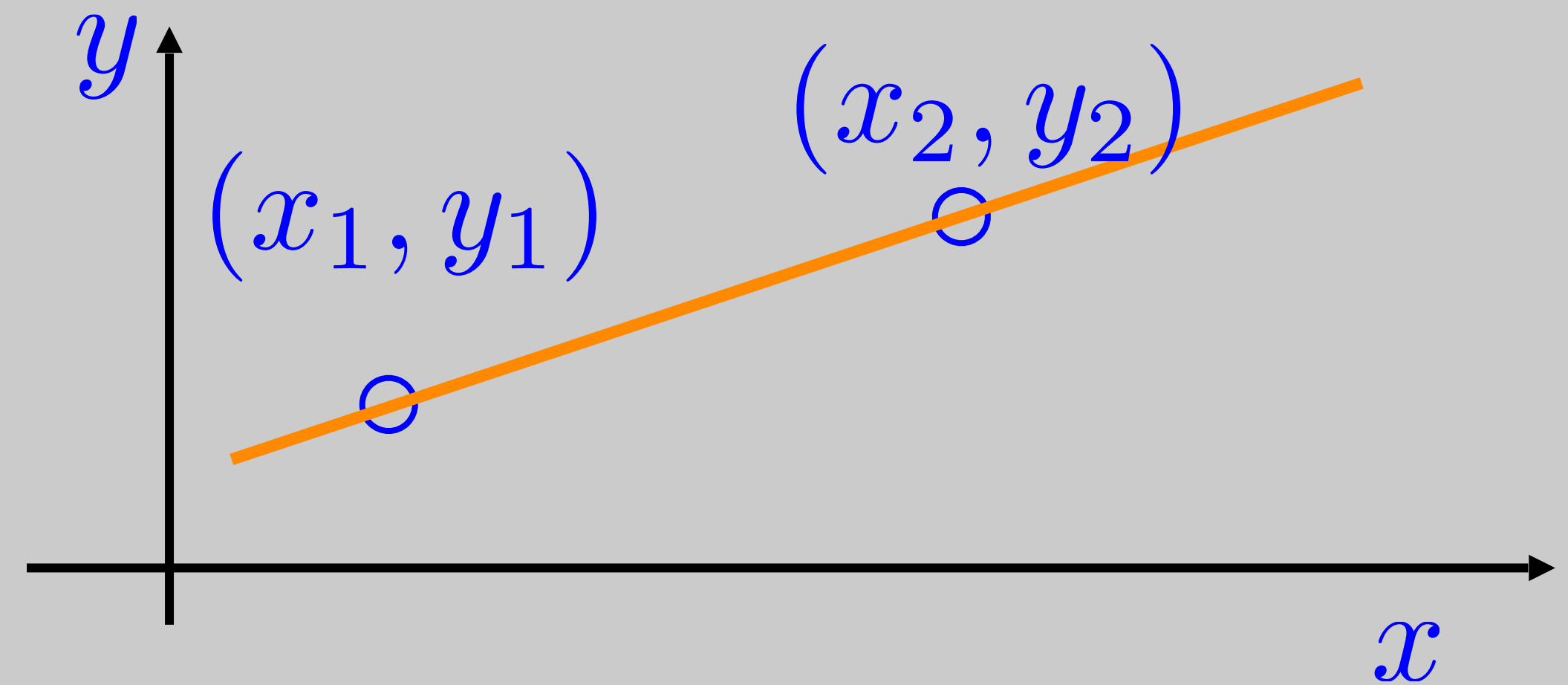


interpolation



# Polynomial Interpolation

- Assumption:
  - Data are samples of a polynomial function (smooth)
  - Lowest order polynomial that exactly fits points



