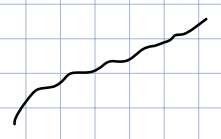


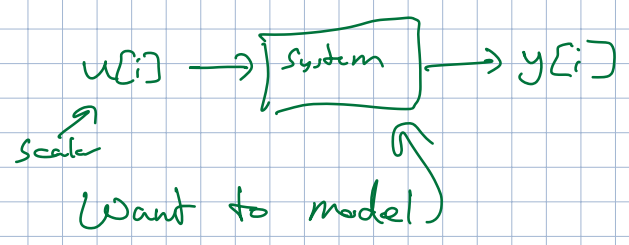
Today: Orthogonalization

Midterm last night, but now redo it like HW.  
 Can collaborate. Go to Bio Party. Due Friday.  
 Read solutions after redo due, and self-grade redo.  
 Can resubmit after seeing solutions to get back redo pts.  
 If redo complete (i.e. Hit 80% after resubmit), unless checker.  
 Our Goal: Your Mastery



Related to HW 6  
 Demo Systemid problem  
 Example 4.

Motivating Example:



The model could be:  $y[i+1] = b u[i]$

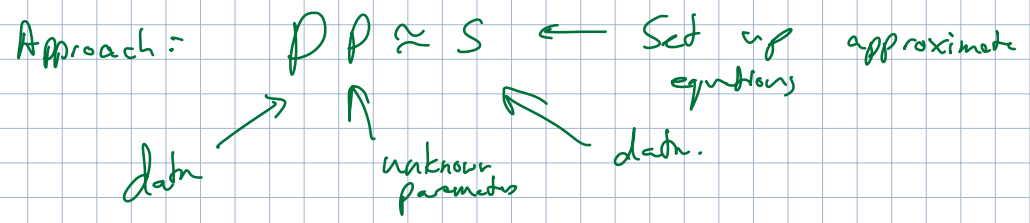
Called  
 auto-regressive  
 models  
 or Markov of order k.

- or  $y[i+1] = b u[i] + a_0 y[i]$
- or  $y[i+1] = b u[i] + a_0 y[i] + a_1 y[i-1]$
- or  $y[i+1] = b u[i] + a_0 y[i] + a_1 y[i-1] + a_2 y[i-2]$
- or  $\vdots$
- $\therefore y[i+1] = b u[i] + \sum_{j=0}^{k-1} a_j y[i-j]$

Want to do System-ID.

- a) Try to fit all the models
  - b) Test them to see which one works best.
- Gold Standard: works in practice  
 Silver Standard: predicts well on test data.  
 "hold-out set"

To do (a), need to solve a family of nested least-squares problems



$$S = \begin{bmatrix} y[1] \\ y[2] \\ \vdots \\ y[i] \end{bmatrix} \quad P = \begin{bmatrix} b \\ \vdots \\ \vdots \\ b \end{bmatrix} \text{ or } \begin{bmatrix} b \\ a_0 \\ \vdots \\ a_1 \end{bmatrix} \dots$$

$$D = \begin{bmatrix} \vec{d}_1 & \vec{d}_2 & \dots & \vec{d}_k \end{bmatrix}$$

# matches dim  $\rho$

Other places this can happen:

1) Polynomial fitting.

2) OMP

LS Problem 0:

$$\begin{bmatrix} \vec{d}_1 \end{bmatrix} \vec{p}_1 \approx \vec{s}$$

LS Problem 1

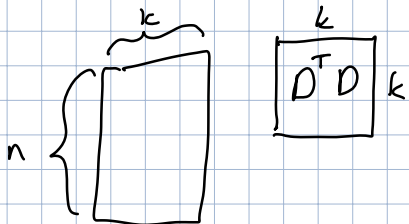
$$\begin{bmatrix} \vec{d}_1 & \vec{d}_2 \end{bmatrix} \vec{p}_2 \approx \vec{s}$$

LS Problem 2

$$\begin{bmatrix} \vec{d}_1 & \vec{d}_2 & \vec{d}_3 \end{bmatrix} \vec{p}_3 \approx \vec{s}$$

How do you solve least-squares?

$$D \vec{p} \approx \vec{s} \implies \hat{\vec{p}} = (D^T D)^{-1} D^T \vec{s}$$



↑  $k \times k$   
keeps growing.