$$y = a_0 + a_1x$$

$$a_0 + a_1x_1 = y_1$$

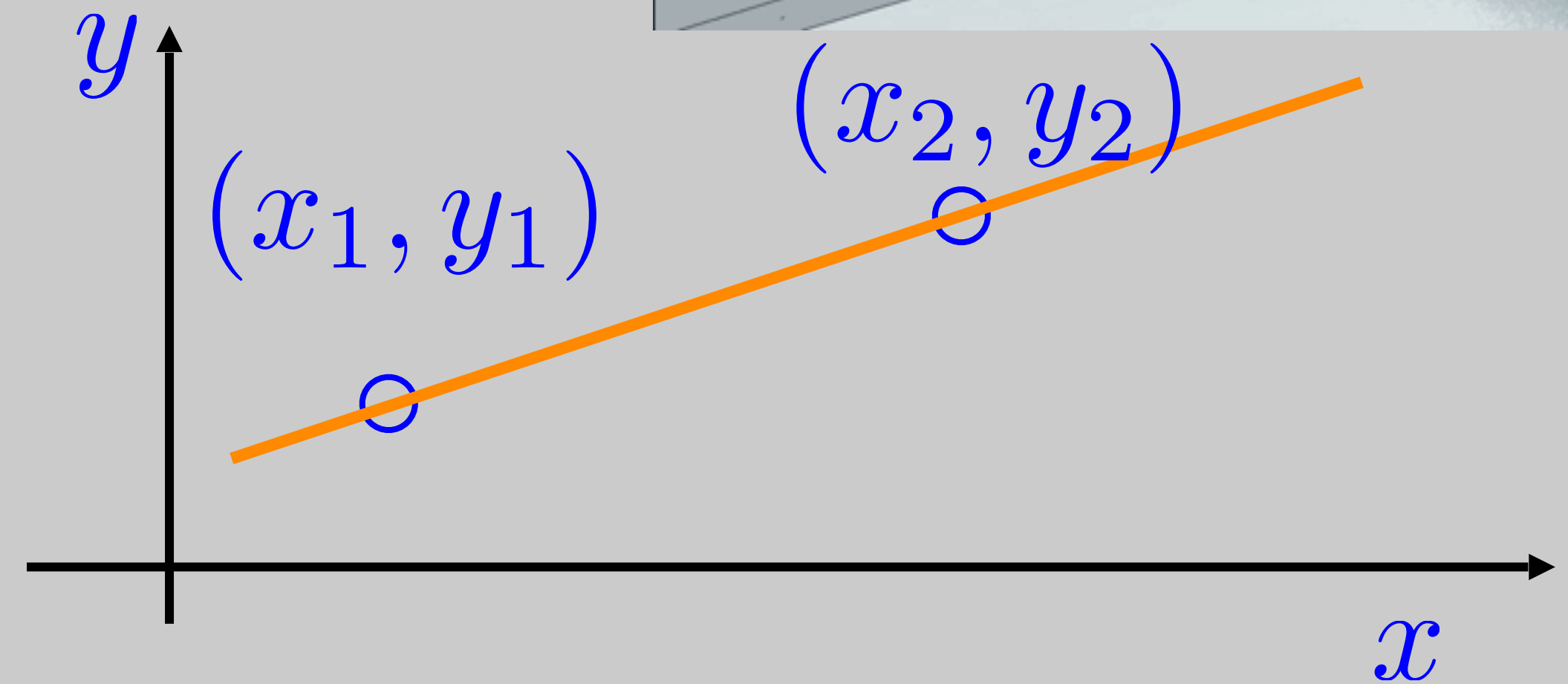
$$a_0 + a_1x_2 = y_2$$

$$\rightarrow \begin{bmatrix} \phantom{0} \\ \phantom{0} \end{bmatrix} \begin{bmatrix} \phantom{0} \\ \phantom{0} \end{bmatrix} = \begin{bmatrix} \phantom{0} \\ \phantom{0} \end{bmatrix}$$



# Polynomial Interpolation

- Assumption:
  - Data are samples of a polynomial function (smooth)
  - Lowest order polynomial that exactly fits points



$$y = a_0 + a_1 x$$

$$\begin{aligned} a_0 + a_1 x_1 &= y_1 \\ a_0 + a_1 x_2 &= y_2 \end{aligned} \rightarrow \begin{bmatrix} 1 & x_1 \\ 1 & x_2 \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$$