Inverting takes  $O(k^3)$  steps.

Gets painful.

Can we make things faster?

To do  $l$  nested least-sq problems

Can we solve a whole nested family of LS problems faster.

$$1 + 2^3 + 3^3 + \dots + l^3 = \frac{l^2(l+1)^2}{4} = O(l^4)$$

Thoughts: (1)  $(D^T D)^{-1}$  is the painful step. When is this fast?

(2) Can we change coordinates somehow?

$$D^T D = \begin{bmatrix} \vec{d}_1^T \\ \vec{d}_2^T \\ \vdots \\ \vec{d}_k^T \end{bmatrix} \begin{bmatrix} \vec{d}_1 & \vec{d}_2 & \dots & \vec{d}_k \end{bmatrix} = \begin{bmatrix} \vec{d}_1^T \vec{d}_1 & \vec{d}_1^T \vec{d}_2 & \dots & \vec{d}_1^T \vec{d}_k \\ \vdots & \vdots & \ddots & \vdots \\ \vec{d}_k^T \vec{d}_1 & \vec{d}_k^T \vec{d}_2 & \dots & \vec{d}_k^T \vec{d}_k \end{bmatrix}$$

$\vec{d}_i \in \mathbb{R}^n \leftarrow n\text{-dim vectors}$

Matrix of pairwise inner products

$$\langle \vec{d}_i, \vec{d}_j \rangle = \vec{d}_i^T \vec{d}_j$$

When is  $D^T D$  the identity?

$$\vec{d}_i^T \vec{d}_i = 1 \leftarrow \text{diagonal entries of}$$

$$\vec{d}_i^T \vec{d}_j = 0 \text{ if } i \neq j \text{ aka. } \langle \vec{d}_j, \vec{d}_i \rangle = 0$$

ie. orthogonal or perpendicular

$$\|\vec{d}_i\|^2 = 1$$

$$\|\vec{d}_i\| = 1 \leftarrow \text{normalized} = \sqrt{(d_i[x])^2 + (d_i[y])^2 + \dots + (d_i[z])^2}$$

Desired property: Orthonormal columns.

$$Q = [\vec{q}_1 \ \vec{q}_2 \ \dots \ \vec{q}_k]$$

Know if  $Q$  is orthonormal matrix then  $Q^T Q = I$

$\hat{=}$  favored letter for a matrix with orthonormal columns.

"Orthonormal matrix"

How can we project onto the span of the columns of  $Q$

Want to find  $\vec{x}$  s.t.

$Q\vec{x}$  is as close as possible to  $\vec{y}$   
 $\hat{y}$  projection of  $\vec{y}$  onto  $\text{span}(Q)$

$$\hat{\vec{x}} = (Q^T Q)^{-1} Q^T \vec{y} = Q^T \vec{y}$$

$$\hat{\vec{y}} = Q \hat{\vec{x}} = \underbrace{Q Q^T}_{\substack{! \\ \text{Projection operator for } Q}} \vec{y}$$

e.g.

$$Q = \begin{bmatrix} 1 & 0 \\ 0 & 1/\sqrt{2} \\ 0 & 1/\sqrt{2} \end{bmatrix}$$

$$Q^T Q = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

$$Q Q^T = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 1/2 \\ 0 & 1/2 & 1/2 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} y_1 \\ 1/2(y_2+y_3) \\ 1/2(y_2+y_3) \end{bmatrix}$$

Given a basis expressed as an orthonormal  $Q$   
for a subspace

Projection onto it is very fast.