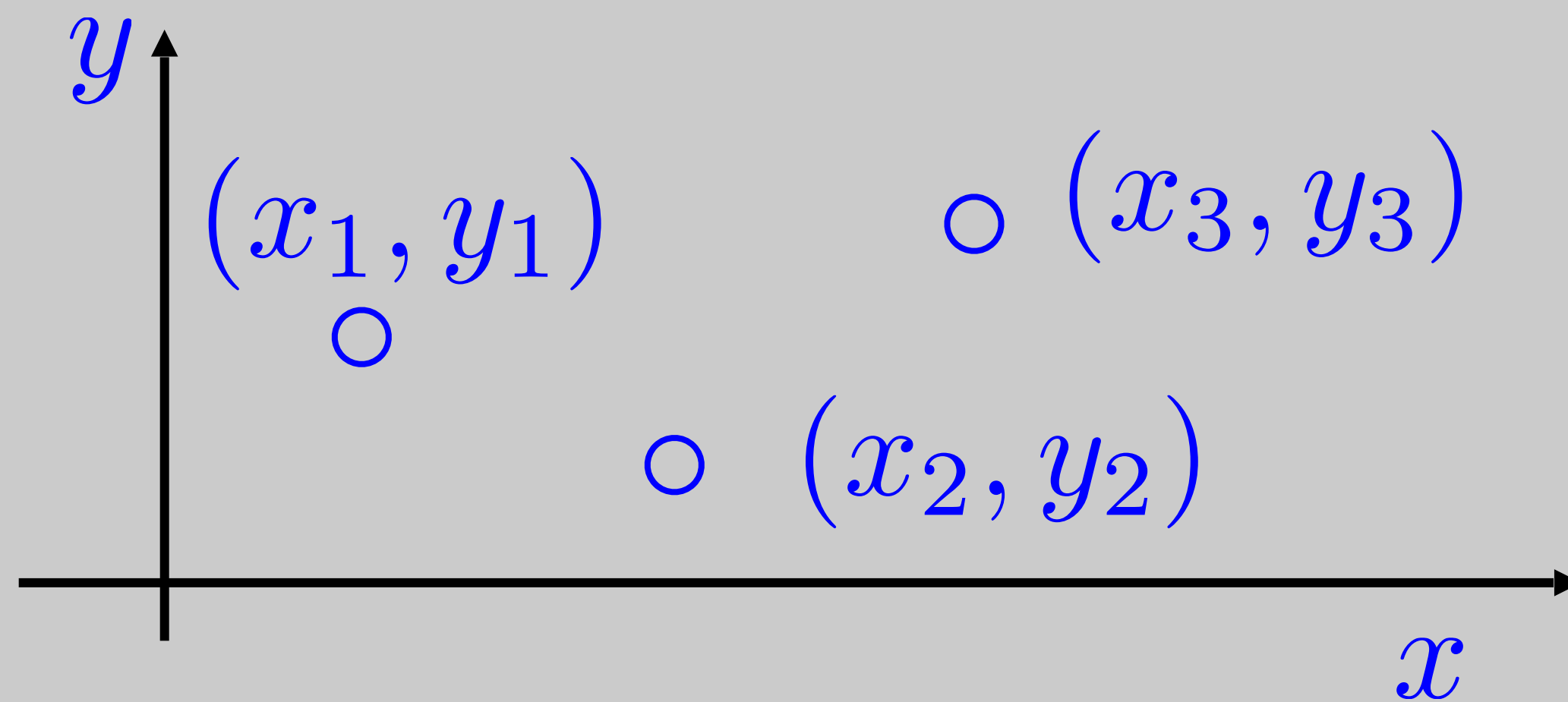
Invertible if  $x_1 \neq x_2$

# Polynomial Interpolation

$(x_1, y_1) (x_2, y_2) (x_3, y_3)$

$$y = a_0 + a_1x + a_2x^2 \rightarrow \begin{bmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}$$

$$x_1 \neq x_j$$



# Polynomial Interpolation

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- Given  $n$  distinct points, then there's exist a unique  $(n-1)$  order polynomial passing through them

$$y = a_0 + a_1x + a_2x^2 + \cdots + a_{n-1}x^{n-1}$$

$$\rightarrow \begin{bmatrix} \phantom{0} \\ \phantom{0} \\ \phantom{0} \\ \phantom{0} \\ \phantom{0} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}$$

# Polynomial Interpolation

- Given  $n$  distinct points, then there's exist a unique  $(n-1)$  order polynomial passing through them

$$y = a_0 + a_1x + a_2x^2 + \cdots + a_{n-1}x^{n-1}$$

$$\rightarrow \begin{bmatrix} 1 & x_1 & x_1^2 & \cdots & x_1^{n-1} \\ 1 & x_2 & x_2^2 & \cdots & x_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^{n-1} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}$$