Goal (Subgoal of making things fast):

Given a sequence of vectors  $\vec{d}_1, \vec{d}_2, \dots, \vec{d}_l$

Return another sequence of vectors  $\vec{q}_1, \vec{q}_2, \dots, \vec{q}_l$

st.  $Q = [\vec{q}_1, \dots, \vec{q}_l]$  is orthonormal

and  $\forall k \leq l \quad \text{span}\{\vec{d}_1, \dots, \vec{d}_k\} = \text{span}\{\vec{q}_1, \dots, \vec{q}_k\}$ .

Orthonormalization is the name given to this process.

Start with simplest case.  $\{\vec{d}_1\} \leftarrow$  just one vector.

Need  $\vec{q}_1 = \alpha \vec{d}_1 \leftarrow$  so it spans the same subspace.

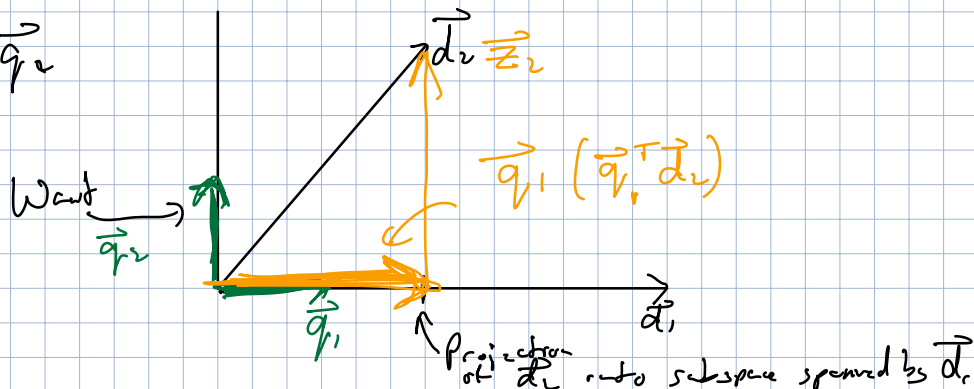
Q: How to pick  $\alpha$ ?

Guess: Divide by norm?

$$\vec{q}_1 = \frac{\vec{d}_1}{\|\vec{d}_1\|} \quad \text{Compute } \vec{q}_1^T \vec{q}_1 = \frac{\vec{d}_1^T \vec{d}_1}{\|\vec{d}_1\| \cdot \|\vec{d}_1\|} = \frac{\|\vec{d}_1\|^2}{\|\vec{d}_1\|^2} = 1$$

First nontrivial case:  $\{\vec{d}_1, \vec{d}_2\}$

Want  $\vec{q}_2$



$$\vec{z}_2 = \vec{d}_2 - \vec{q}_1 (\vec{q}_1^T \vec{d}_2)$$

↑ Parallel to desired  $\vec{q}_2$

$$\text{So let } \vec{q}_2 = \frac{\vec{z}_2}{\|\vec{z}_2\|}$$

Breaks if  $\|\vec{z}_2\| = 0$  or  $\vec{z}_2 = \vec{0}$

Happens only if  $\vec{d}_2$  is a multiple of  $\vec{d}_1$ .

Steps

1) Project  $\vec{d}_2$  onto  $\vec{d}_1$

2) Subtract this projection to get a residual that is  $\perp$

3) Normalize it.

Generalize Given  $\{\vec{d}_1, \dots, \vec{d}_k, \vec{d}_{k+1}\}$

How can we get  $\vec{q}_{k+1}$  given that we already have  $Q_k = [\vec{q}_1, \dots, \vec{q}_k]$

with  $\text{colspan}(Q_k) = \text{colspan}(D_k)$

$$\text{Let } D_k = [\vec{d}_1, \dots, \vec{d}_k]$$

Want  $\vec{z}_{k+1} = \vec{d}_{k+1}$  - Projection of  $\vec{d}_{k+1}$  onto subspace spanned by columns of  $D_k$

$$= \vec{d}_{k+1} - D_k (D_k^T D_k)^{-1} D_k^T \vec{d}_{k+1}$$

Since Projection onto  $\text{colspan}(D_k)$

is the same as projection onto  $\text{colspan}(Q_k)$

Projection Operator for  $\text{colspan } D_k$

$$= \vec{d}_{k+1} - Q_k Q_k^T \vec{d}_{k+1}$$

$$= \vec{d}_{k+1} - \sum_{j=1}^k \vec{q}_j (\vec{q}_j^T \vec{d}_{k+1})$$

Can normalize to set  $\vec{q}_{k+1} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|}$

Is  $Q_{k+1} = [Q_k; \vec{q}_{k+1}]$  orthonormal?

Need to check:  $\|\vec{q}_{k+1}\| = 1$ ? Yes because

Is  $\vec{q}_{k+1} \perp q_i$  for  $i=1, \dots, k$

Check: 
$$\vec{q}_i^T \vec{q}_{k+1} = \frac{\vec{q}_i^T \vec{z}_{k+1}}{\|\vec{z}_{k+1}\|} = \frac{1}{\|\vec{z}_{k+1}\|} \left( \vec{q}_i^T \vec{d}_{k+1} - \sum_{j=1}^k \vec{q}_i^T \vec{q}_j (\vec{q}_j^T \vec{d}_{k+1}) \right)$$

$$= \frac{1}{\|\vec{z}_{k+1}\|} \left( \vec{q}_i^T \vec{d}_{k+1} - 0 - 0 - 0 - \vec{q}_i^T \vec{d}_{k+1} - 0 - 0 - 0 \right)$$

$$= 0 \quad \text{So } \vec{q}_{k+1} \perp \vec{q}_i \text{ for } i=1, \dots, k.$$



Subgoal Achieved: Gram-Schmidt Orthonormalization

Is this actually any faster?!

New Way

How many operations to set  $\vec{q}_k$ ?

$\textcircled{n} + \textcircled{kn}$   
 ↑  
 For normalization to sub  $\vec{z}_k$

For  $k=1, \dots, l$

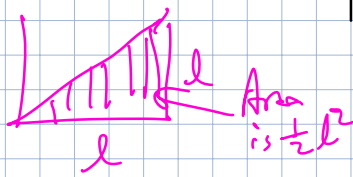
So total cost of setting  $\vec{q}_1, \dots, \vec{q}_l$  is  $O(l^2 n)$

$$= \vec{d}_{k+1} - \sum_{j=1}^k \vec{q}_j (\vec{q}_j^T \vec{d}_{k+1})$$


to set  $\vec{q}_{k+1} = \vec{z}_{k+1} / \|\vec{z}_{k+1}\|$

Recall  $\vec{d}_i$  are  $n$ -dim as an  $\vec{q}_i$

Because



What about the old way?  $(D^T D)^T D^T \vec{y}$   
For  $D$  being  $l$

$$\underbrace{l^3}_{\text{From } (D^T D)^T} + \underbrace{l^2 n}_{\text{to compute } D^T D} + l n$$


$n > l$ . So Both ways for the  $l$ -dim case  
alone cost the same.

Savings exist relative to the  $l^4$  term  
is computing all the LS. problem

Get significant speedup-

$$(100)^4 = 10^8 \leftarrow 100 \text{ million}$$

$$(100)^3 = 10^6 \leftarrow 1 \text{ million.}$$

Final Approach: To solve nested least square problems:

1) Orthonormalize the columns  $[\vec{d}_1, \dots, \vec{d}_l]$   
to get  $Q = [\vec{q}_1, \dots, \vec{q}_l]$

2) Compute  $Q^T \vec{y}$  to get solutions to all nested problems  
at once.

So in the new coordinates: to see the solution for just  $k$  cols,  
look at  $\vec{q}_1^T \vec{y}, \dots, \vec{q}_k^T \vec{y}$ .

This gives you a projection  $\sum_{i=1}^k \vec{q}_i (\vec{q}_i^T \vec{y})$   
that best approximates  $\vec{y}$  in the subspace spanned by  
 $\vec{d}_1, \dots, \vec{d}_k$ .