“Vandermonde” Matrix  $\det(v) = \prod_{1 \leq i < j \leq n} (x_j - x_i)$

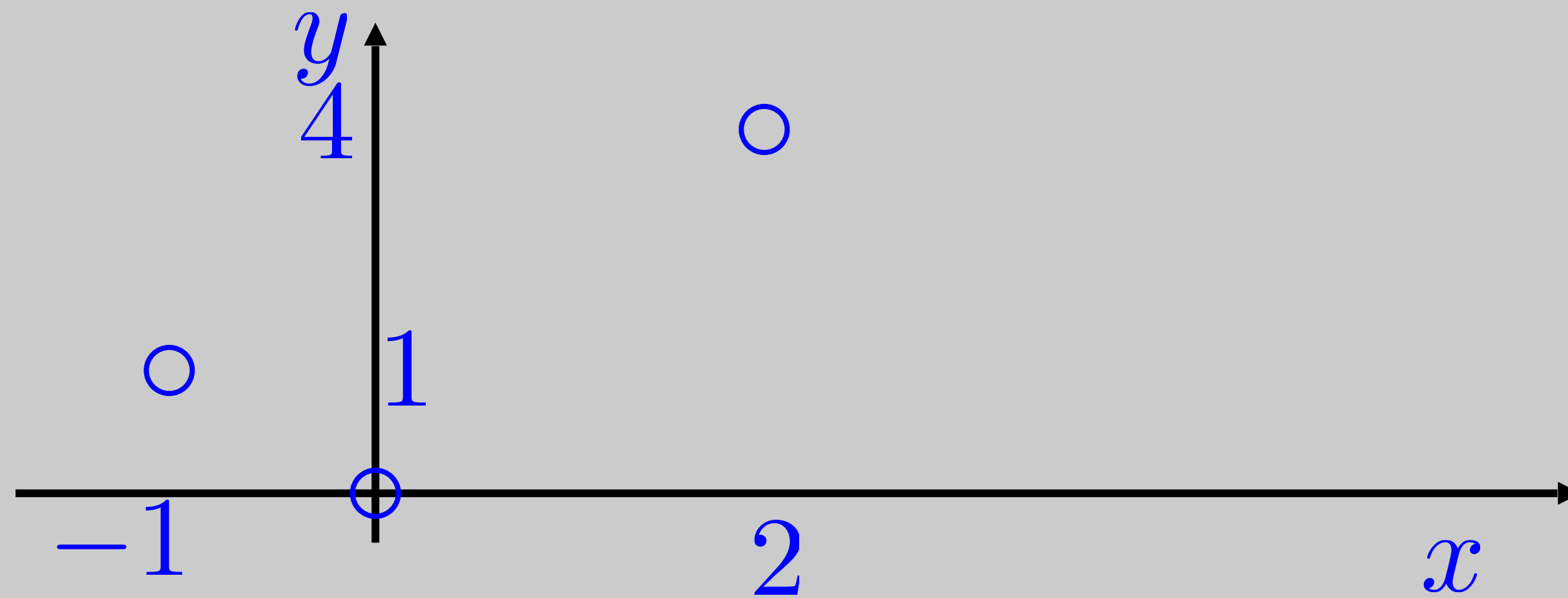


# Quiz

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- What's the polynomial that passes through these points:

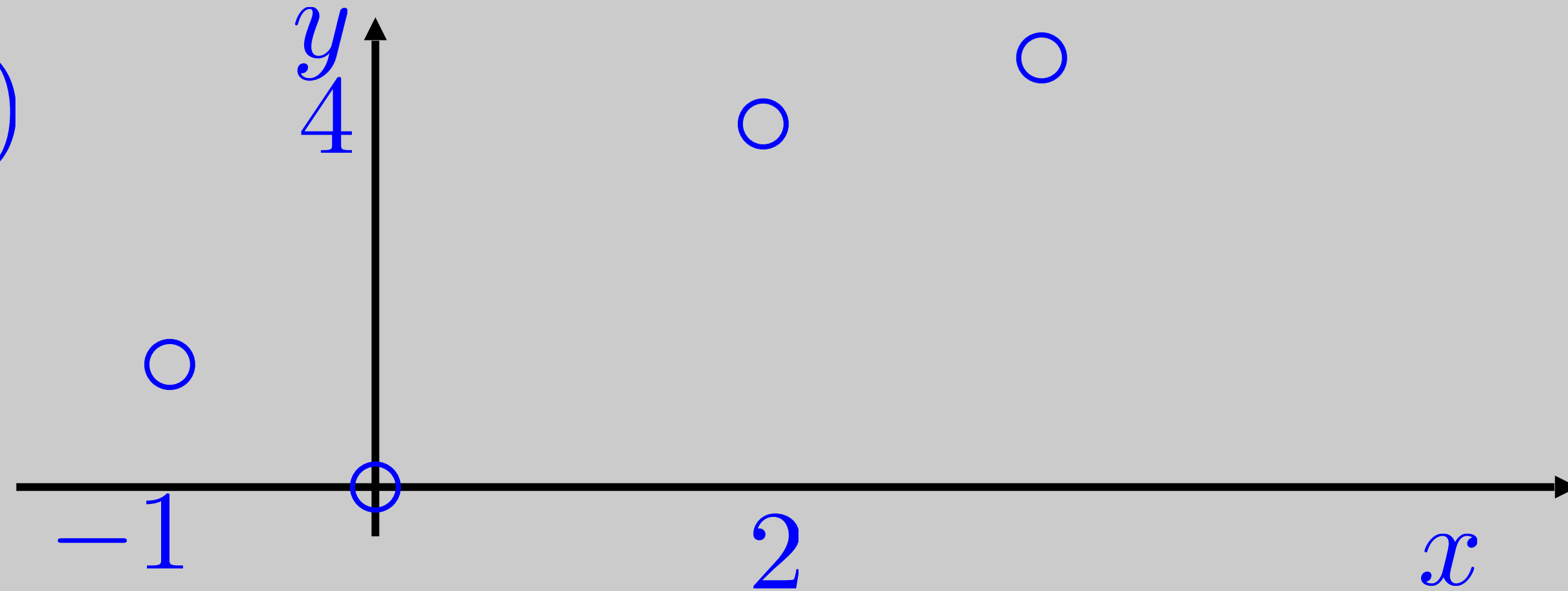
$$(-1, 1), (0, 0), (2, 4)$$



# Polynomial Regression

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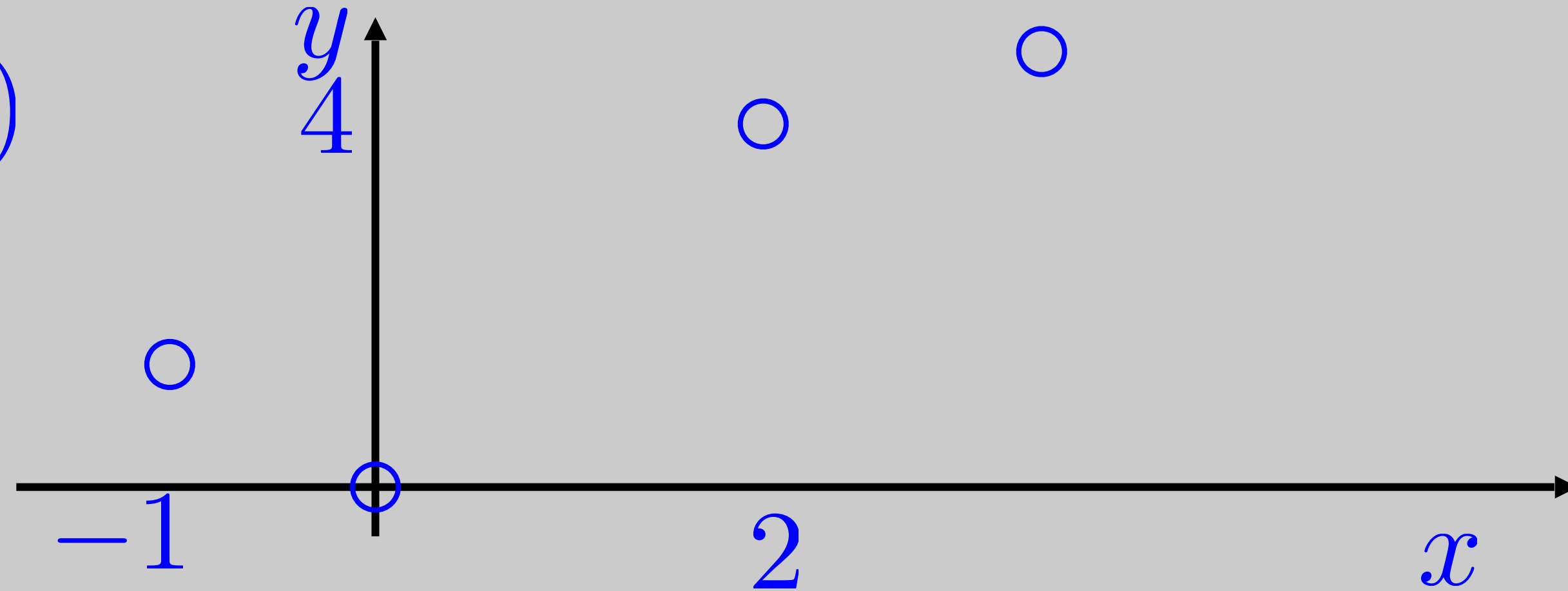
$(-1, 1), (0, 0), (2, 4), (3, 5)$



- What is the “best” quadratic polynomial that passes through the points?

# Polynomial Regression

$(-1, 1), (0, 0), (2, 4), (3, 5)$



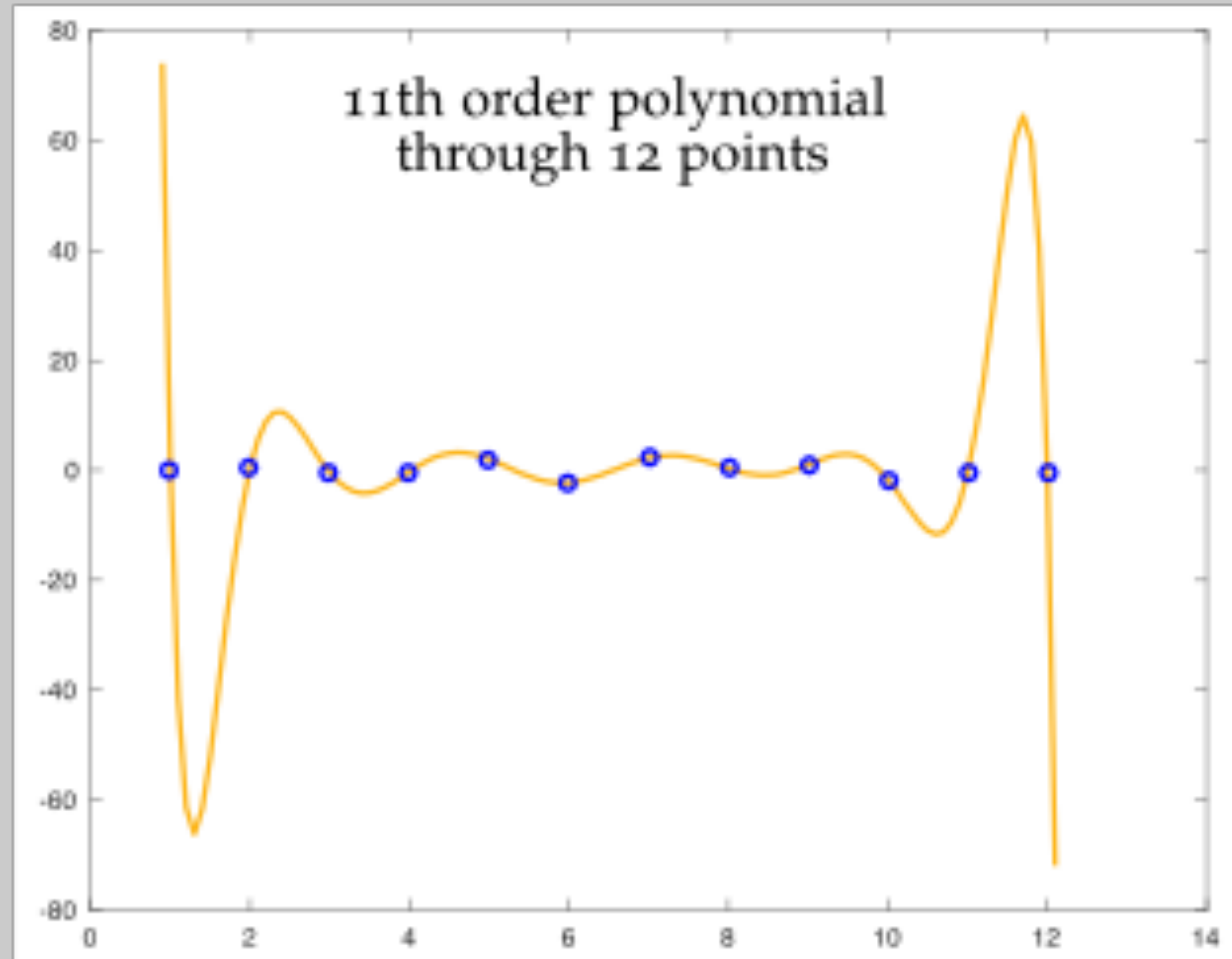
What is the “best” quadratic polynomial that passes through the points?

$$y = a_0 + a_1x + a_2x^2 + \cdots + a_{n-1}x^{n-1}$$

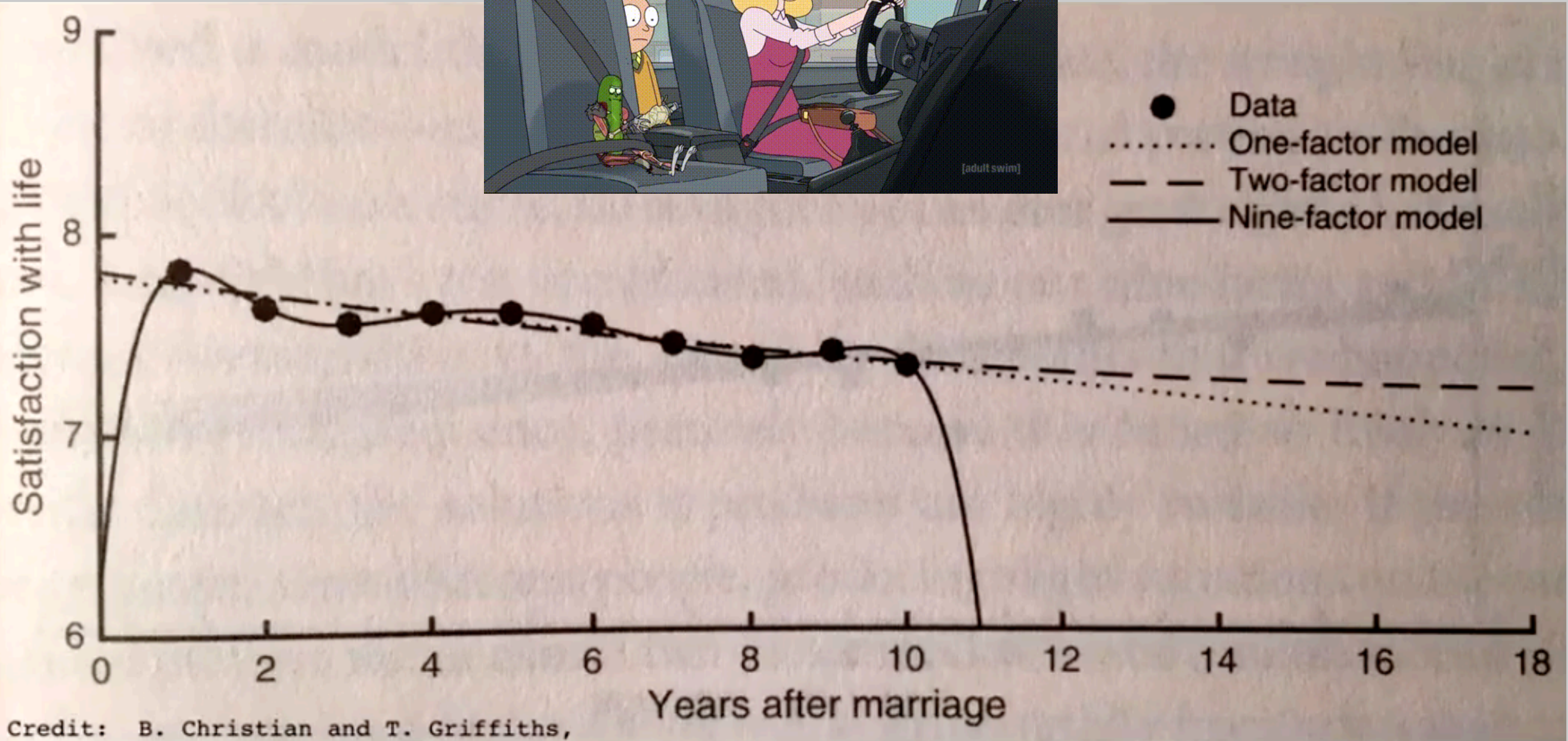
$$\begin{bmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \\ 1 & x_4 & x_4^2 \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \end{bmatrix}$$

# Issue with Polynomial Interpolation

- Tend to be oscillatory in high order interp/regression
- Not numerically stable!  $x^{n-1}$



# Example



Credit: B. Christian and T. Griffiths,